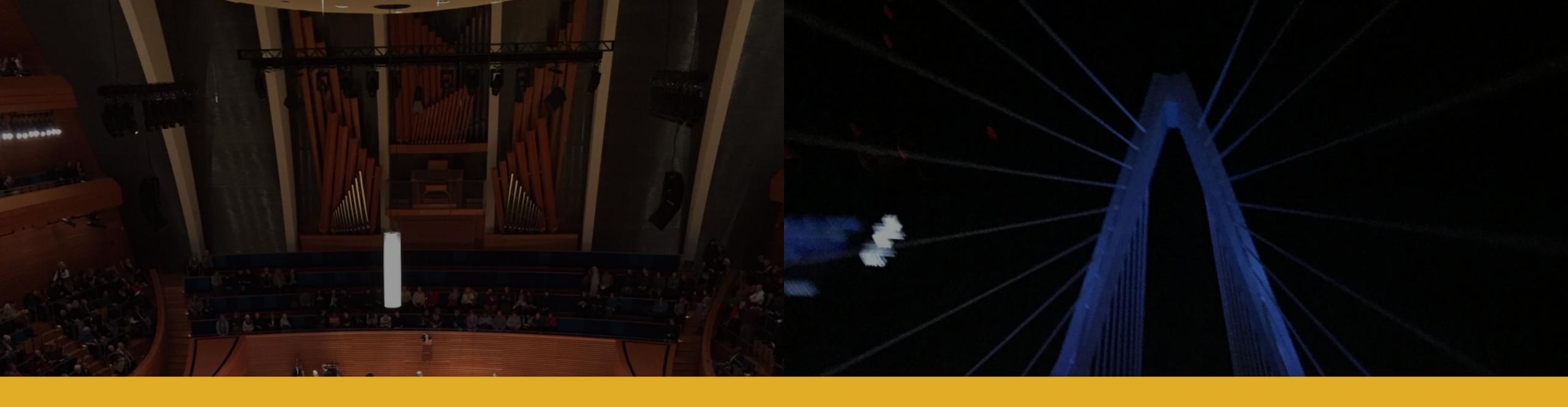
GETTING THE WORD OUT:

MEMBERSHIP, DISSEMINATION & POPULATION PROTOCOLS

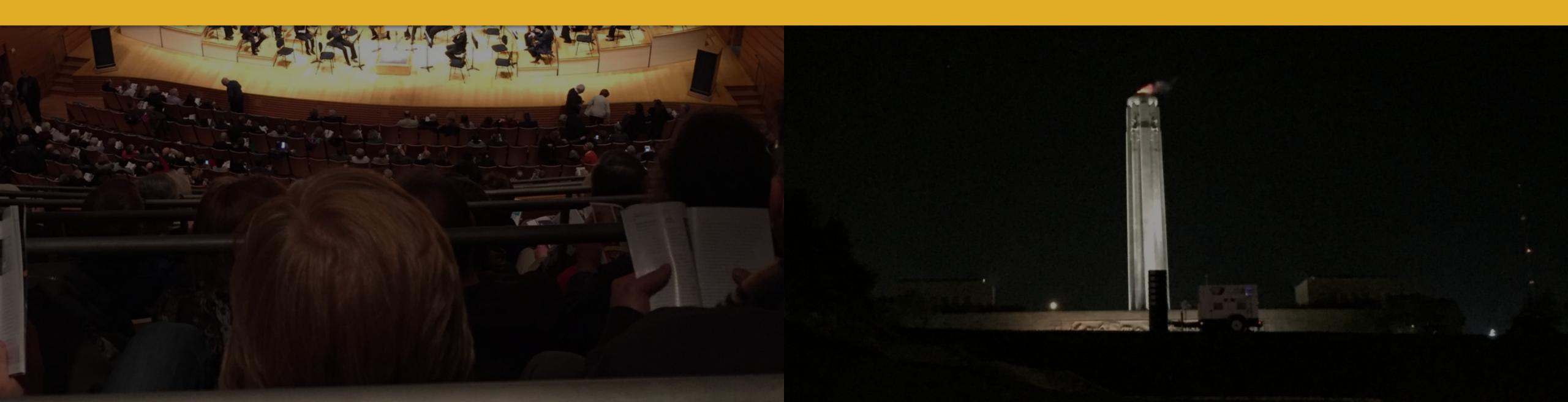




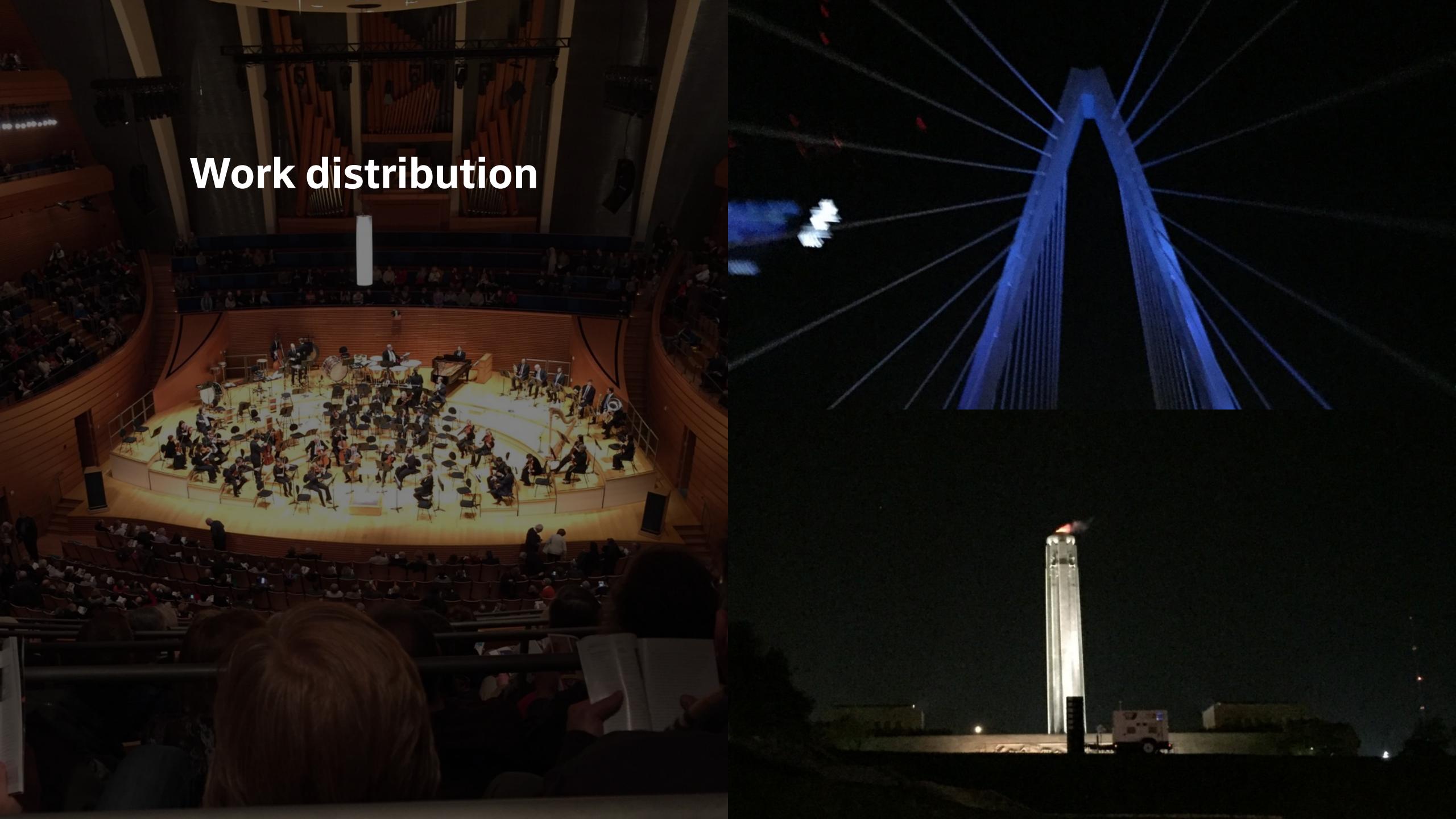


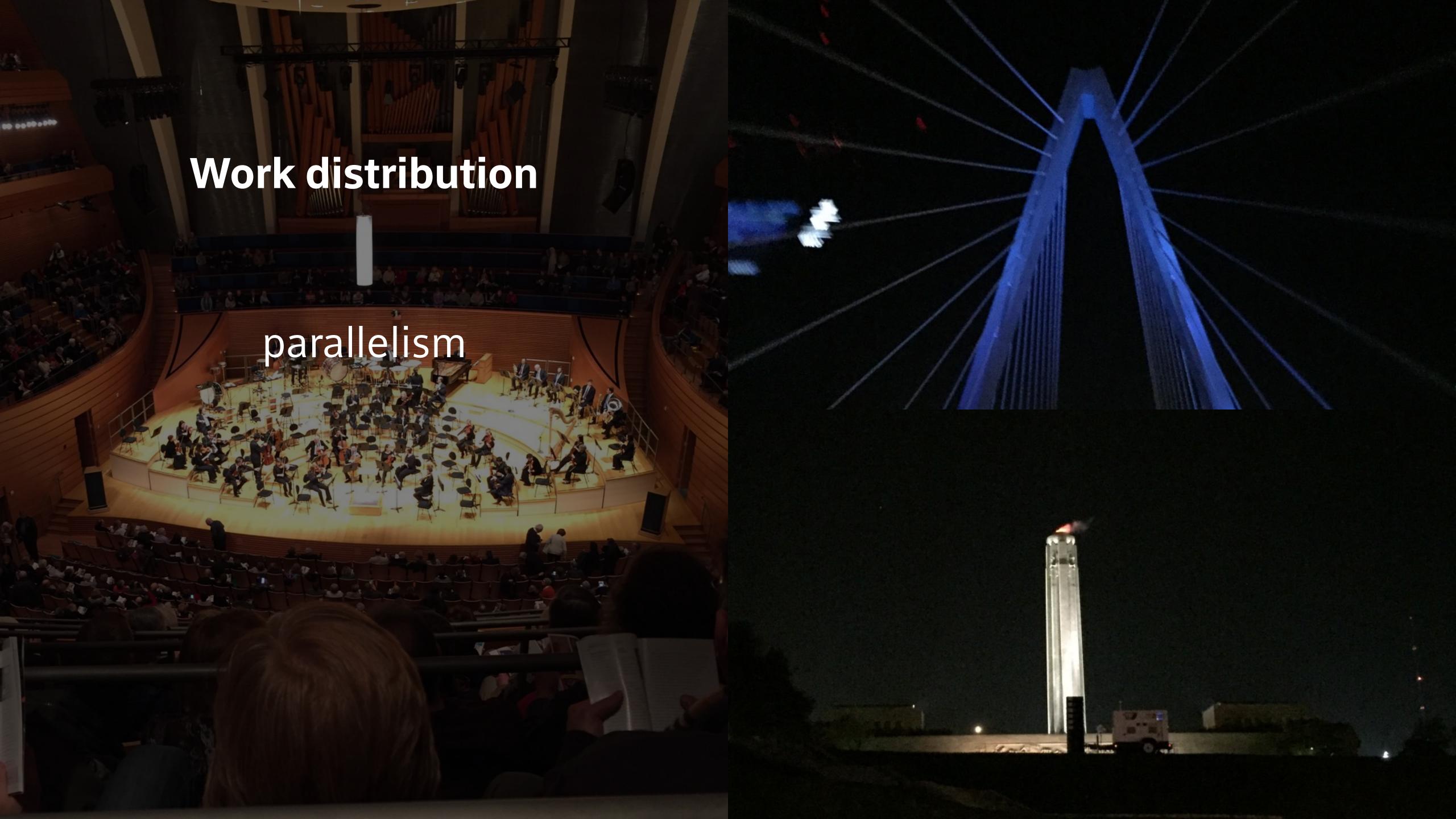


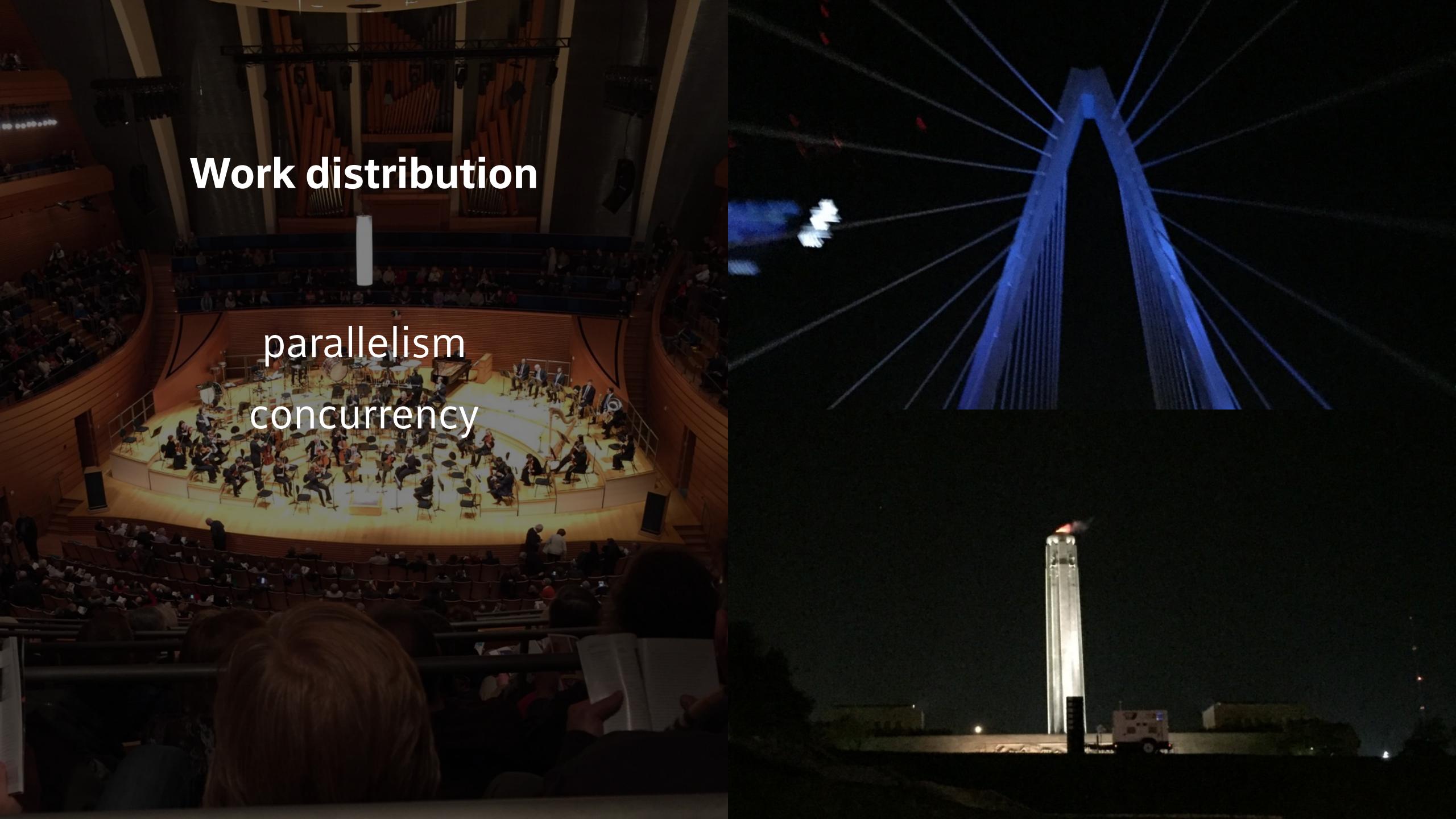
WHY BUILD PEER-TO-PEER SYSTEMS?

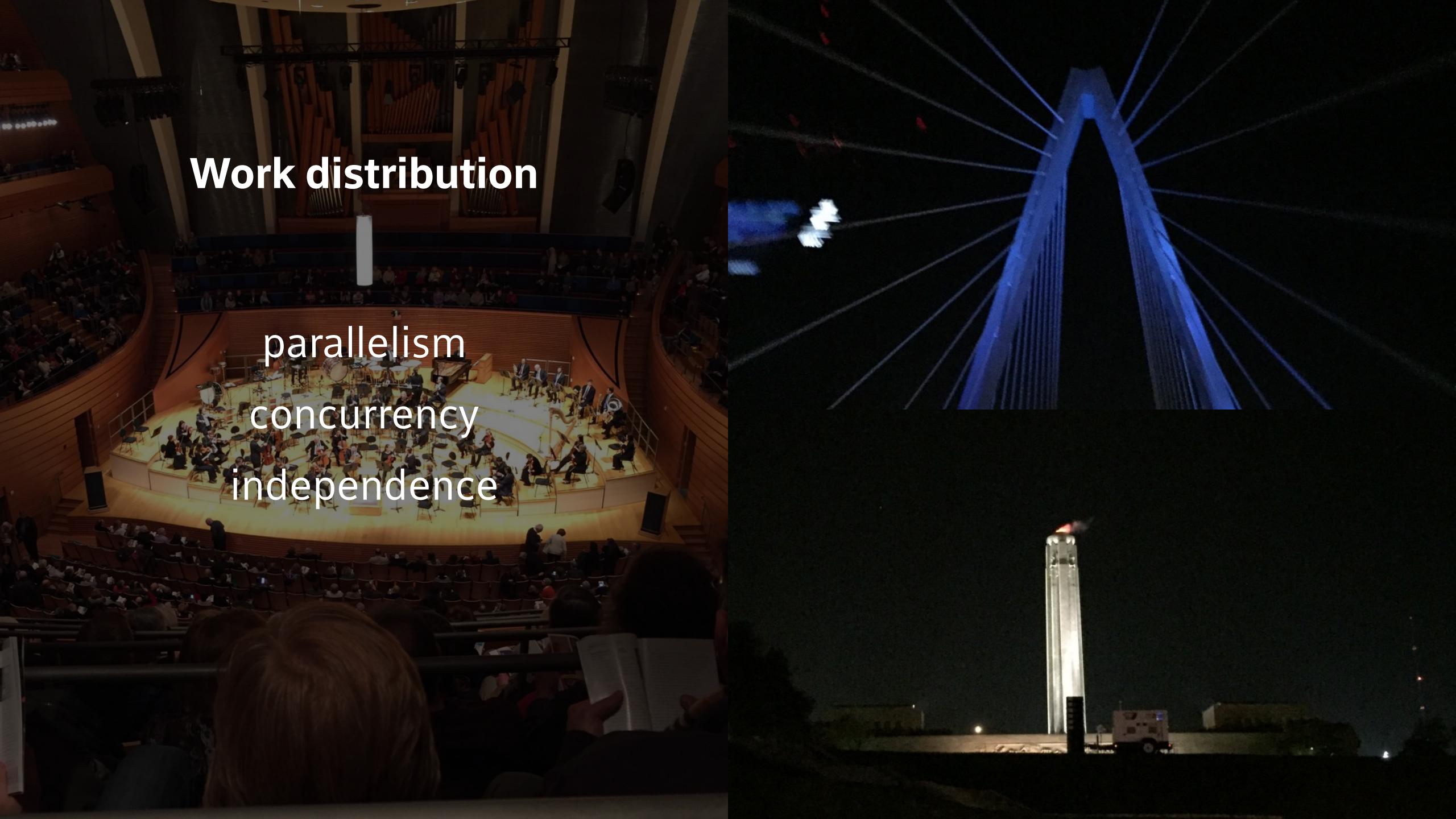


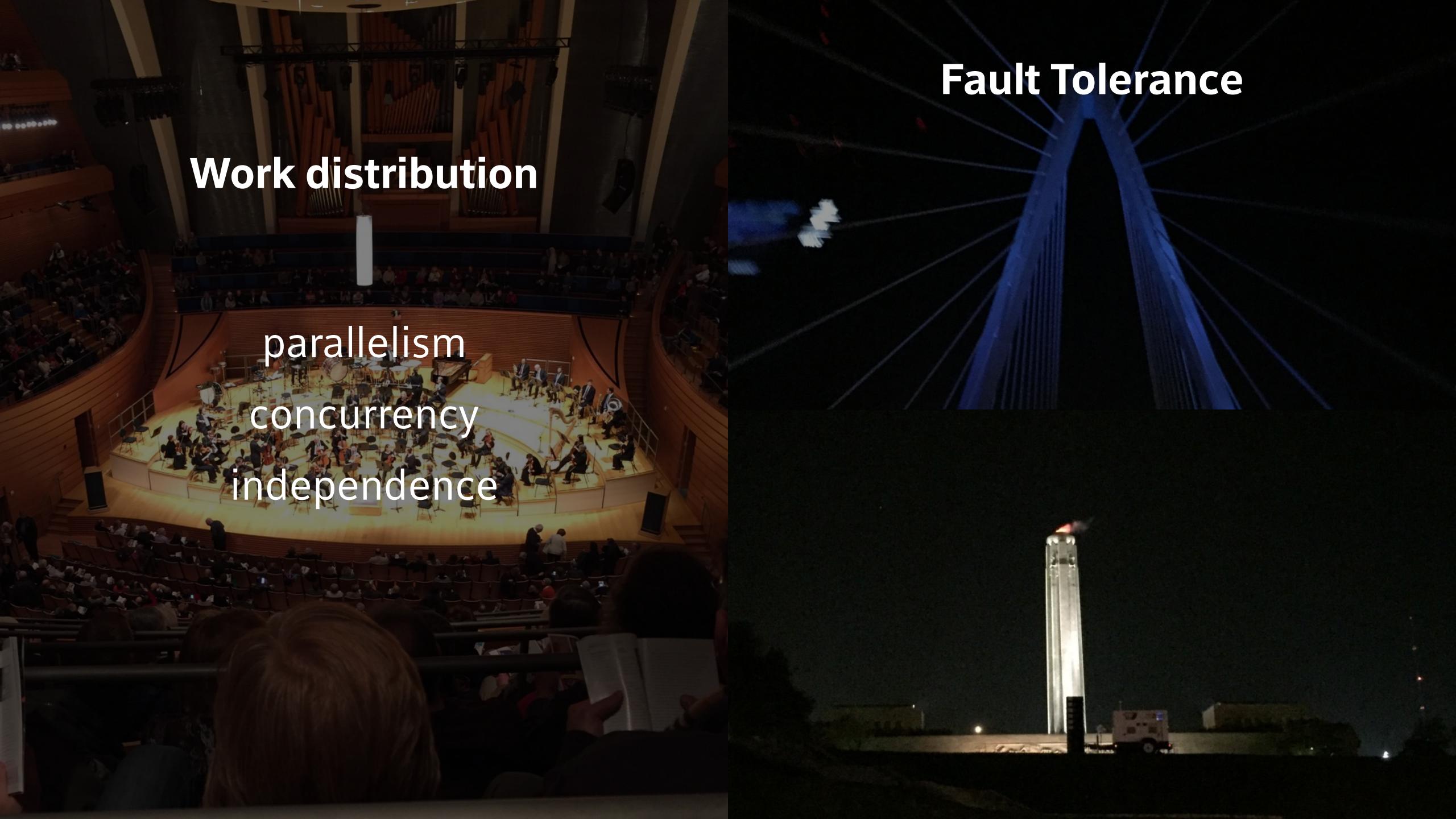


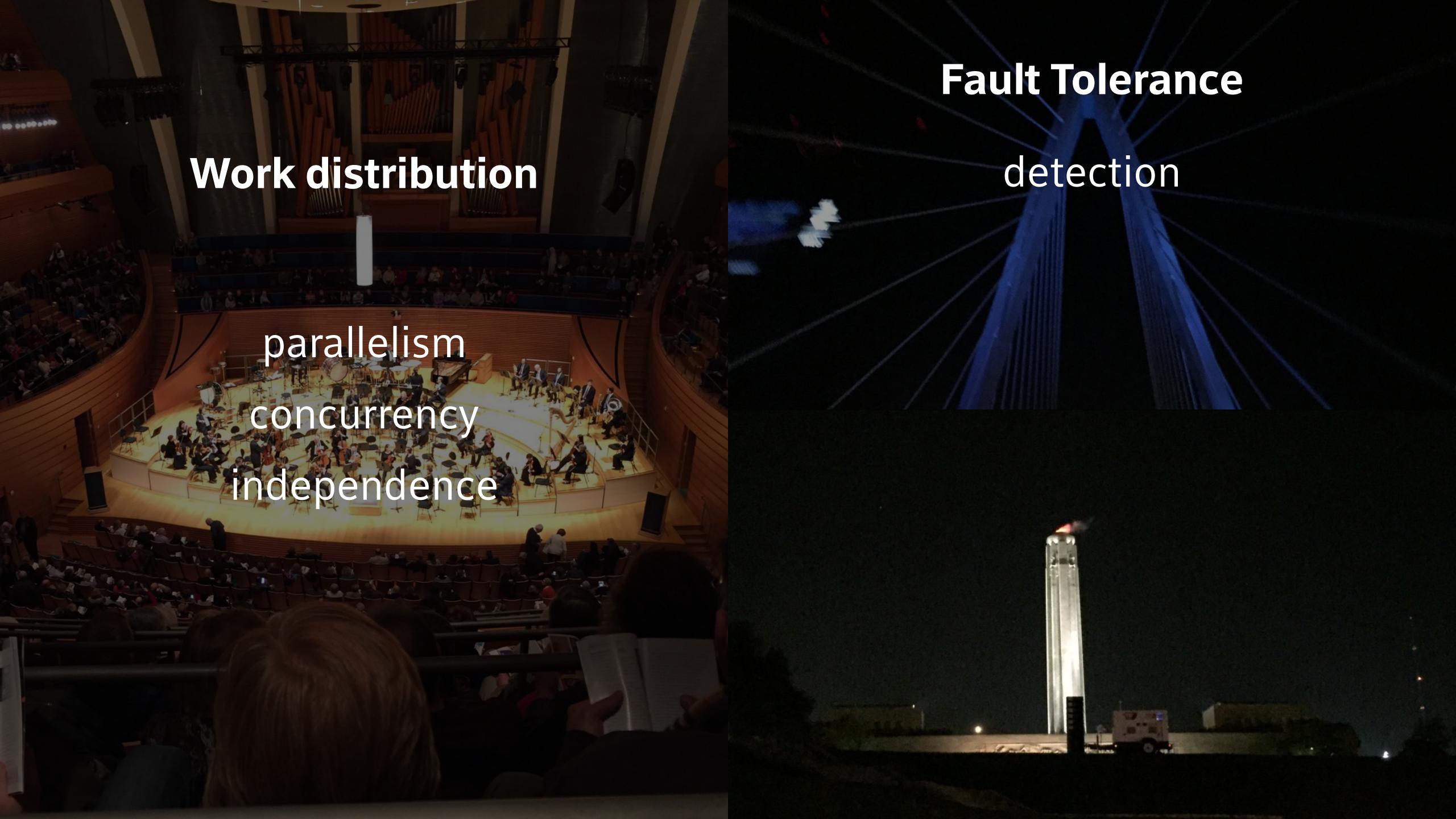


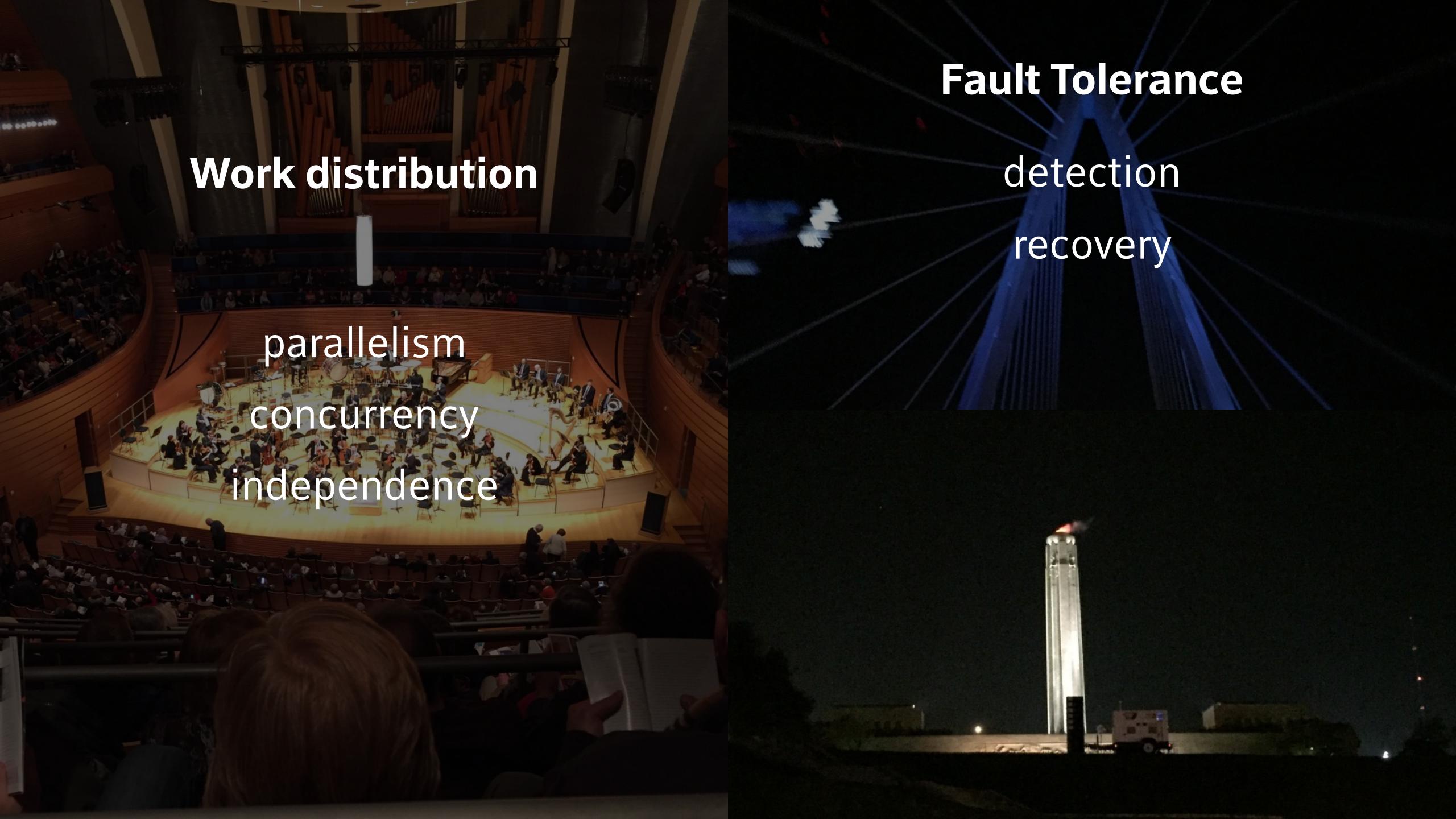


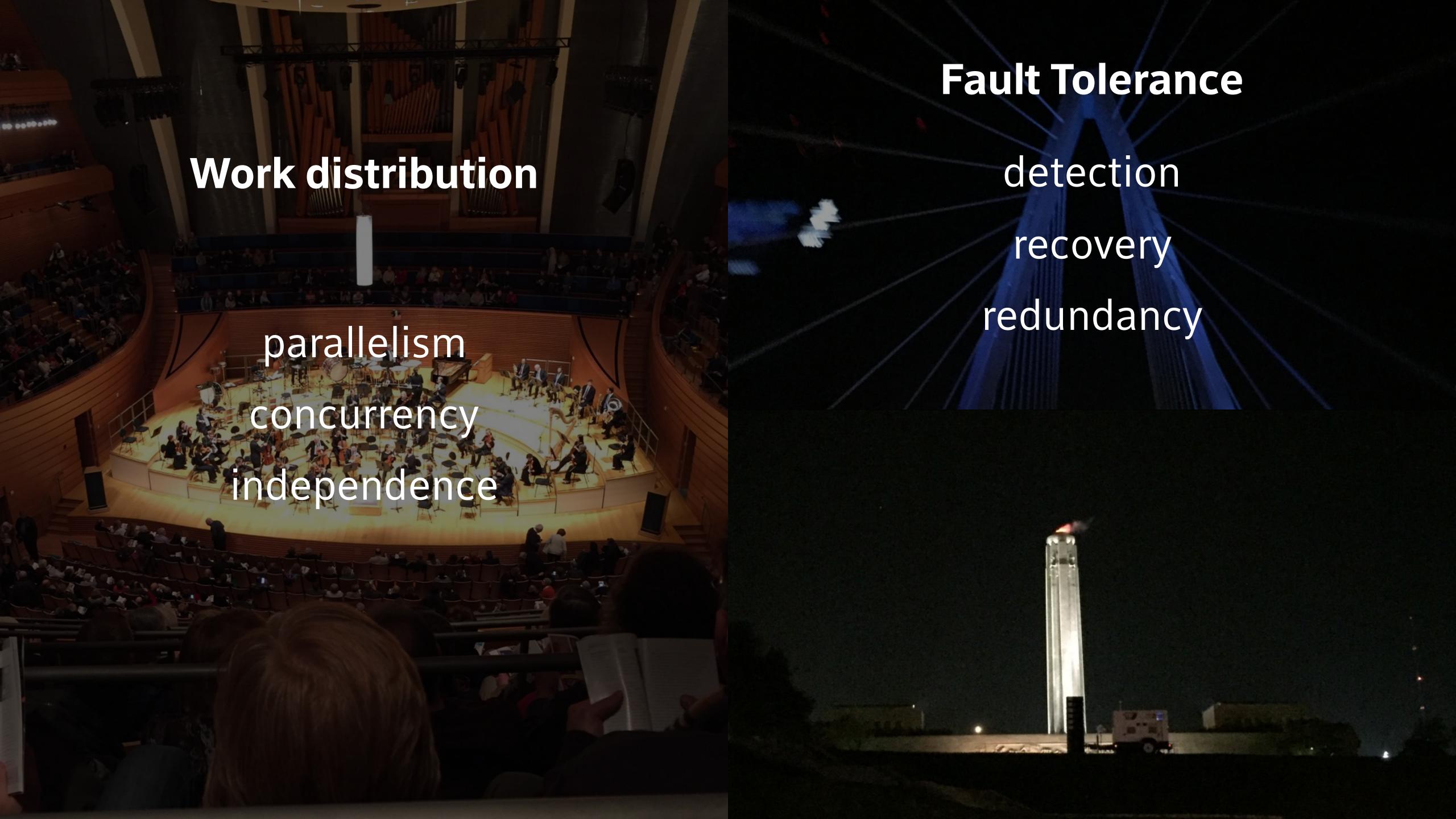


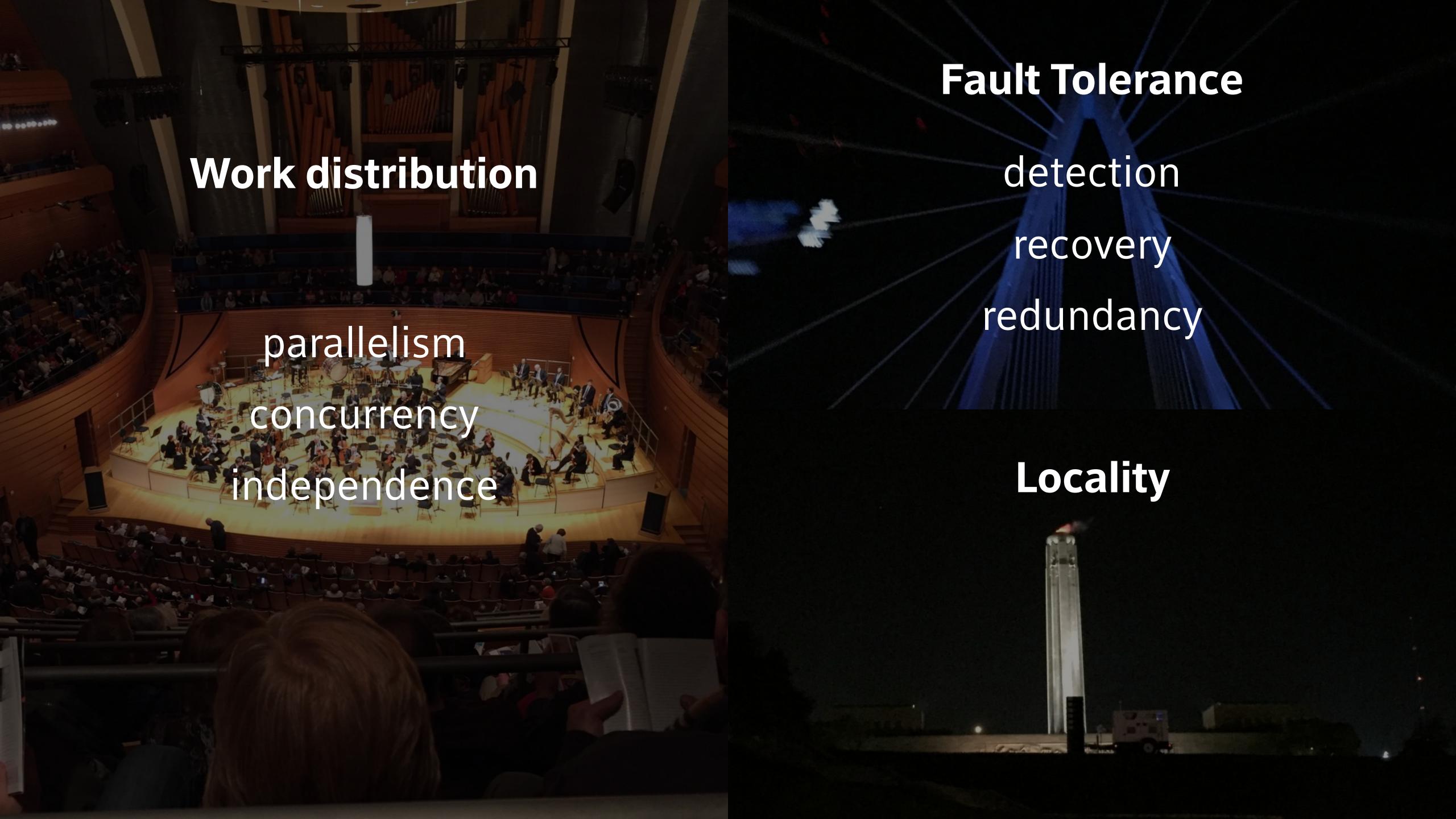


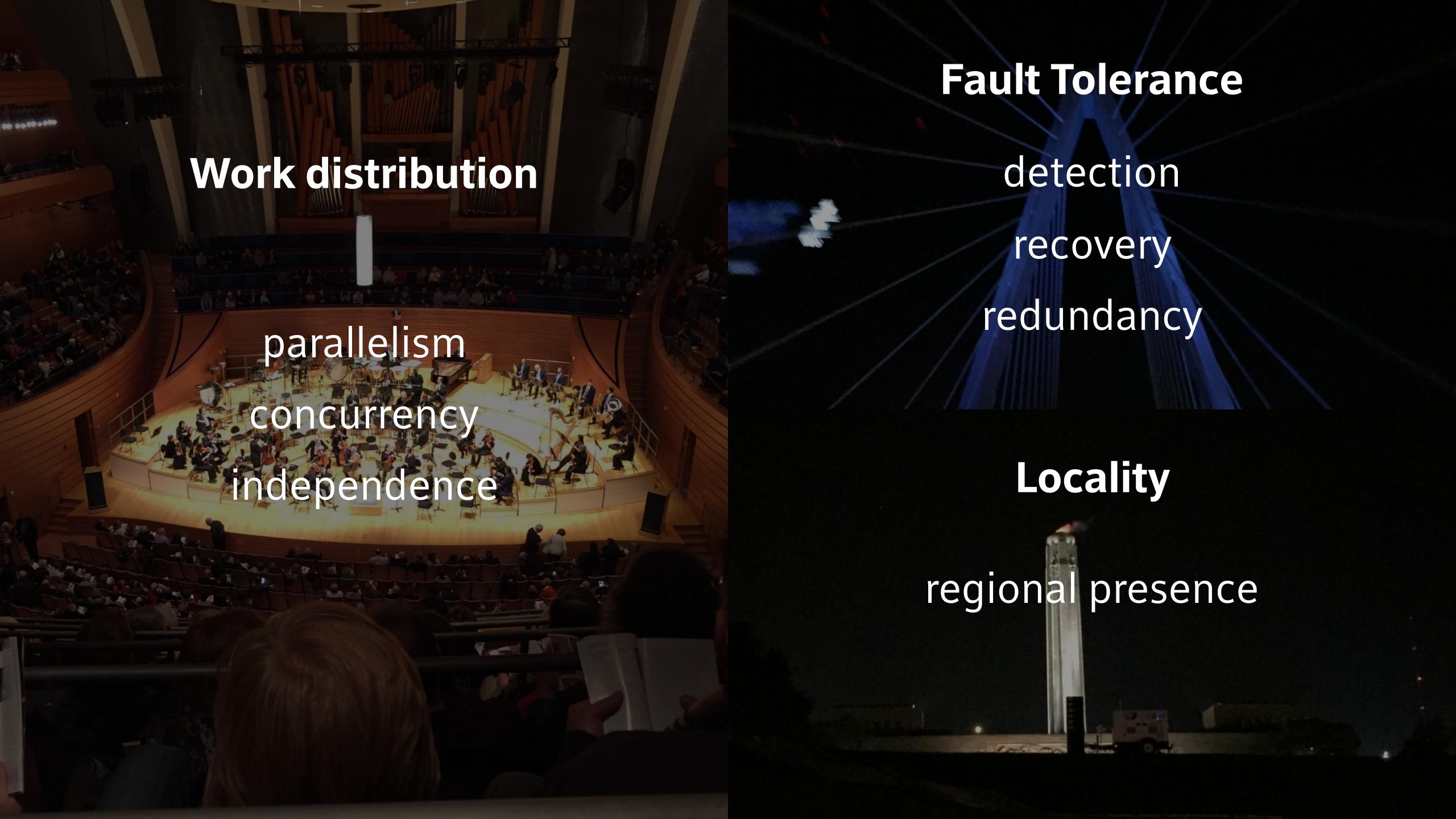


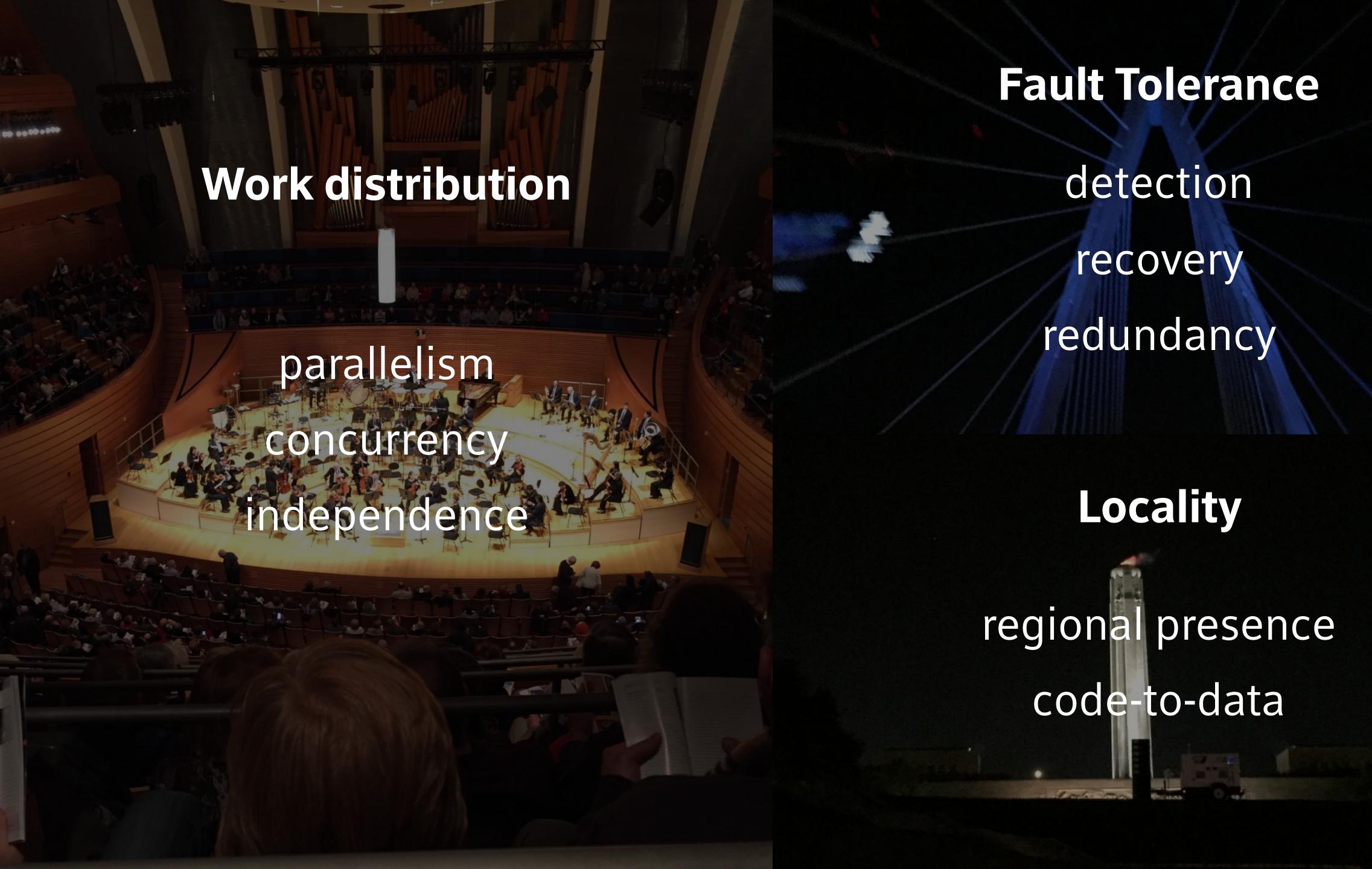












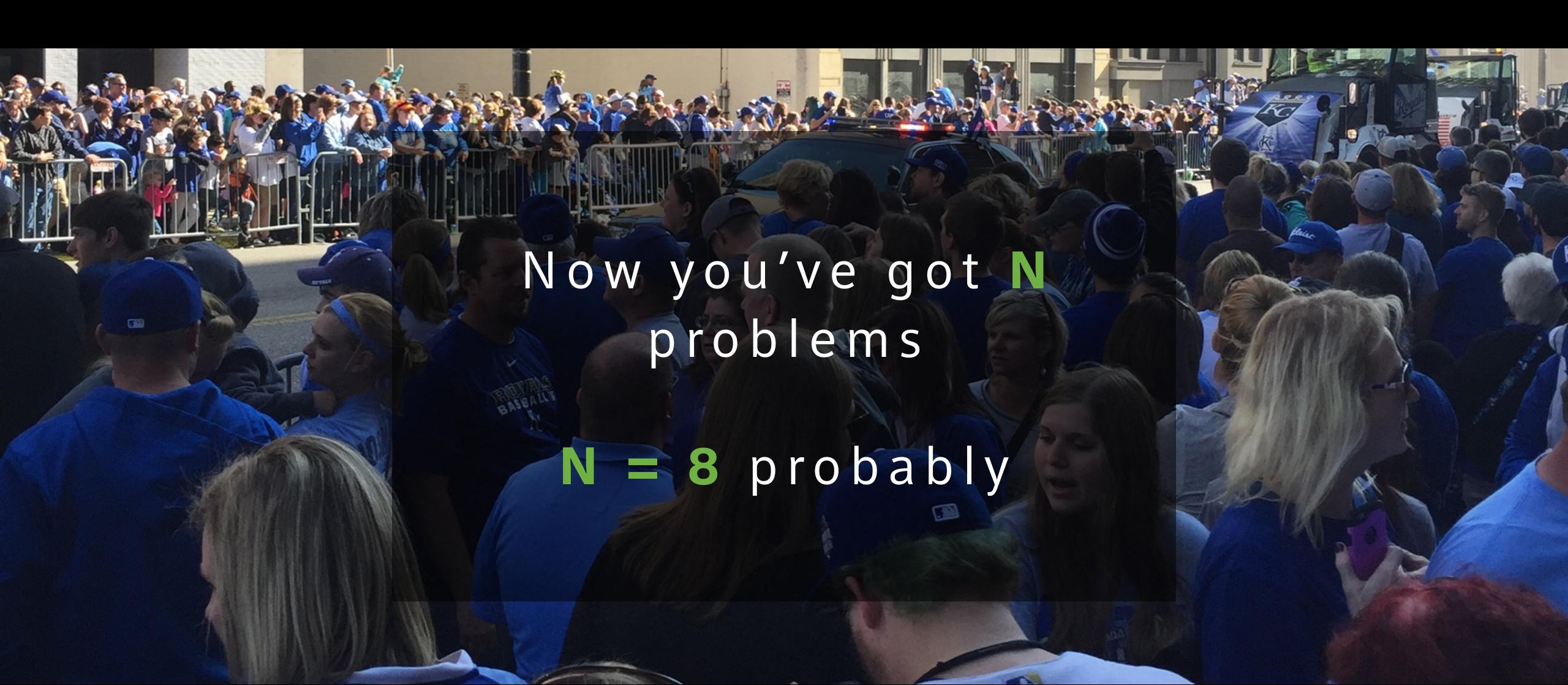
WHY NOT PEER-TO-PEER SYSTEMS?



WHY NOT PEER-TO-PEER SYSTEMS?

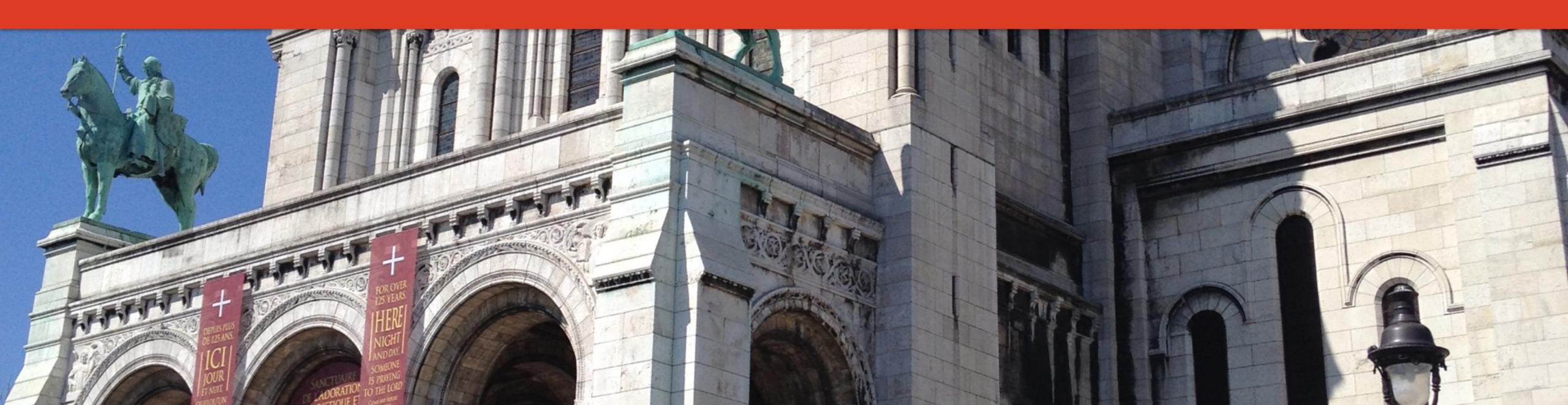


WHY NOT PEER-TO-PEER SYSTEMS?





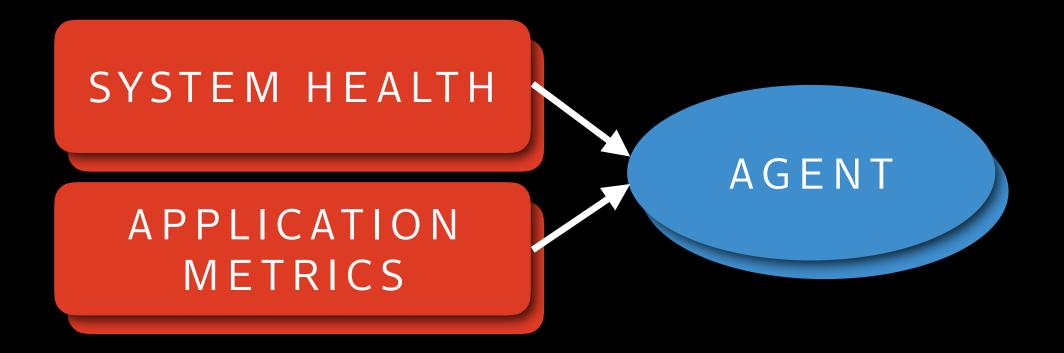
WHAT ARE WE BUILDING?



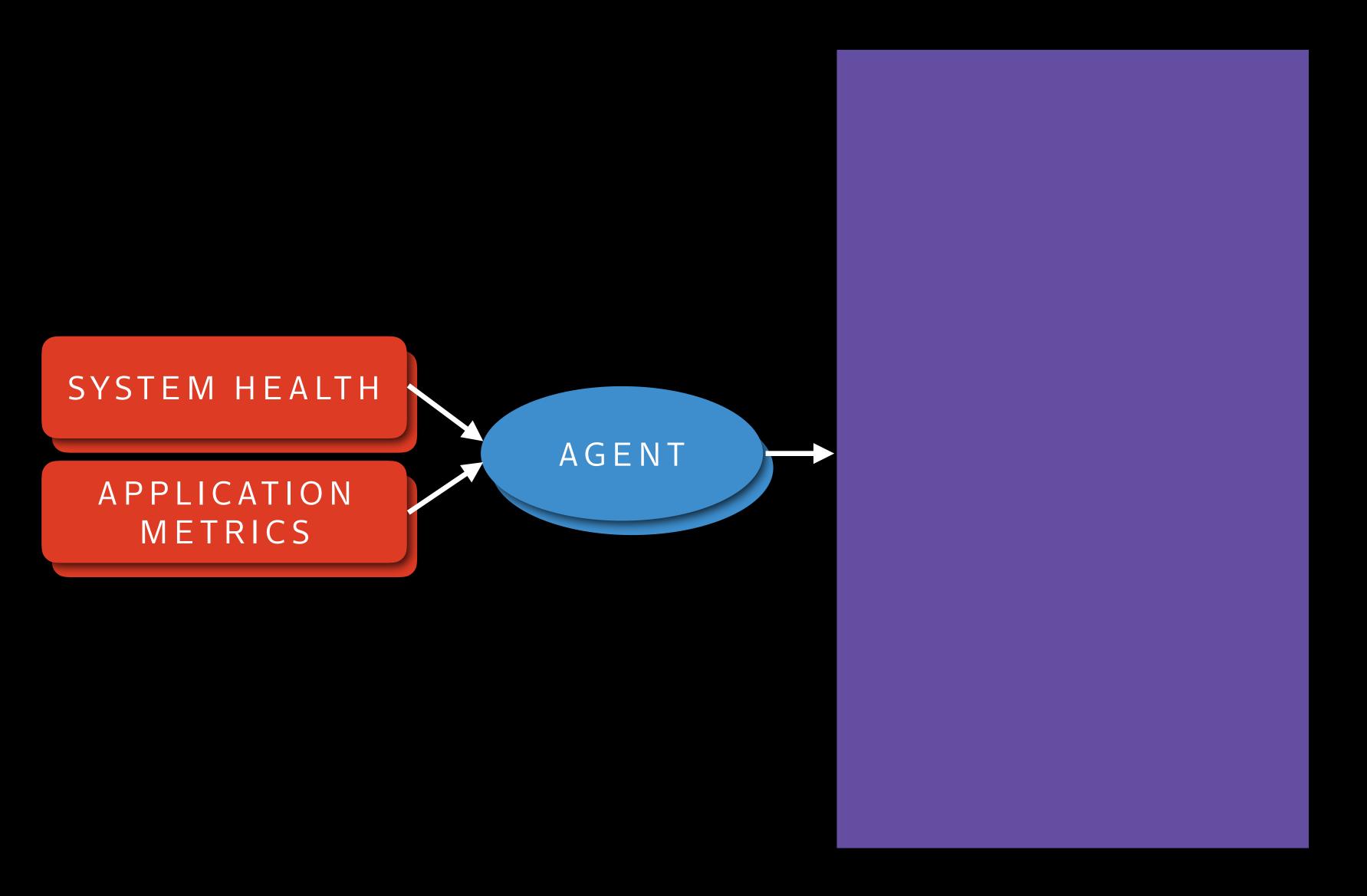
SYSTEM HEALTH

APPLICATION METRICS

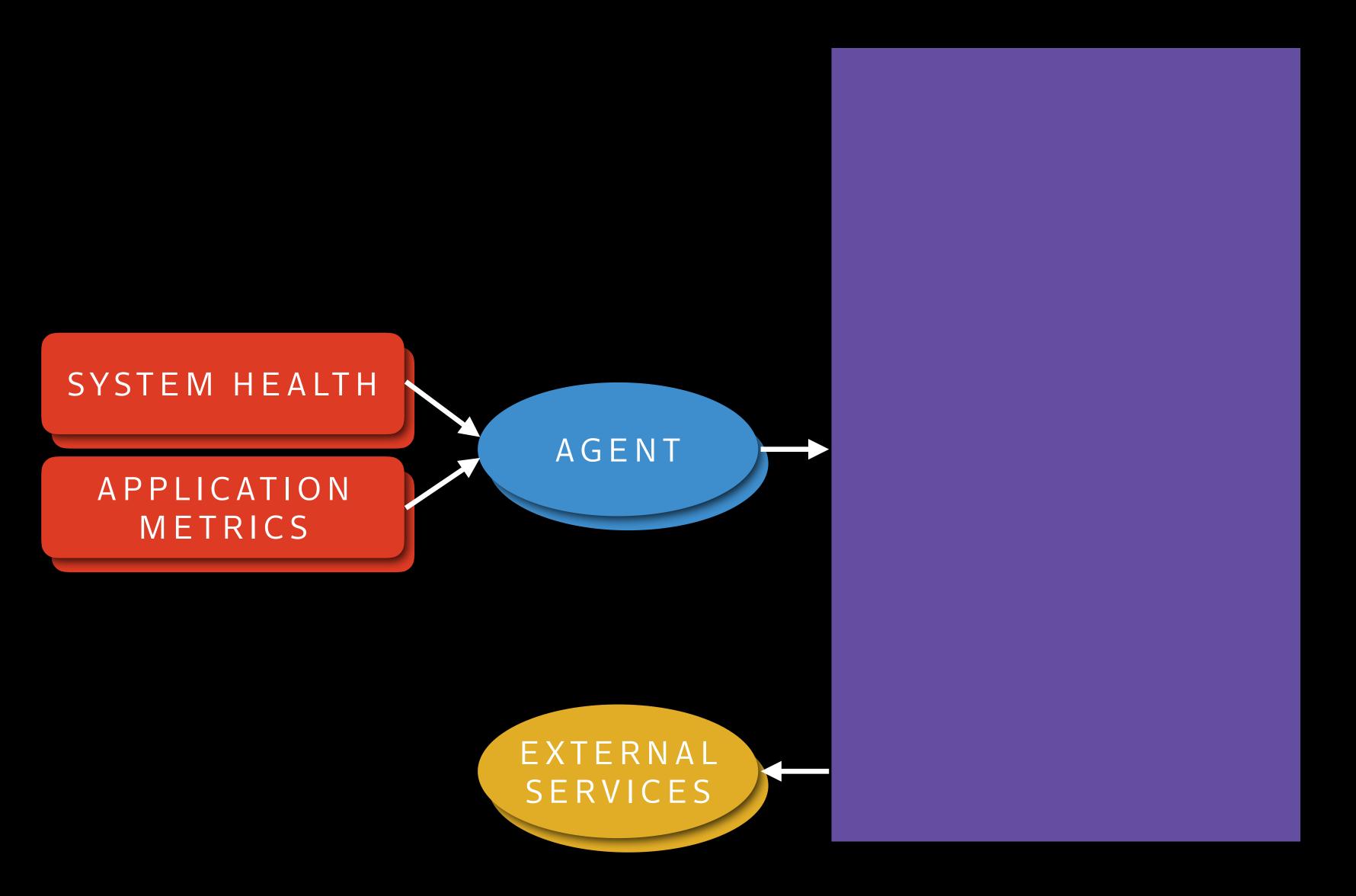




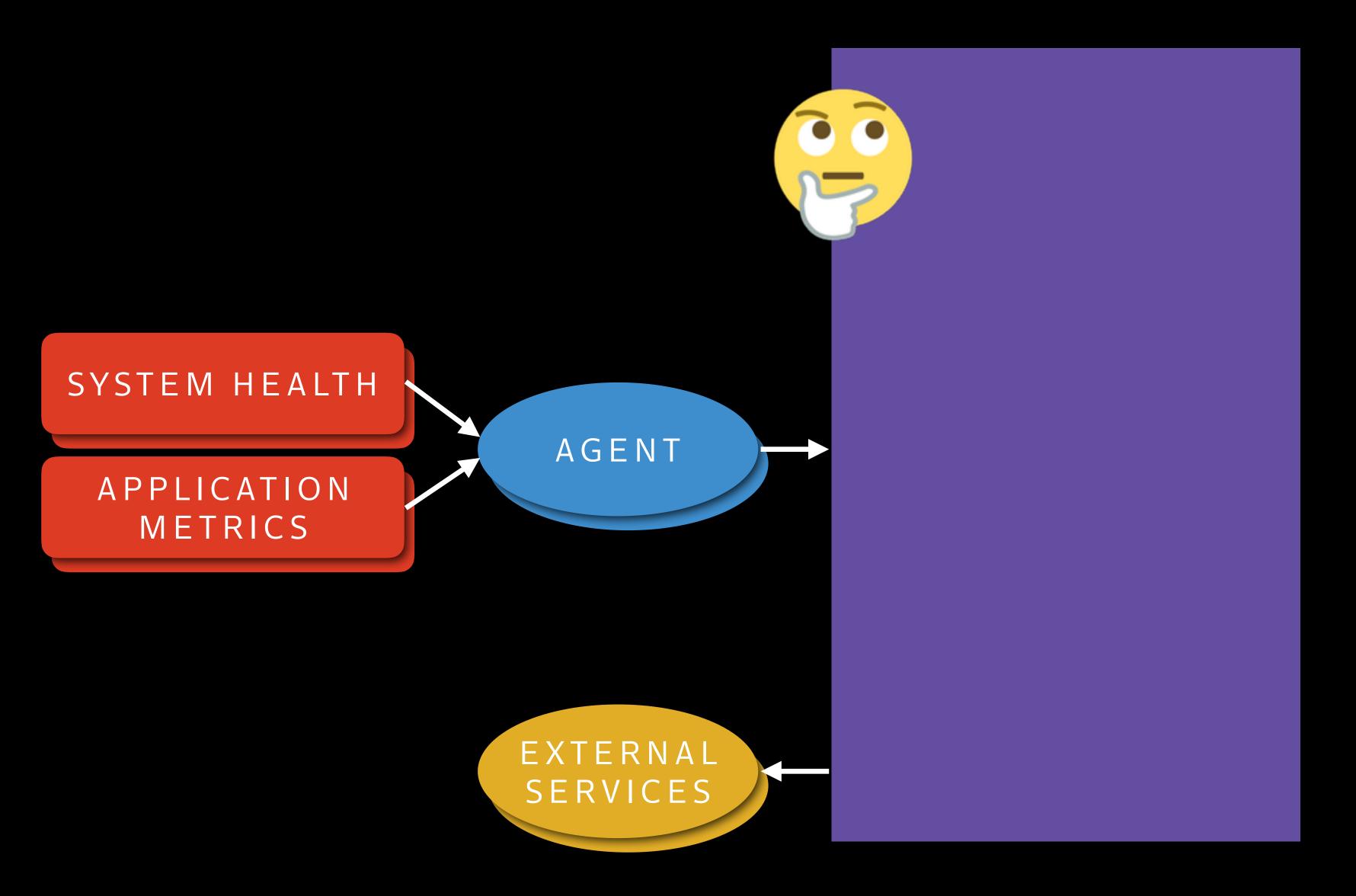




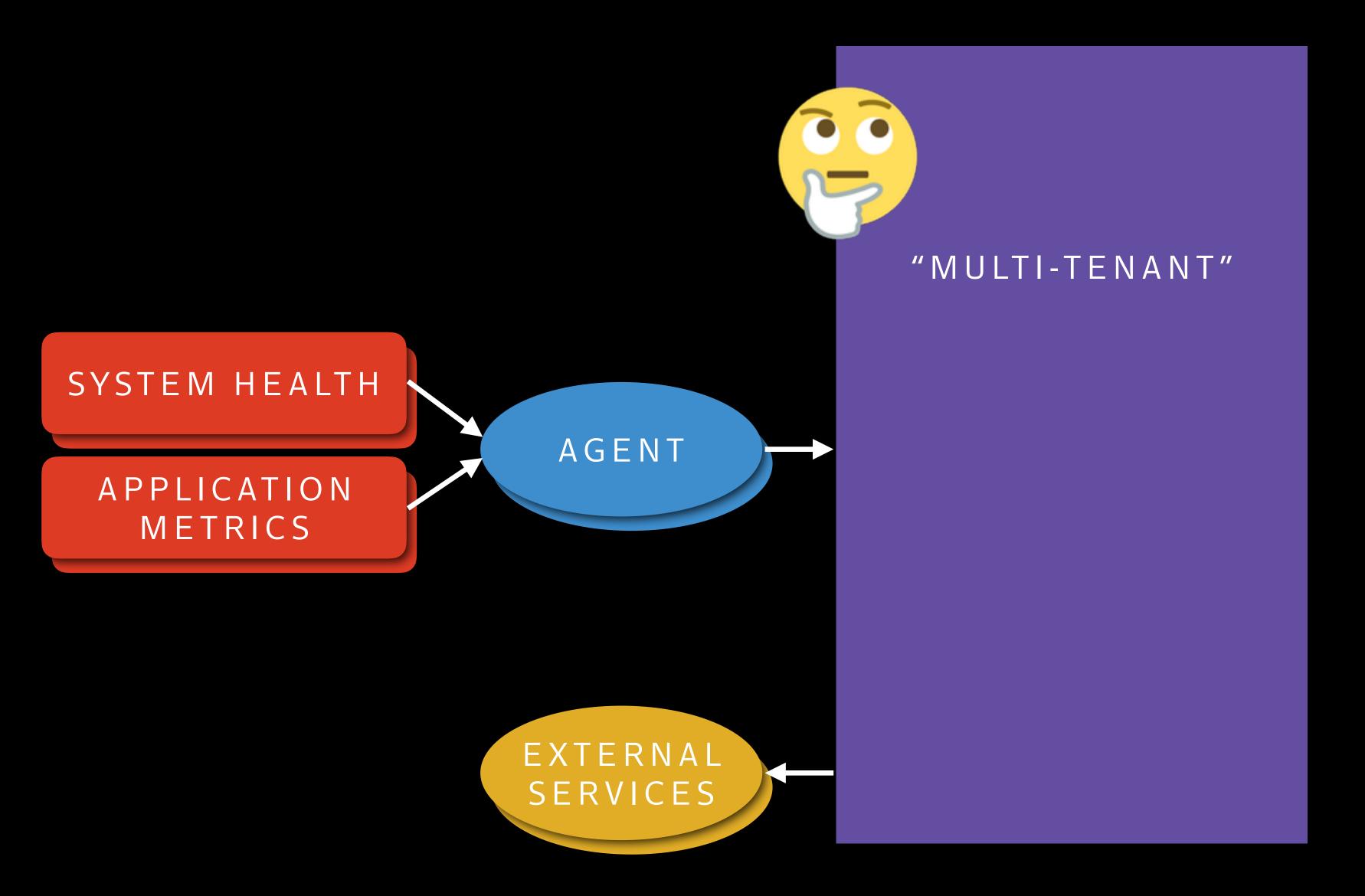




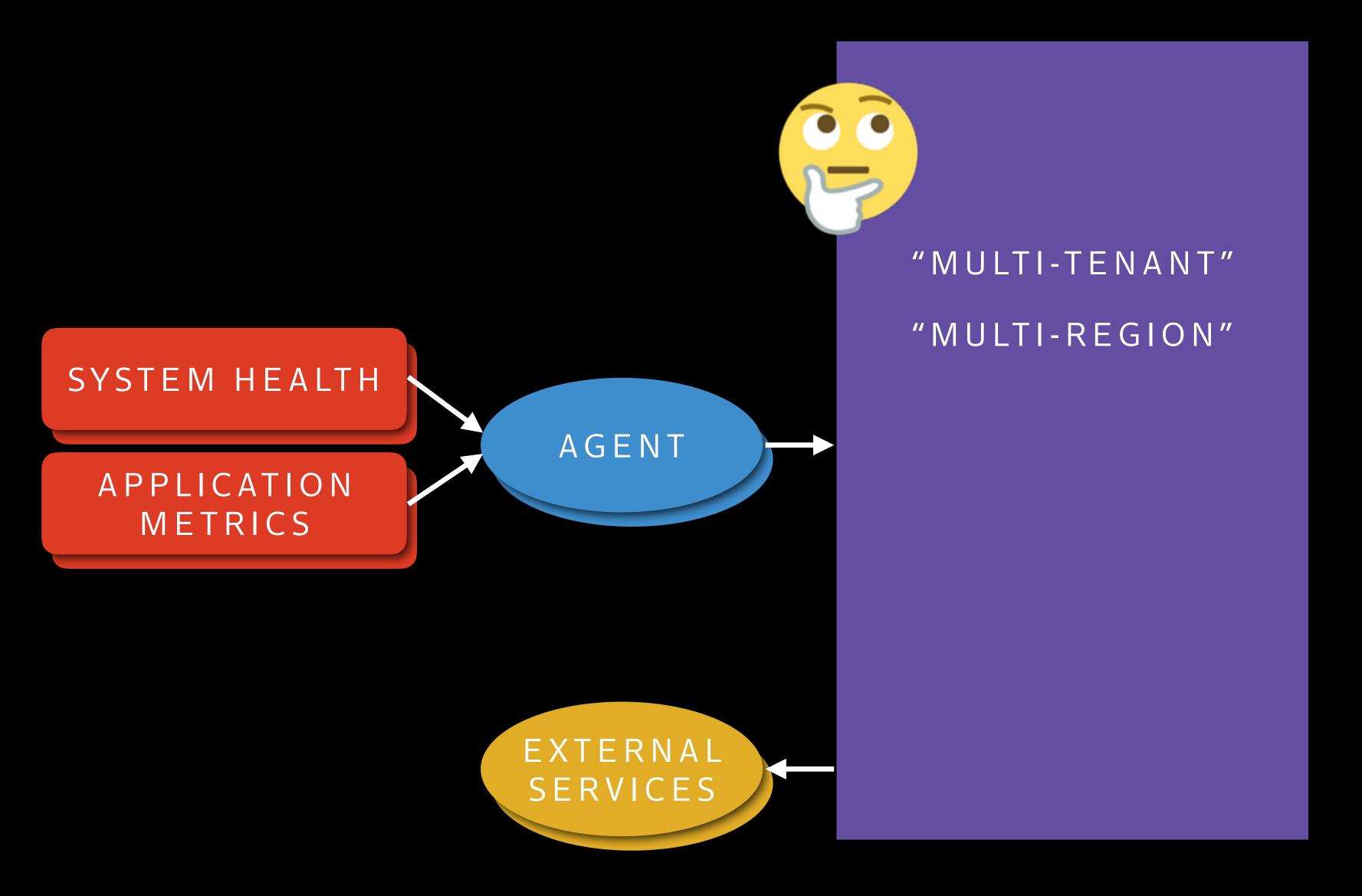




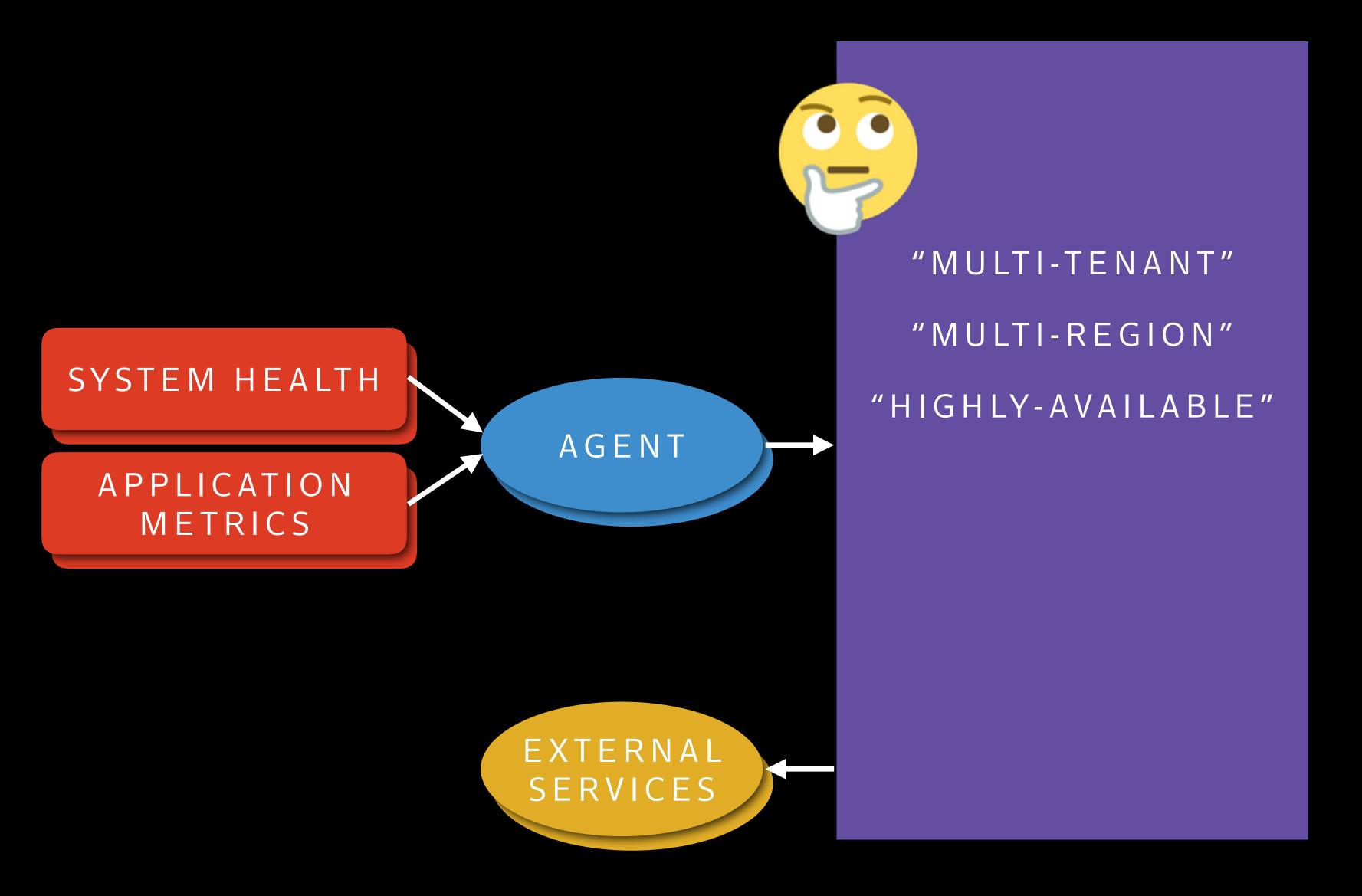




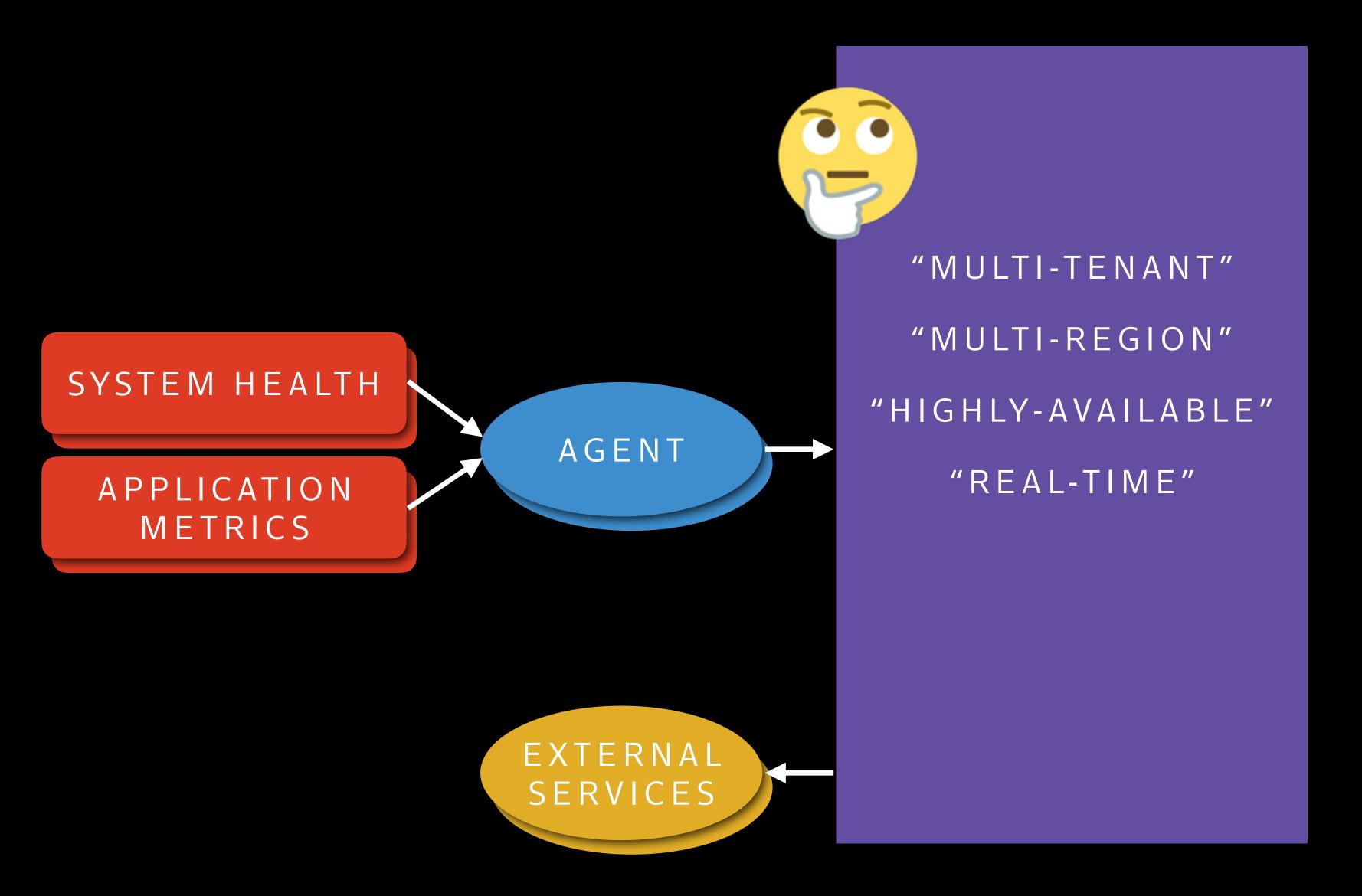




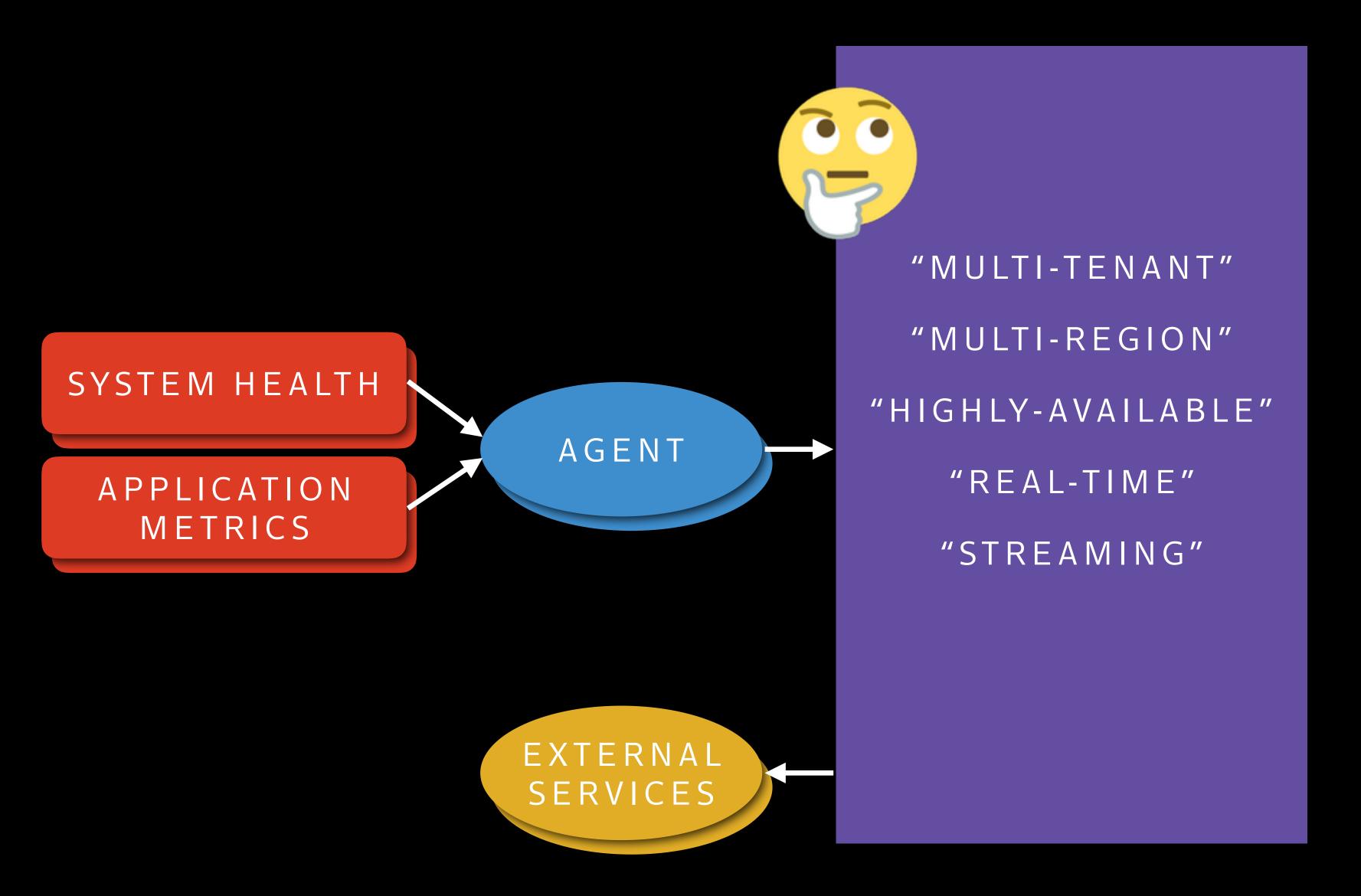




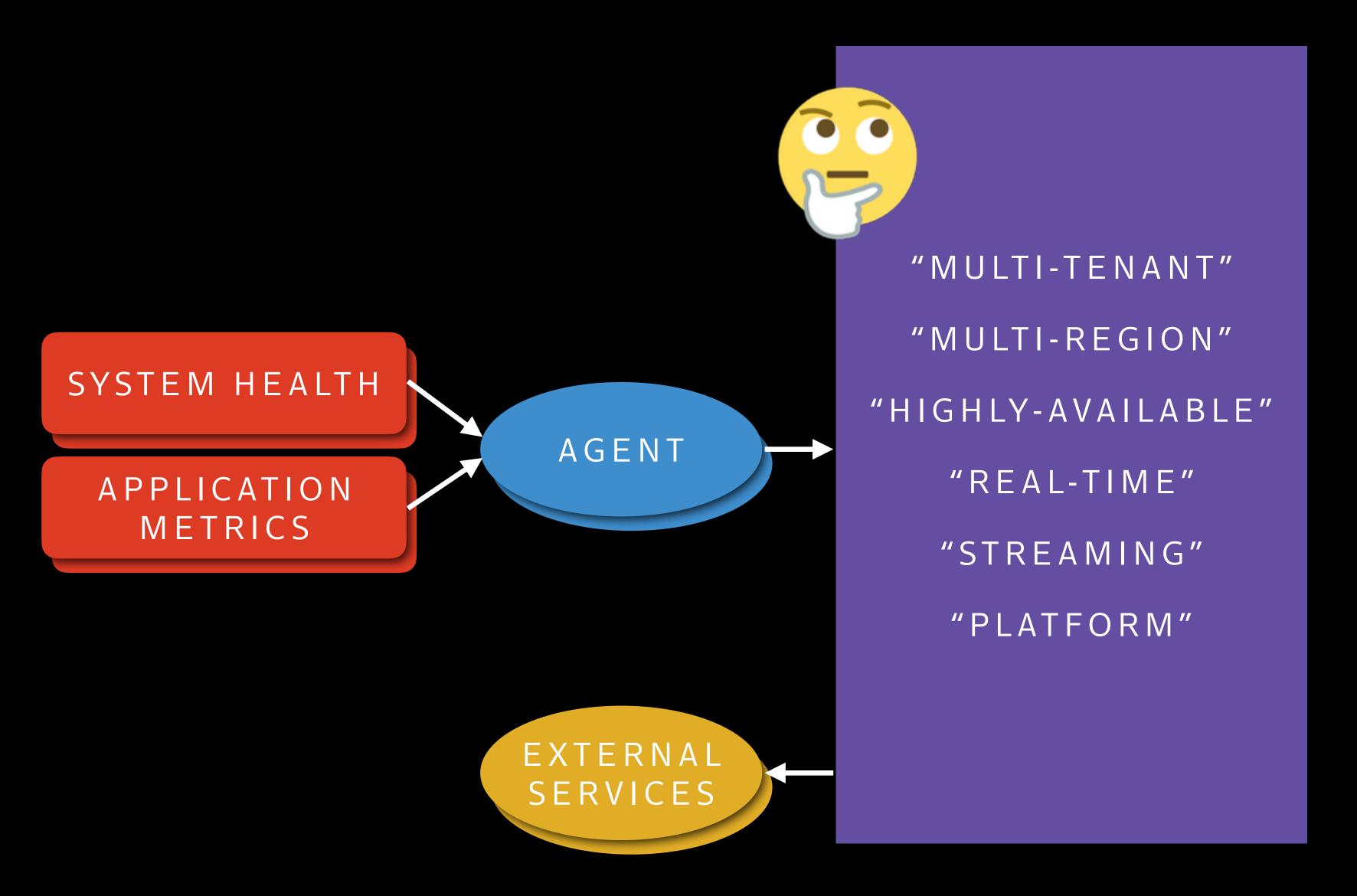




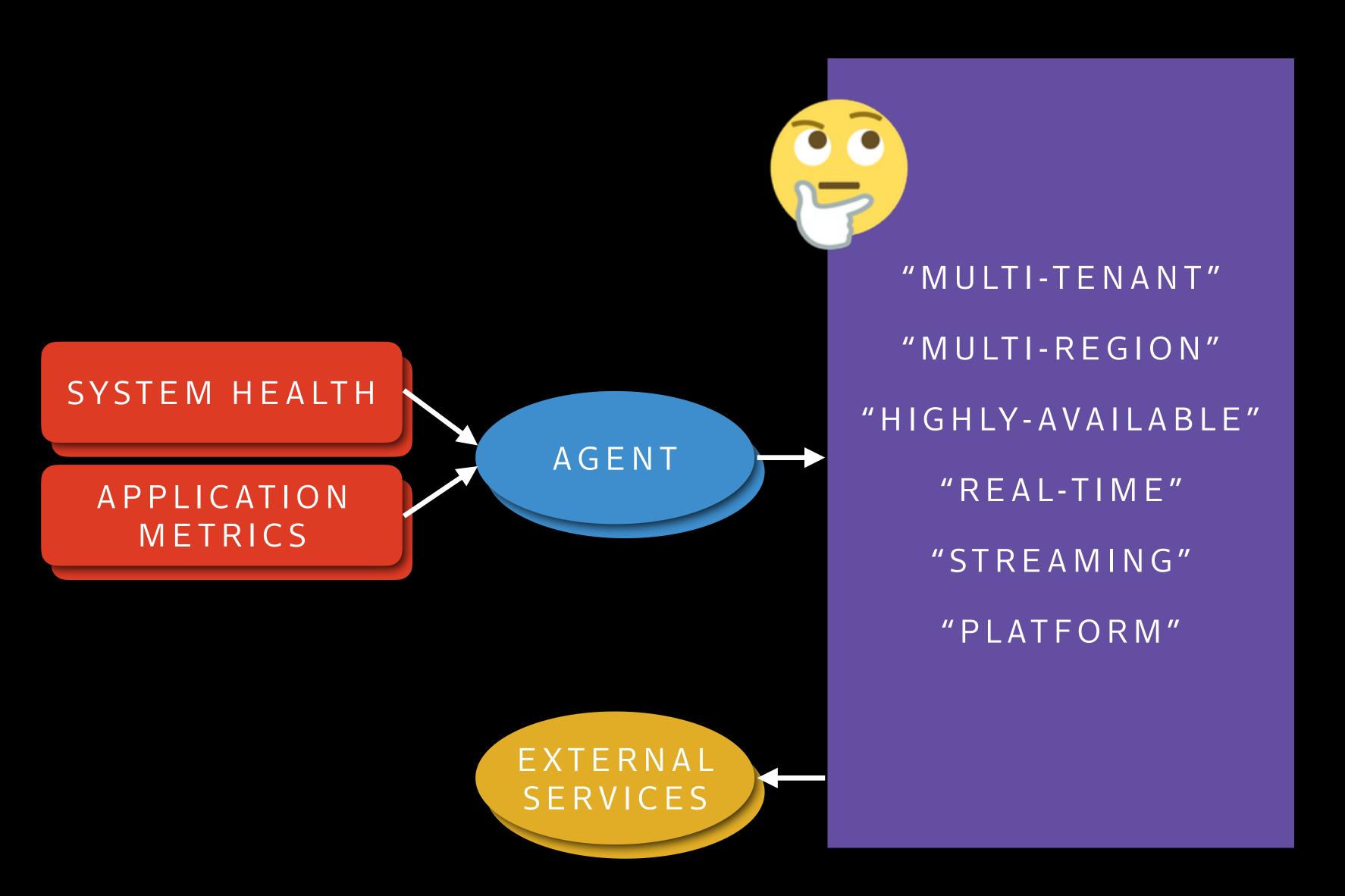






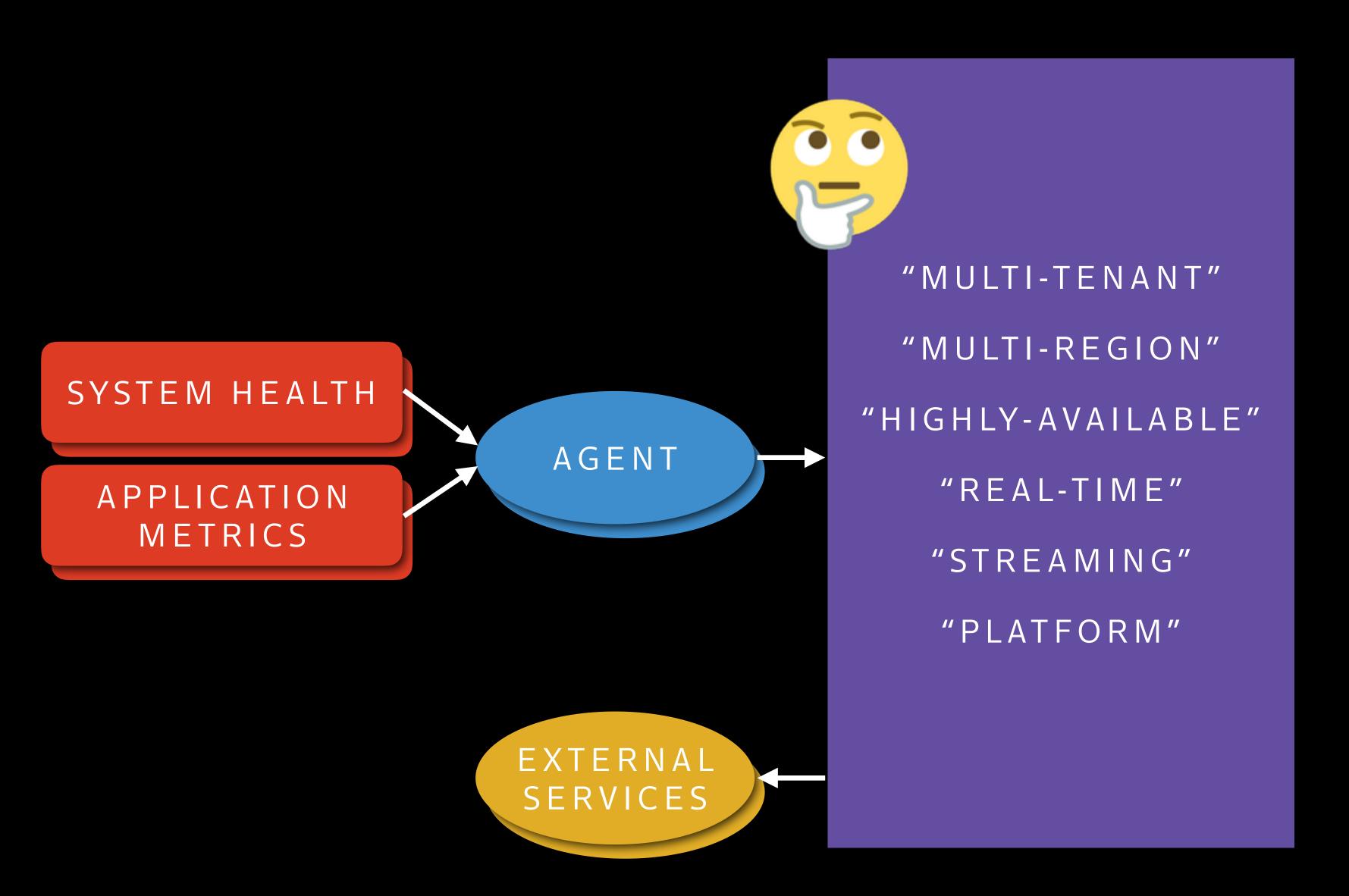






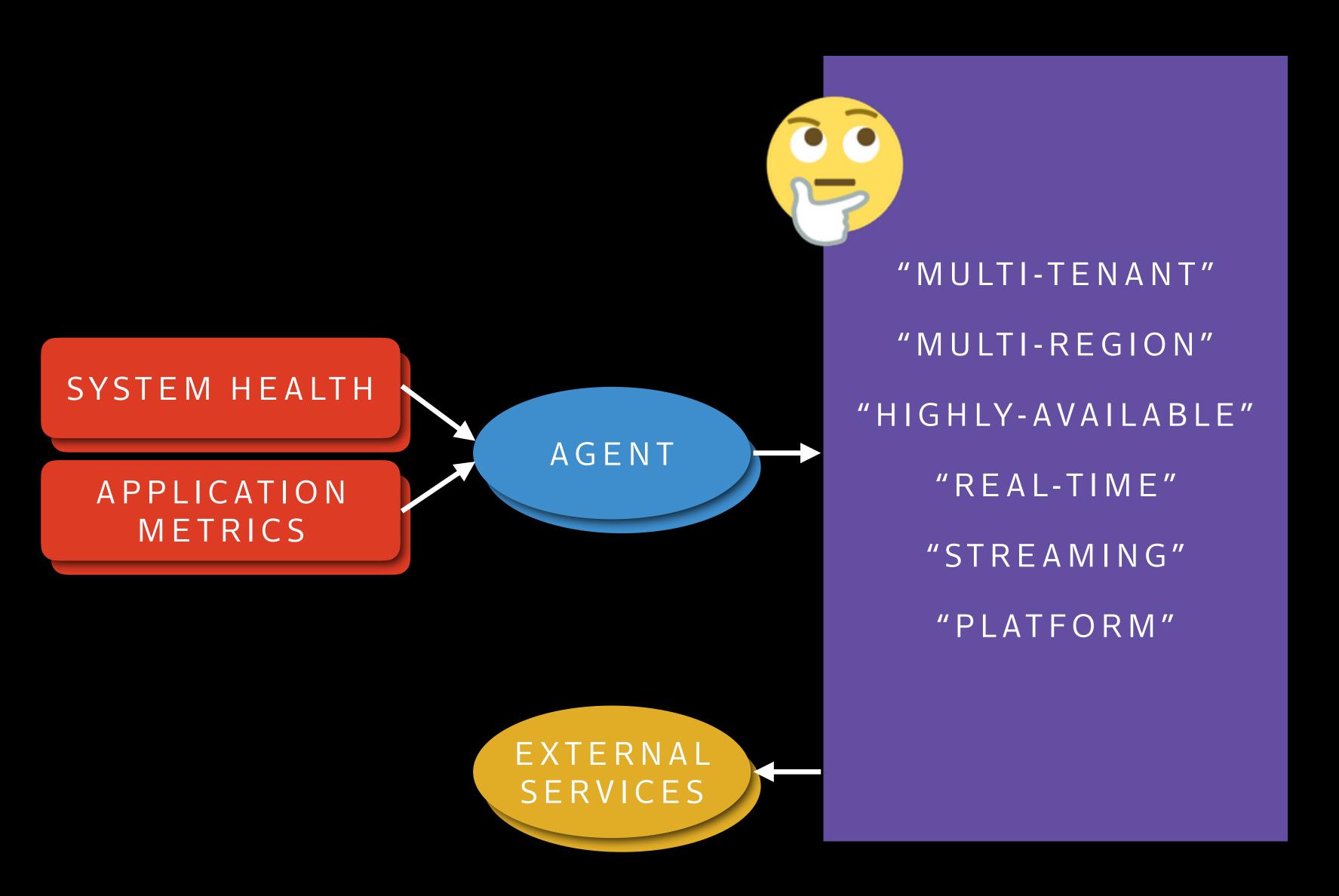
Work distribution





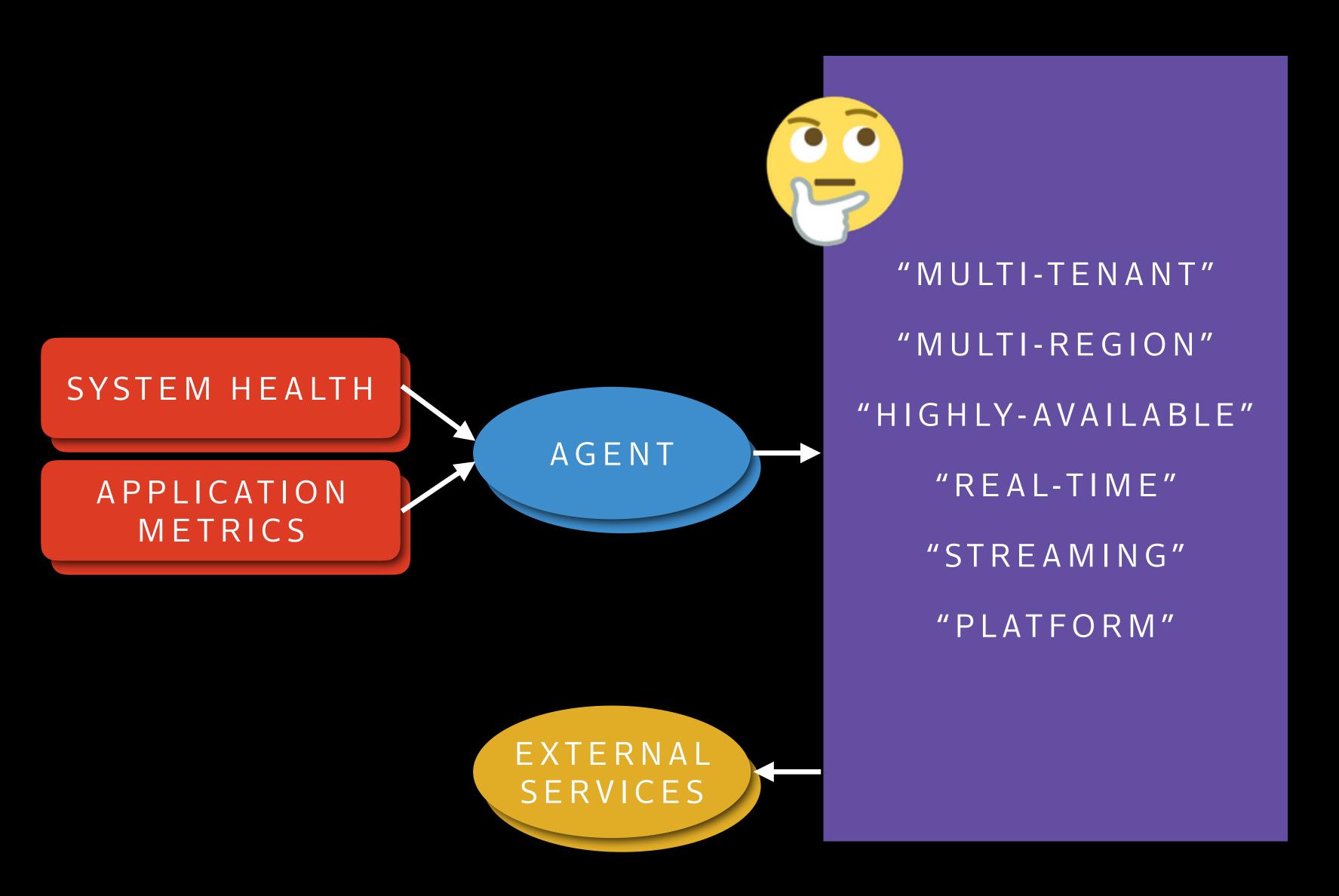
- Work distribution
- → Fault-tolerance





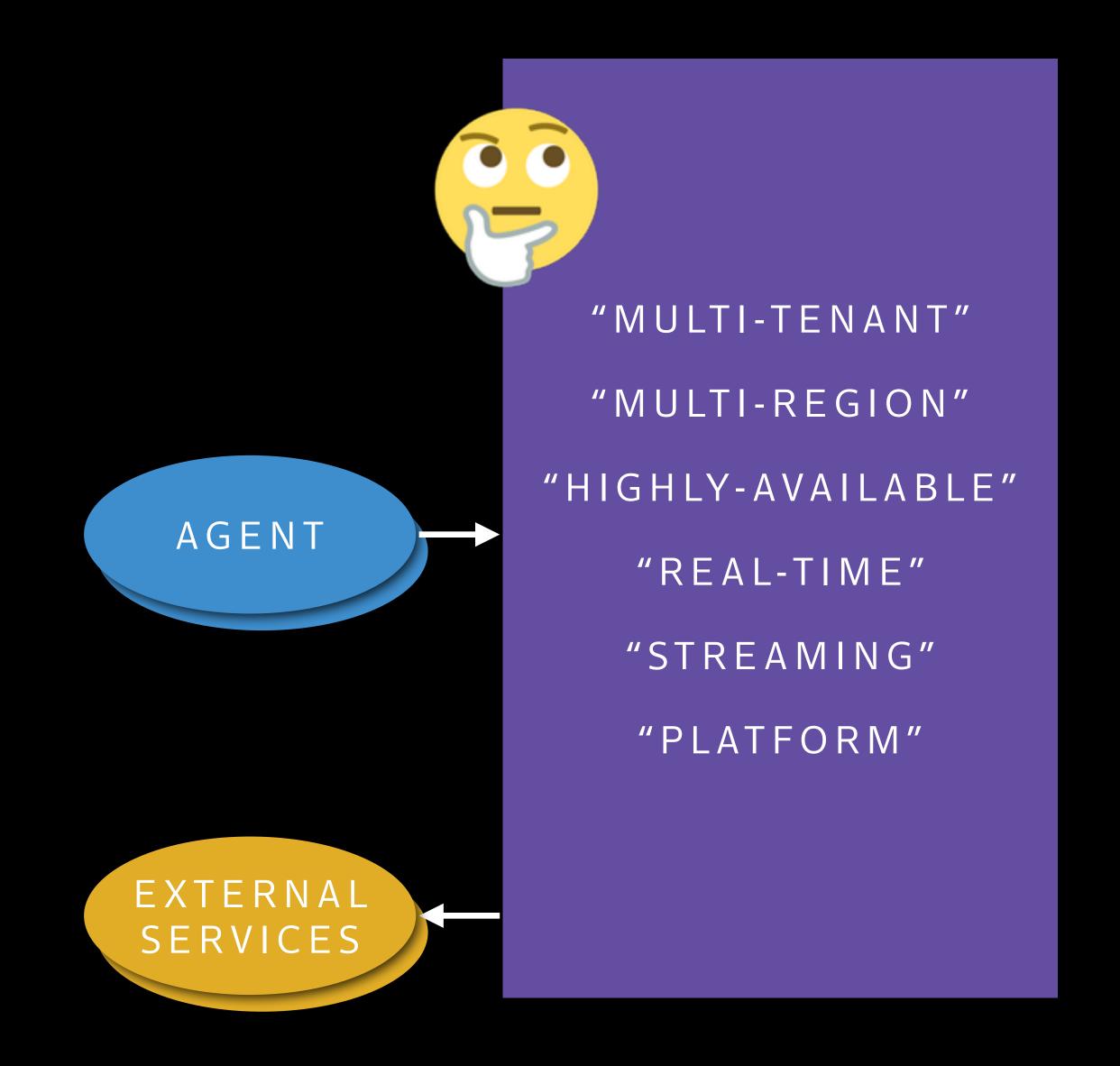
- Work distribution
- → Fault-tolerance
- → Locality



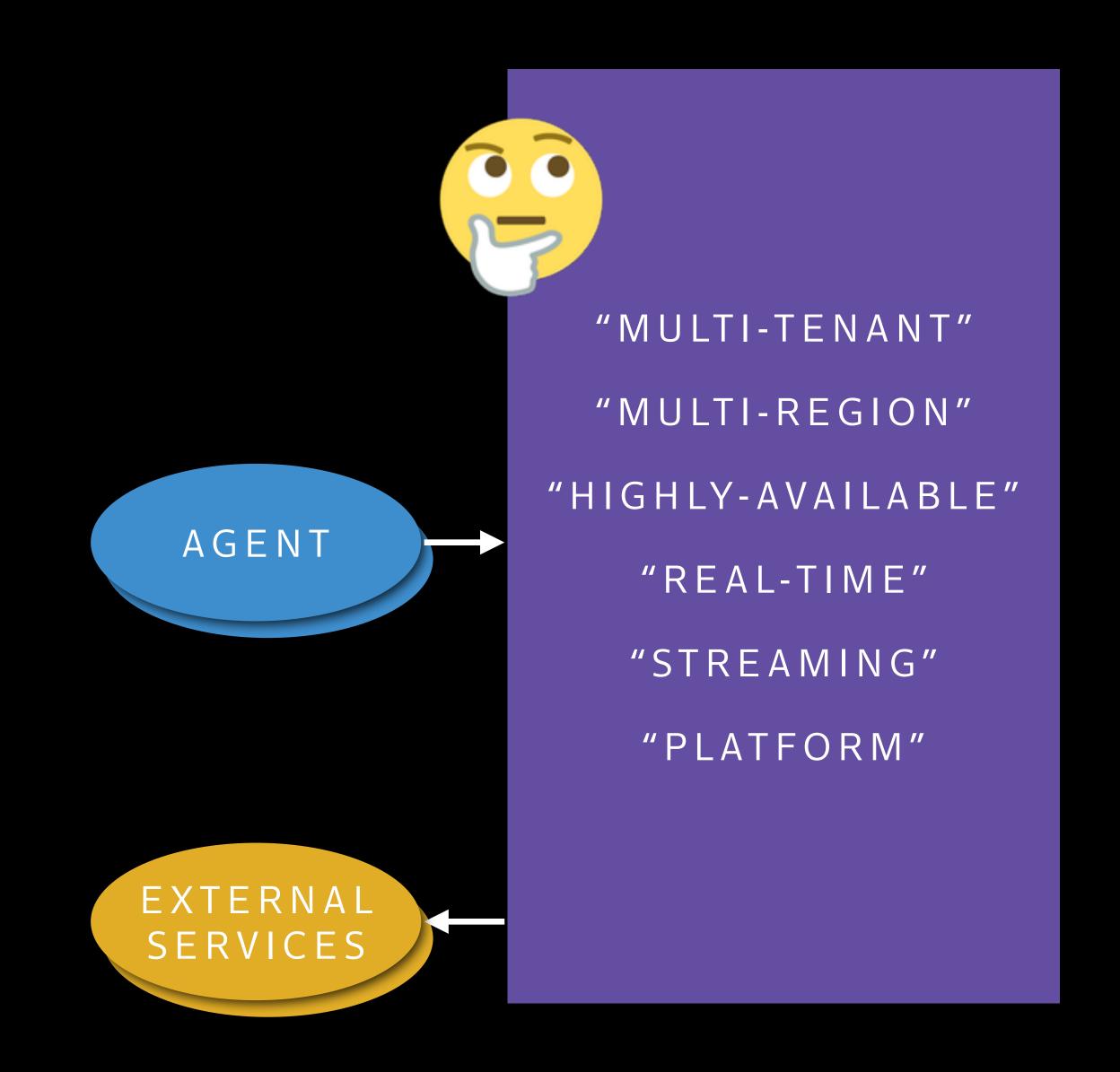


- Work distribution
- Fault-tolerance
- Locality
- ✓ Peer to Peer!



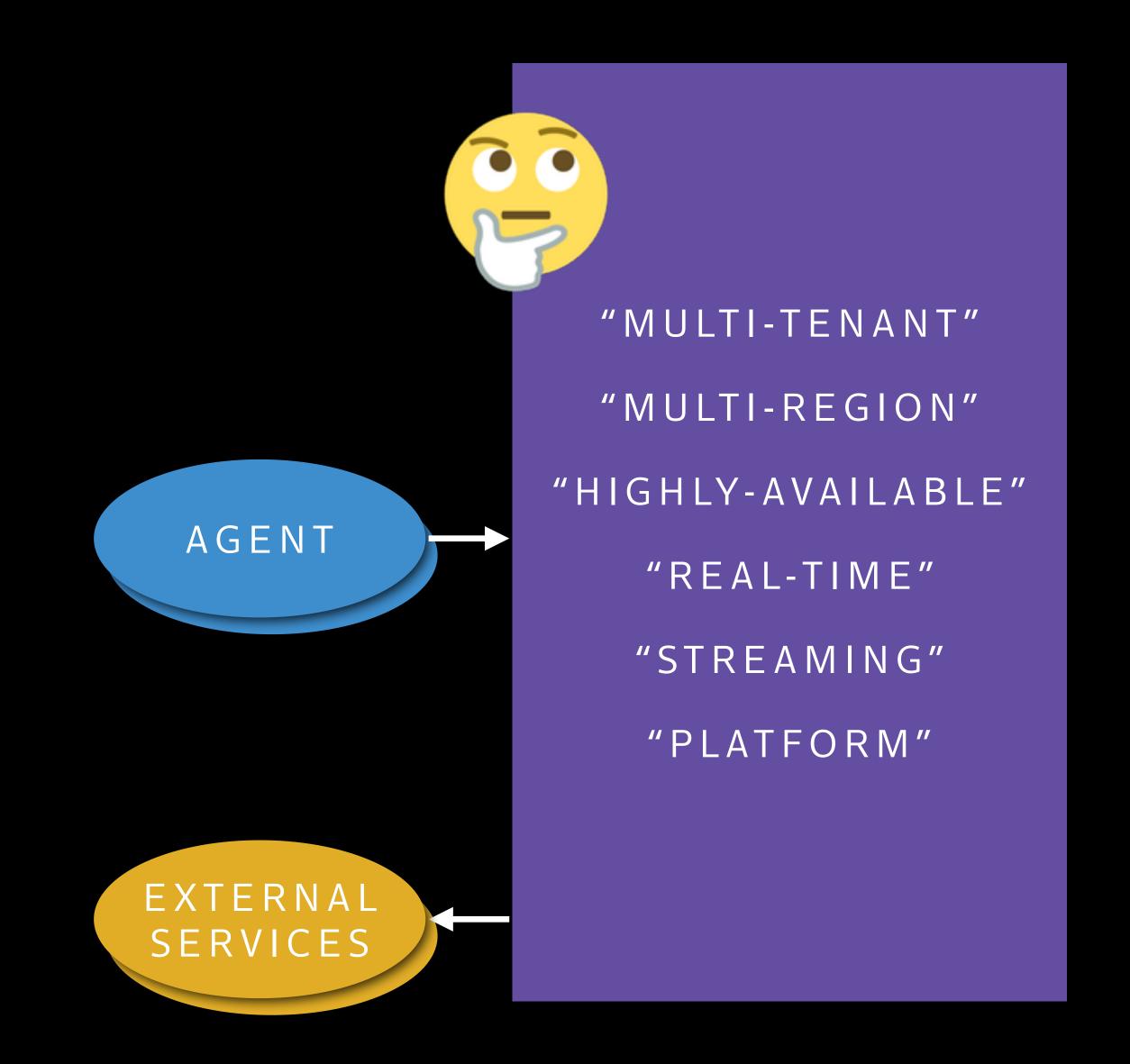






How do cluster nodes find each other?

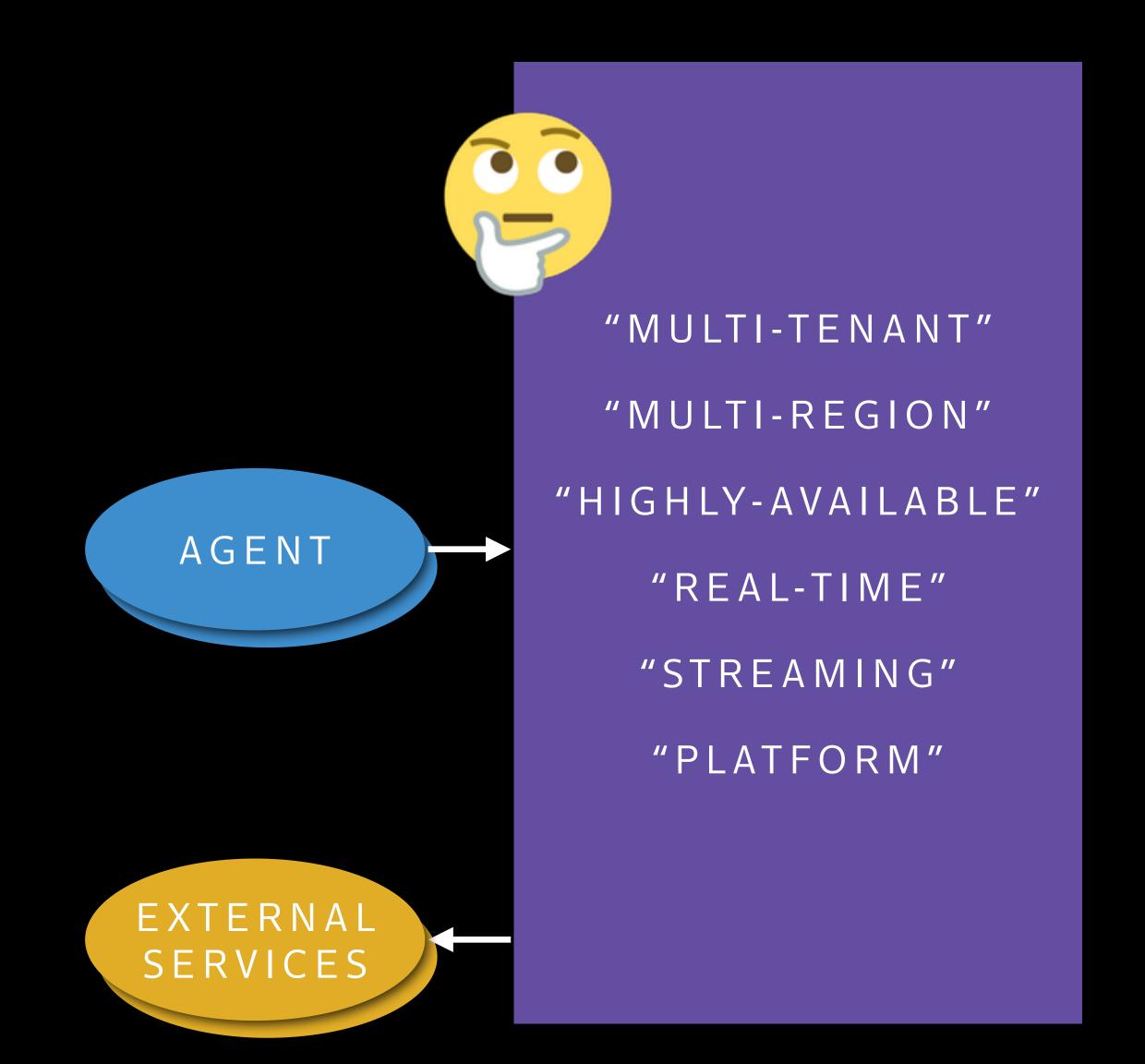




How do cluster nodes find each other?

Distribute code and configuration?





How do cluster nodes find each other?

Distribute code and configuration?

Know what happened when?



Where do agents send data?

AGENT

EXTERNAL
SERVICES

"MULTI-TENANT"

"MULTI-REGION"

"HIGHLY-AVAILABLE"

"REAL-TIME"

"STREAMING"

"PLATFORM"

How do cluster nodes find each other?

Distribute code and configuration?

Know what happened when?



Where do agents send data?

AGENT

How to get fault-tolerance without spam?

EXTERNAL
SERVICES

"MULTI-TENANT"

"MULTI-REGION"

"HIGHLY-AVAILABLE"

"REAL-TIME"

"STREAMING"

"PLATFORM"

How do cluster nodes find each other?

Distribute code and configuration?

Know what happened when?





Cluster membership and discovery



- Cluster membership and discovery
- Code and configuration dissemination

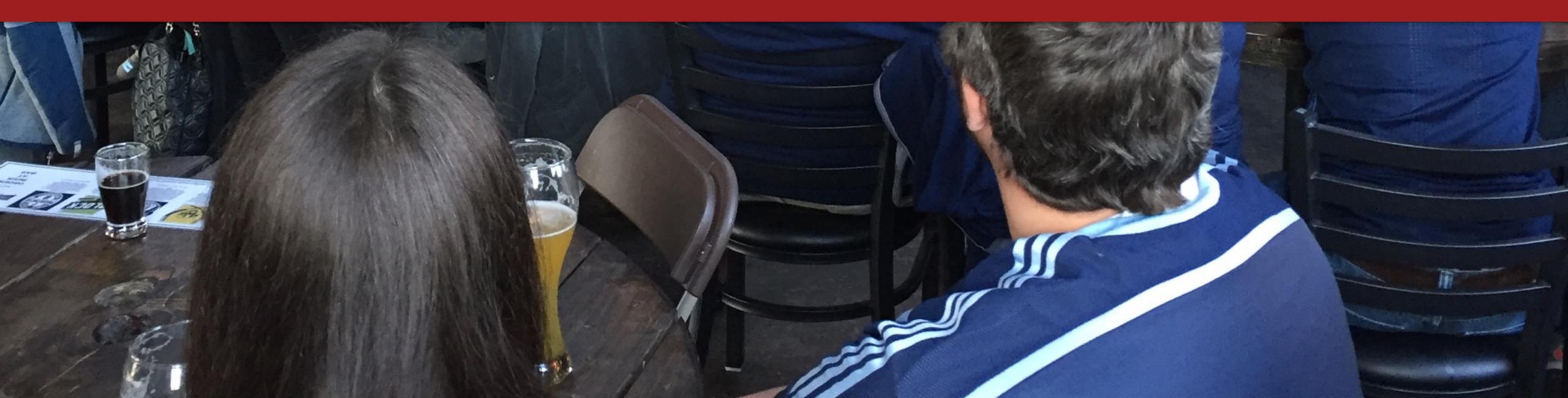


- Cluster membership and discovery
- Code and configuration dissemination
- Relative and convergent time





MEMBERSHIP PROTOCOLS



WHY NOT ZOOKEEPER/CONSUL/ETCD?





Connectedness



- Connectedness
- Balance



- Connectedness
- Balance
- Short path-length



Connectedness

Low clustering

- Balance
- Short path-length



- Connectedness
- Balance
- Short path-length

- Low clustering
- Scalability



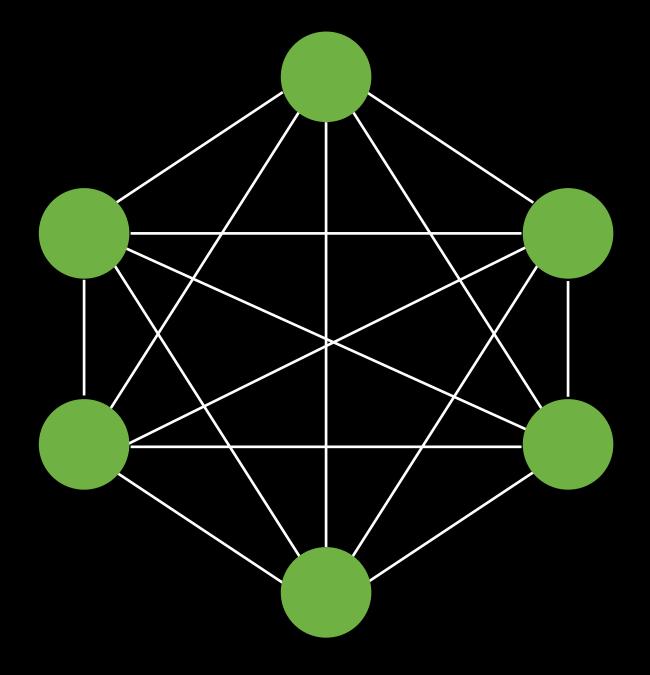
- Connectedness
- Balance
- Short path-length

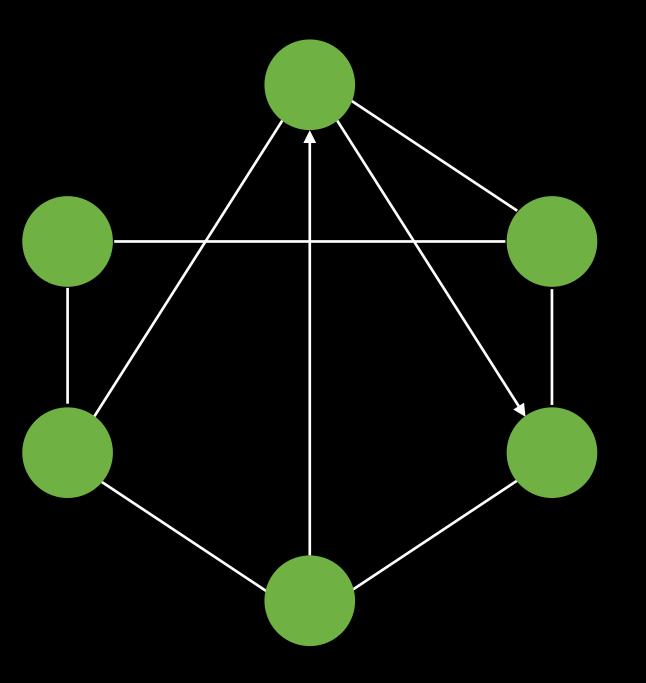
- Low clustering
- Scalability
- Accuracy



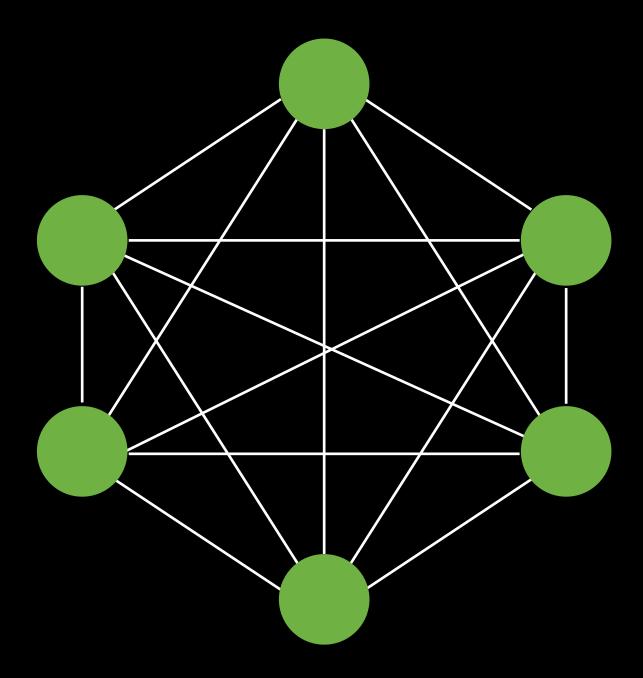
MEMBERSHIP: "VIEW" FLAVORS

Full

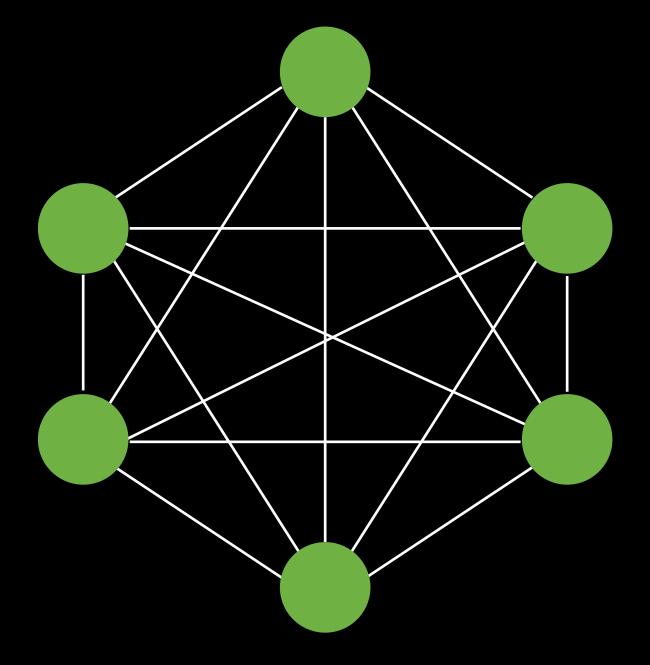






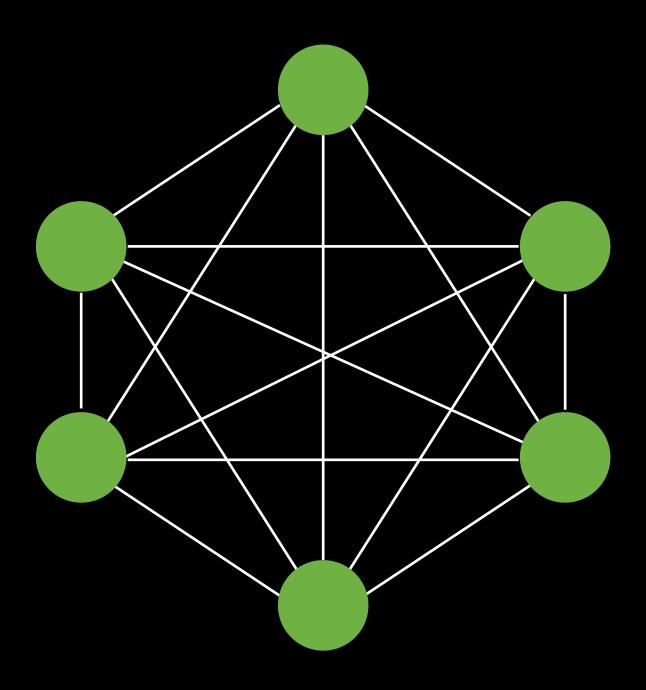






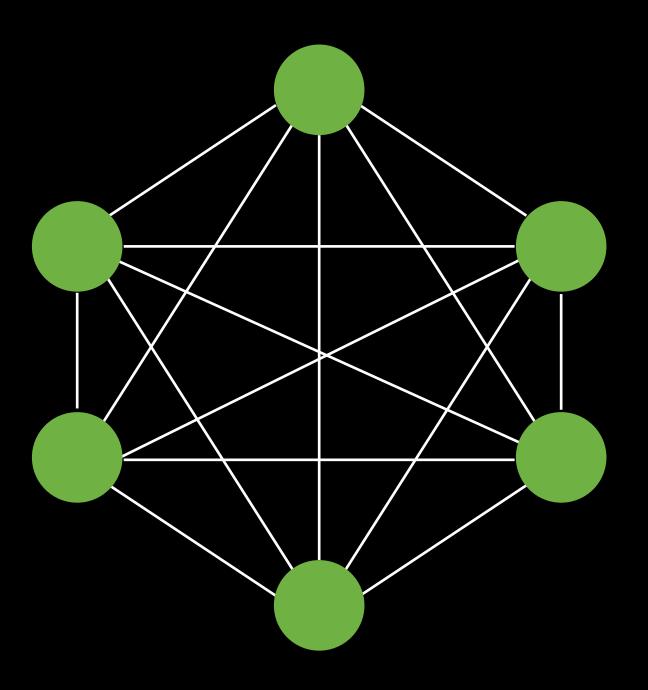
✓ Connectedness





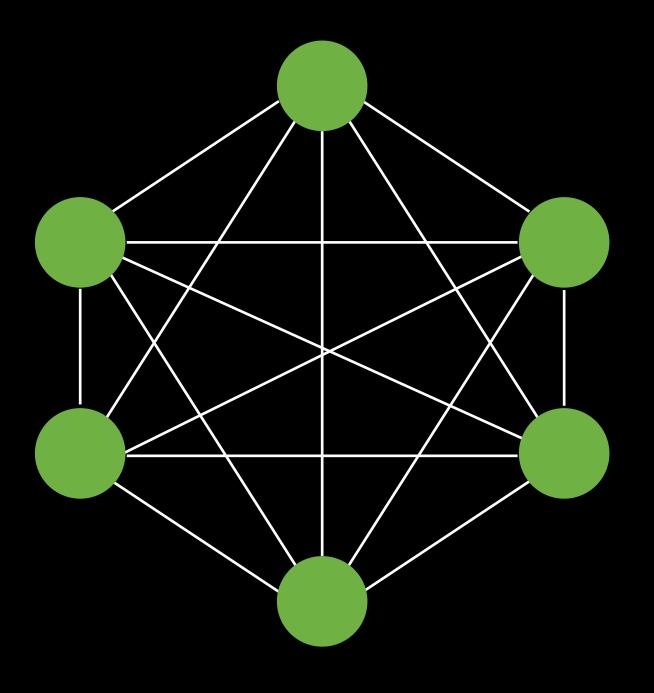
- ✓ Connectedness
- ✓ Short path-length





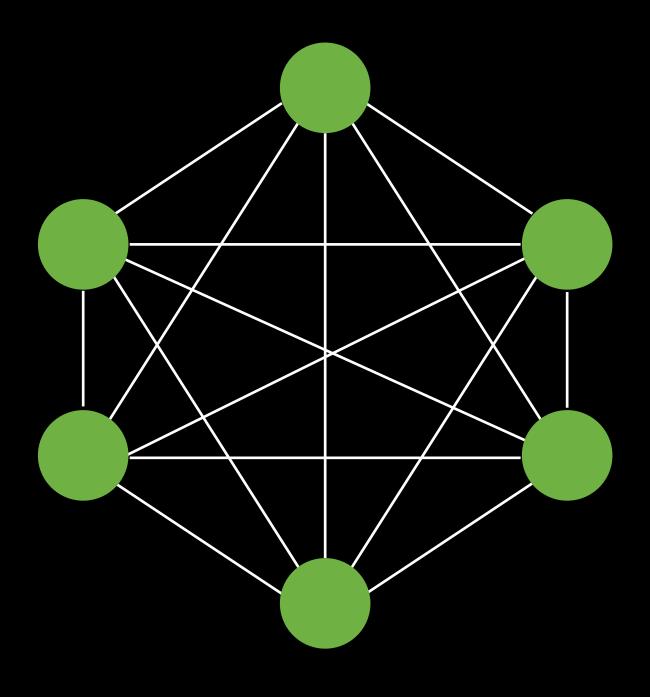
- ✓ Connectedness
- ✓ Short path-length
- Accuracy





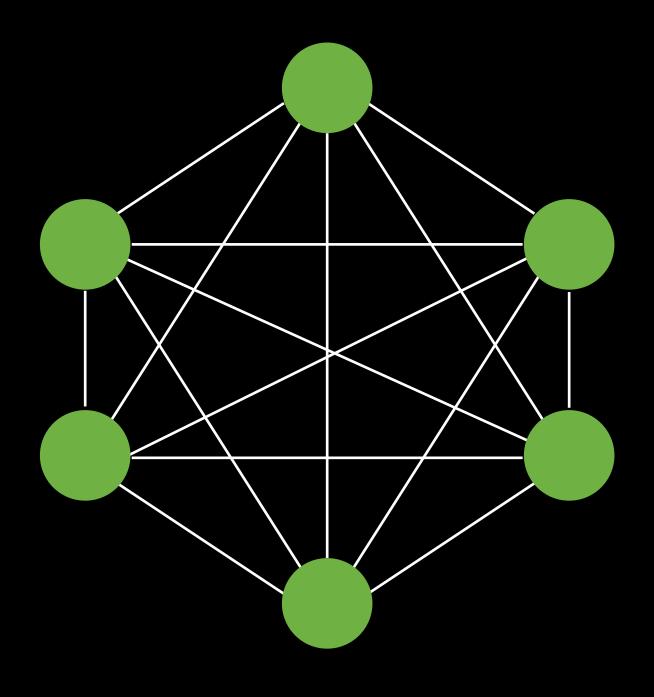
- ✓ Connectedness
- ✓ Short path-length
- Accuracy
- Balance





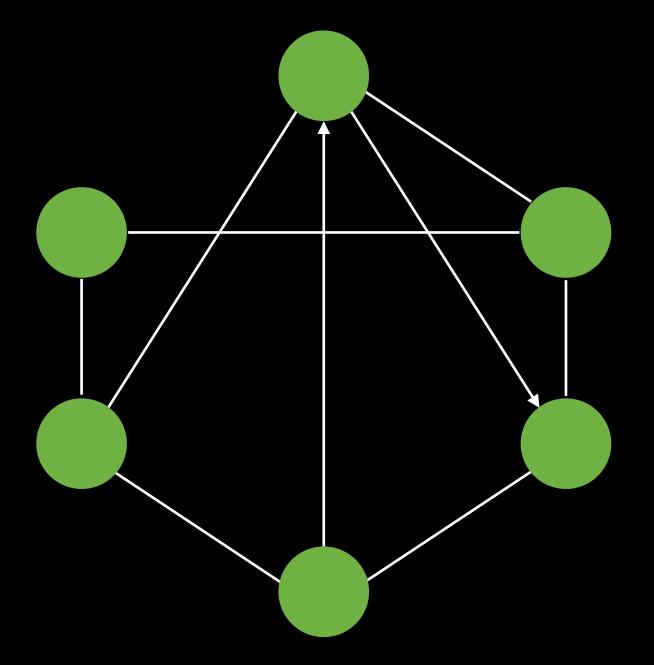
- ✓ Connectedness
- ✓ Short path-length
- Accuracy
- Balance
- High Clustering





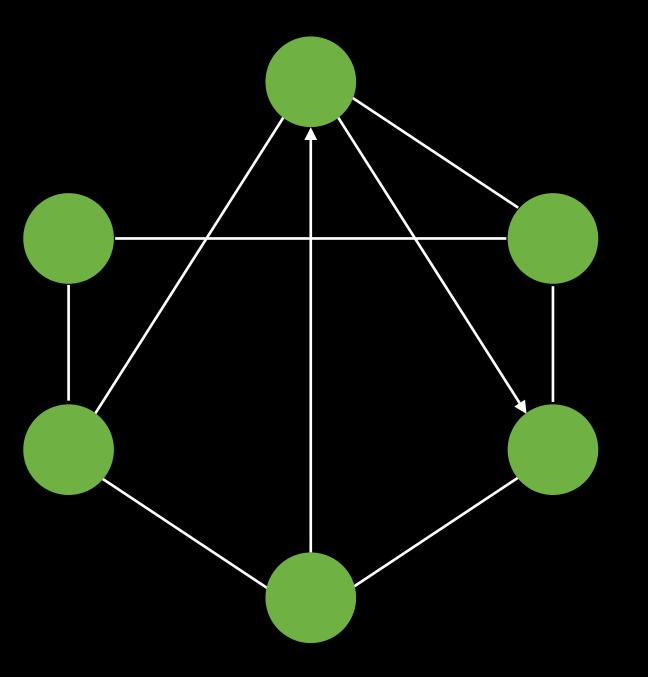
- ✓ Connectedness
- ✓ Short path-length
- Accuracy
- Balance
- High Clustering
- Low Scalability





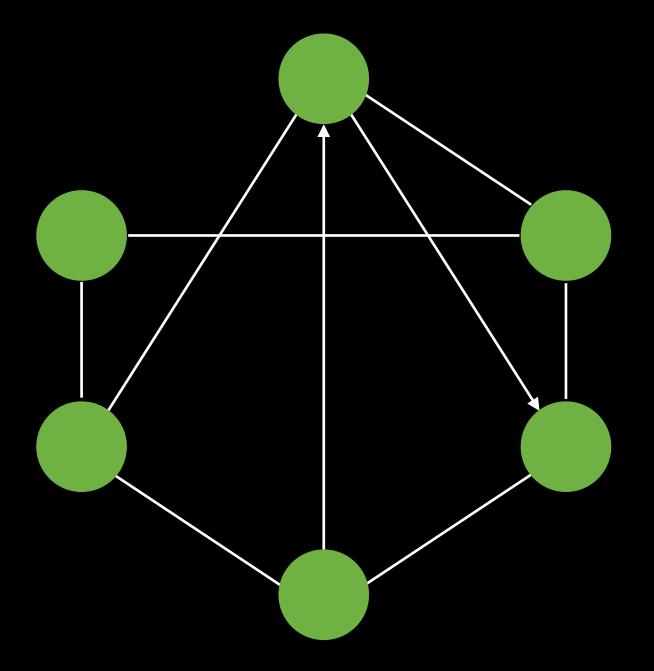


✓ Low Clustering



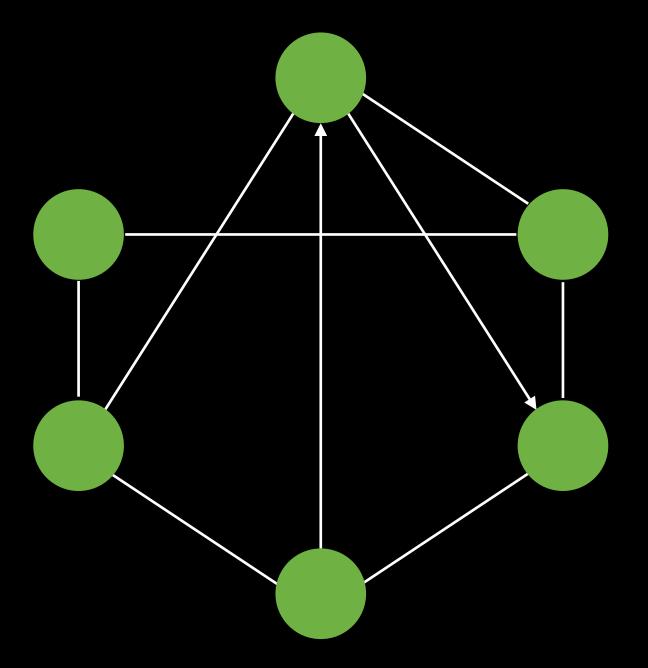


- ✓ Low Clustering
- ✓ High scalability



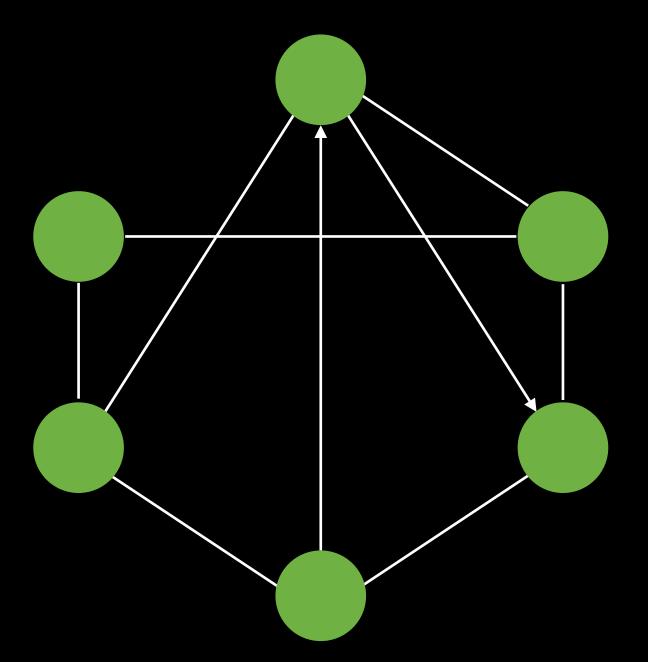


- ✓ Low Clustering
- ✓ High scalability
- Connectedness



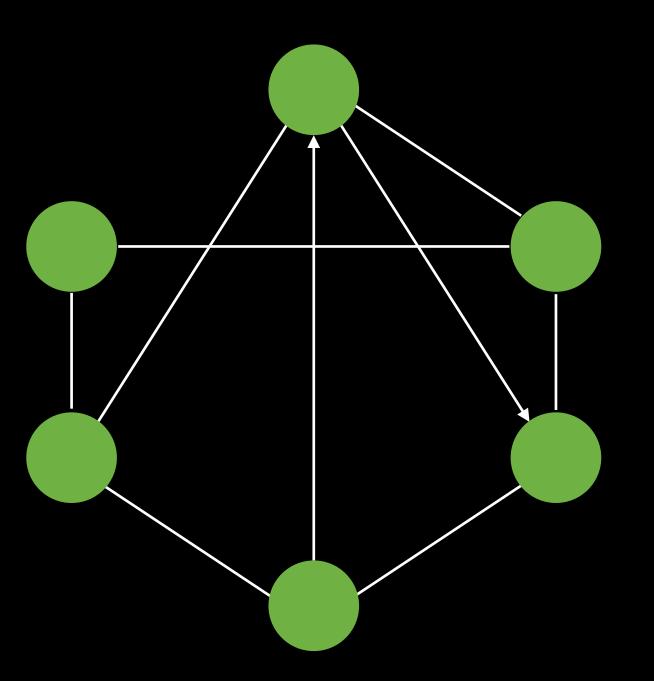


- ✓ Low Clustering
- High scalability
- Connectedness
- Balance



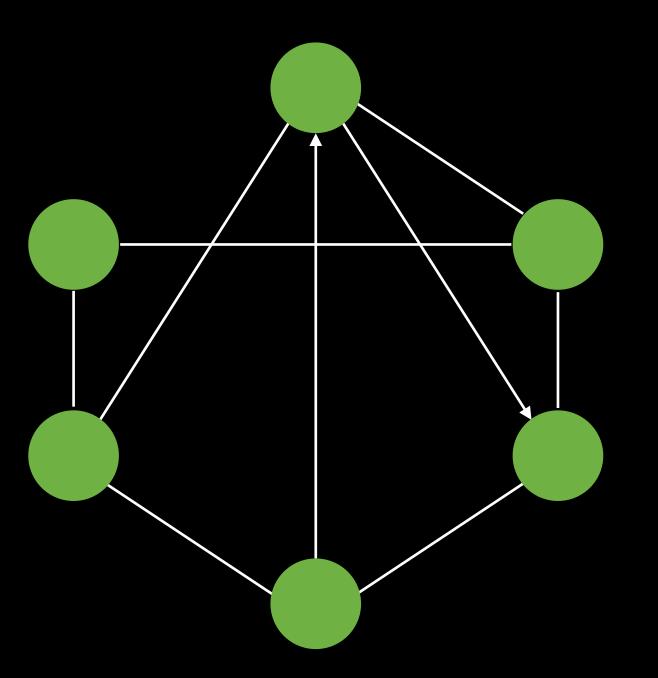


- ✓ Low Clustering
- High scalability
- Connectedness
- Balance
- Path-length





- ✓ Low Clustering
- High scalability
- Connectedness
- Balance
- Path-length
- Accuracy







SWIM: Scalable Weakly-consistent Infection-style Process Group Membership Protocol

Abhinandan Das, Indranil Gupta, Ashish Motivala*
Dept. of Computer Science, Cornell University
Ithaca NY 14853 USA

{asdas,gupta,ashish}@cs.cornell.edu















Quadratic load







- Quadratic load
- Failure detection







- Quadratic load
- Failure detection
 - Response times







- Quadratic load
- Failure detection
 - Response times
 - False positives









- Failure detection
 - Response times
 - False positives



SWIM solutions







- Quadratic load
- Failure detection
 - Response times
 - False positives



SWIM solutions

Separate membership and failure detection







- Quadratic load
- Failure detection
 - Response times
 - False positives



SWIM solutions

- Separate membership and failure detection
- Randomized probing







- Quadratic load
- Failure detection
 - Response times
 - False positives



SWIM solutions

- Separate membership and failure detection
- Randomized probing
- Piggyback membership on probes



SWIM - 2002

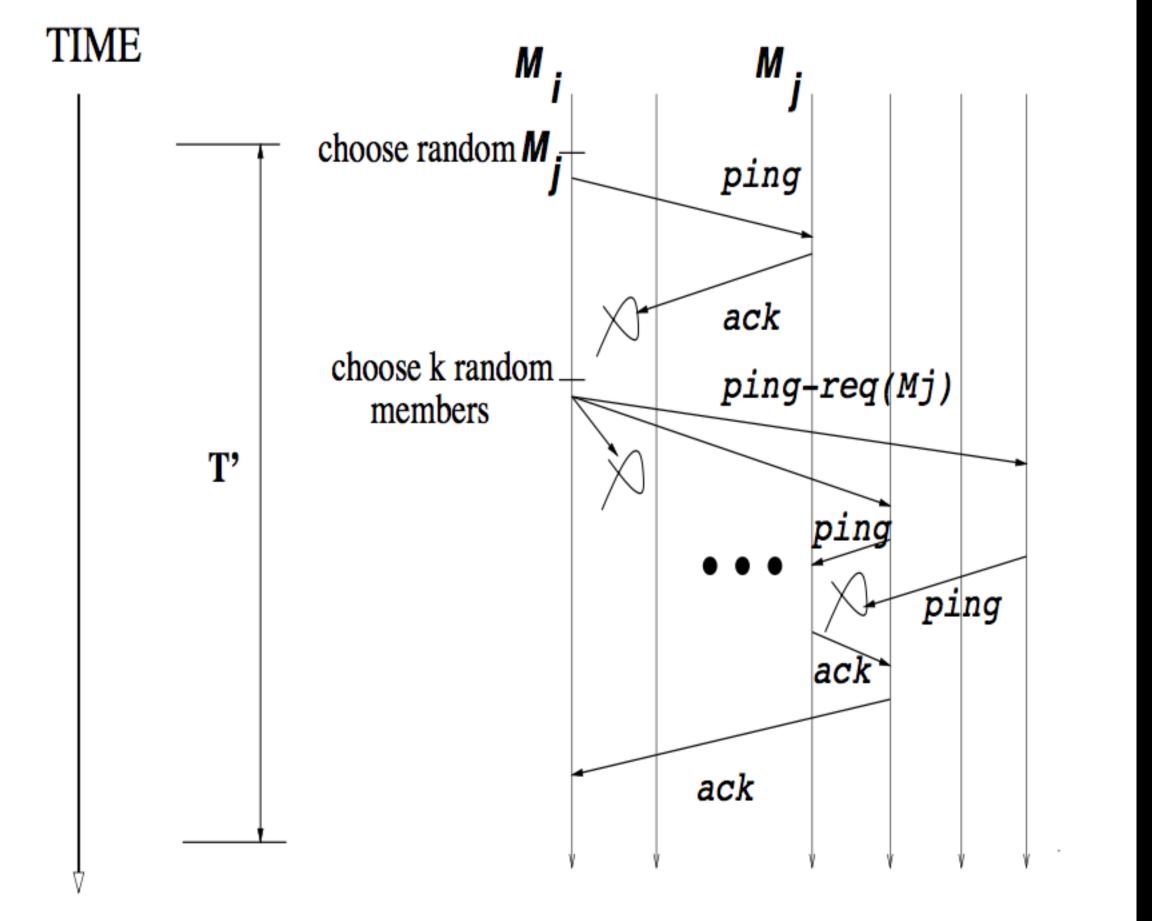


Figure 1. SWIM failure detection: Example protocol period at M_i . This shows all the possible messages that a protocol period may initiate. Some message contents excluded for simplicity.





Peer-to-Peer Membership Management for Gossip-Based Protocols

Ayalvadi J. Ganesh, Anne-Marie Kermarrec, and Laurent Massoulié



SCAMP - 2003





Full views limit scalability



- Full views limit scalability
- Flexible partial-view size, asymmetric



SCAMP - 2003

- Full views limit scalability
- → Flexible partial-view size, asymmetric
- Reactive view management



SCAMP - 2003

- Full views limit scalability
- → Flexible partial-view size, asymmetric
- Reactive view management
- Join ("subscribe") via random walk



- Full views limit scalability
- → Flexible partial-view size, asymmetric
- Reactive view management
- Join ("subscribe") via random walk
- Automatic balancing via indirection and leases



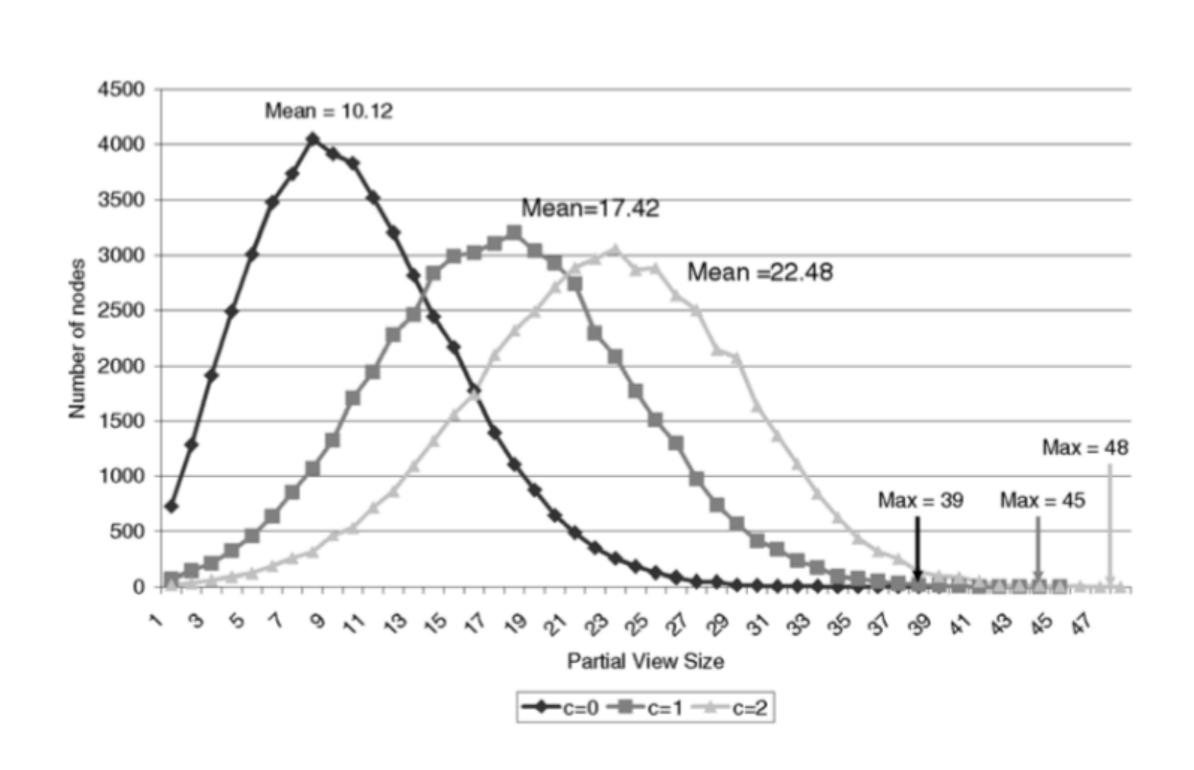


Fig. 5. Impact of c on the partial view size distribution in a 50,000 node group.



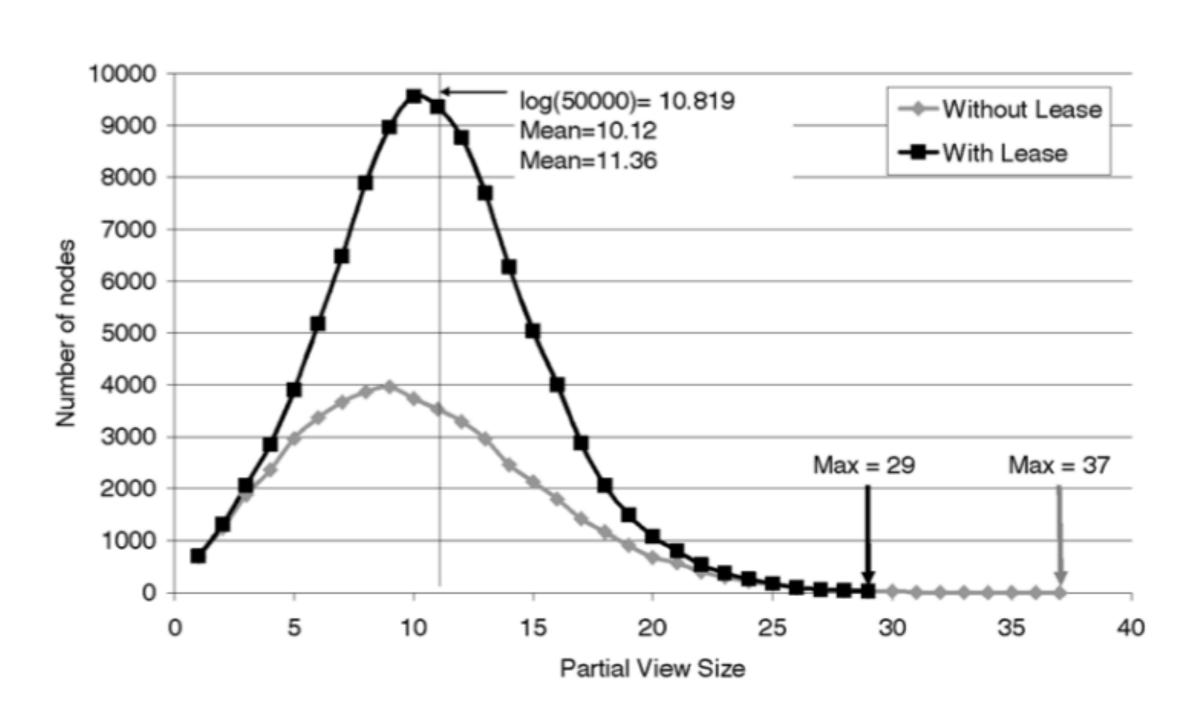


Fig. 8. Distribution of the partial views by size in a 50,000 node group with and without the lease mechanism.





CYCLON: Inexpensive Membership Management for Unstructured P2P Overlays

Spyros Voulgaris,^{1,2} Daniela Gavidia,¹ and Maarten van Steen¹



CYCLON - 2005





Random shuffling doesn't create good balance



- Random shuffling doesn't create good balance
- → Fixed partial-view size, symmetric



- Random shuffling doesn't create good balance
- → Fixed partial-view size, symmetric
- Cyclic view management



- Random shuffling doesn't create good balance
- → Fixed partial-view size, symmetric
- Cyclic view management
- Join via random walk



CYCLON - 2005

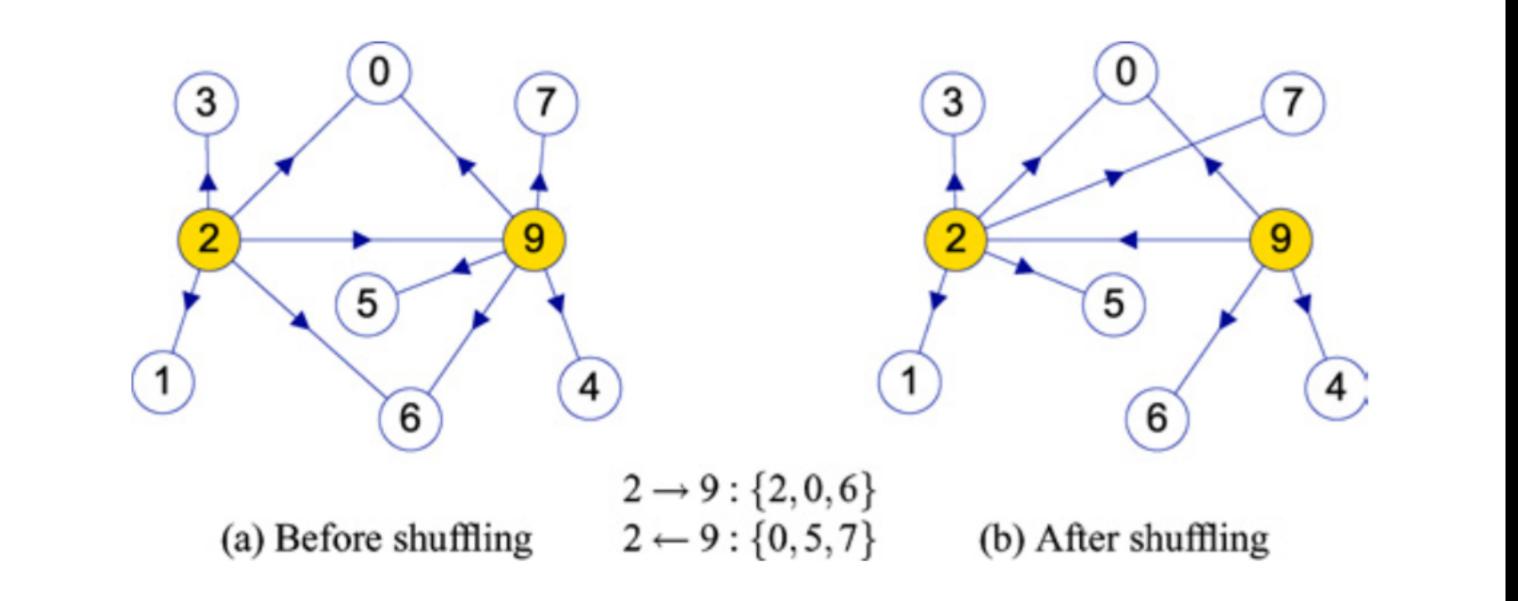


Fig. 1. An example of shuffling between nodes 2 and 9. Note that, among other changes, the link between 2 and 9 reverses direction.



CYCLON - 2005

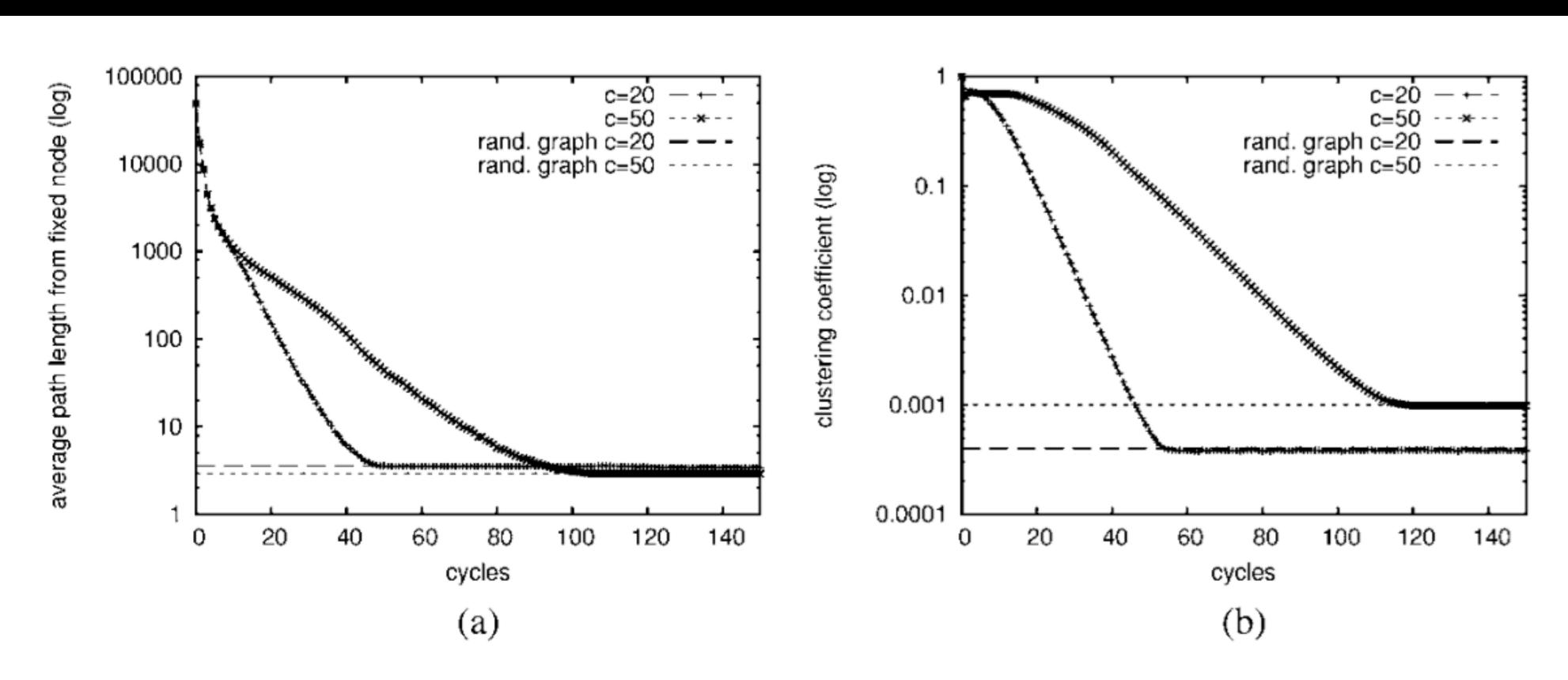


Fig. 2. (a) Average shortest path length between two nodes for different cache sizes. (b) Average clustering coefficient taken over all nodes.



- No failure detectors
- SCAMP: asymmetric views \Rightarrow disconnection
- SCAMP: unbounded view size ⇒ imbalance





HyParView: a membership protocol for reliable gossip-based broadcast

João Leitão José Pereira Luís Rodrigues



HYPARVIEW - 2007





Fanout is related to reliability





- Fanout is related to reliability
- High failure rates decrease quality





- Fanout is related to reliability
- High failure rates decrease quality



- Fanout is related to reliability
- High failure rates decrease quality
- TCP for transport and failure detector



- Fanout is related to reliability
- High failure rates decrease quality
- TCP for transport and failure detector
- → Small reactive view ("active")



- Fanout is related to reliability
- High failure rates decrease quality
- → TCP for transport and failure detector
- → Small reactive view ("active")
- → Larger cyclic view ("passive")



- Fanout is related to reliability
- High failure rates decrease quality
- → TCP for transport and failure detector
- Small reactive view ("active")
- → Larger cyclic view ("passive")
- Join and shuffle via random walk



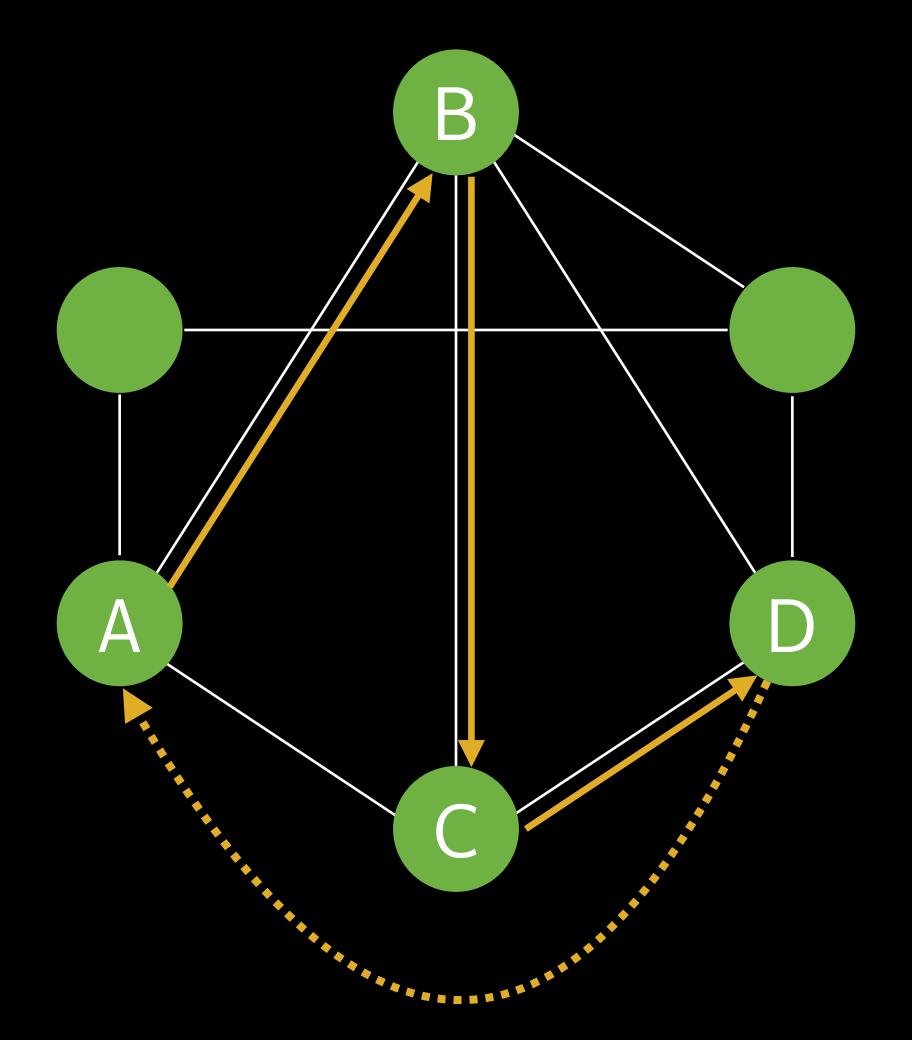


Algorithm 1: Membership Operations

```
upon init do
     Send(Join, contactNode, myself);
upon Receive(Join, newNode) do
     if isfull(activeView) then
         trigger dropRandomElementFromActiveView
     activeView \leftarrow activeView \cup newNode
         foreach n \in \text{activeView and } n \neq \text{newNode do}
               Send(ForwardJoin, n, newNode, ARWL, myself)
upon Receive(FORWARDJOIN, newNode, timeToLive, sender) do
     if timeToLive== 0|#activeView== 0 then
         trigger addNodeActiveView(newNode)
     else
         if timeToLive==PRWL then
              trigger addNodePassiveView(newNode)
         n \leftarrow n \in \text{activeView and } n \neq \text{sender}
         Send(ForwardJoin, n, newNode, timeToLive-1, myself)
upon dropRandomElementFromActiveView do
     n \leftarrow n \in \text{activeView}
     Send(DISCONNECT, n, myself)
     activeView \leftarrow activeView \setminus \{n\}
     passiveView \leftarrow passiveView \cup \{n\}
```



Passive view maintenance





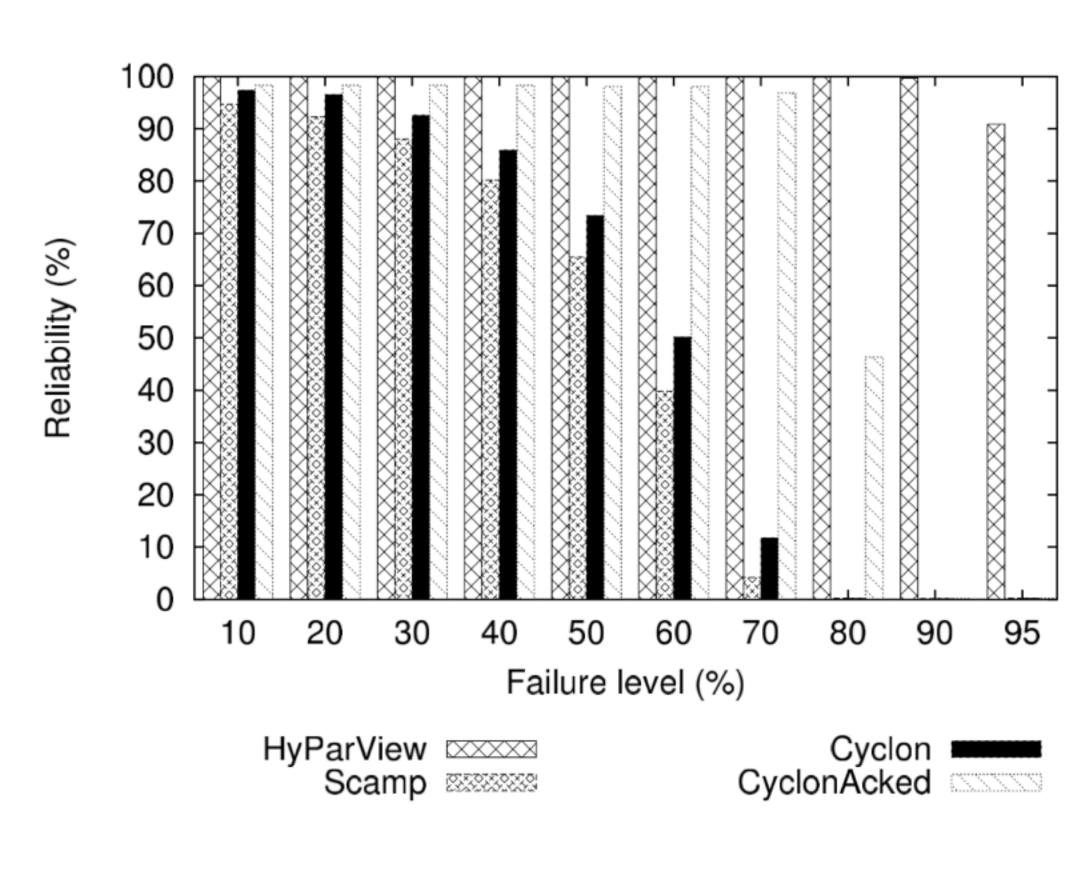
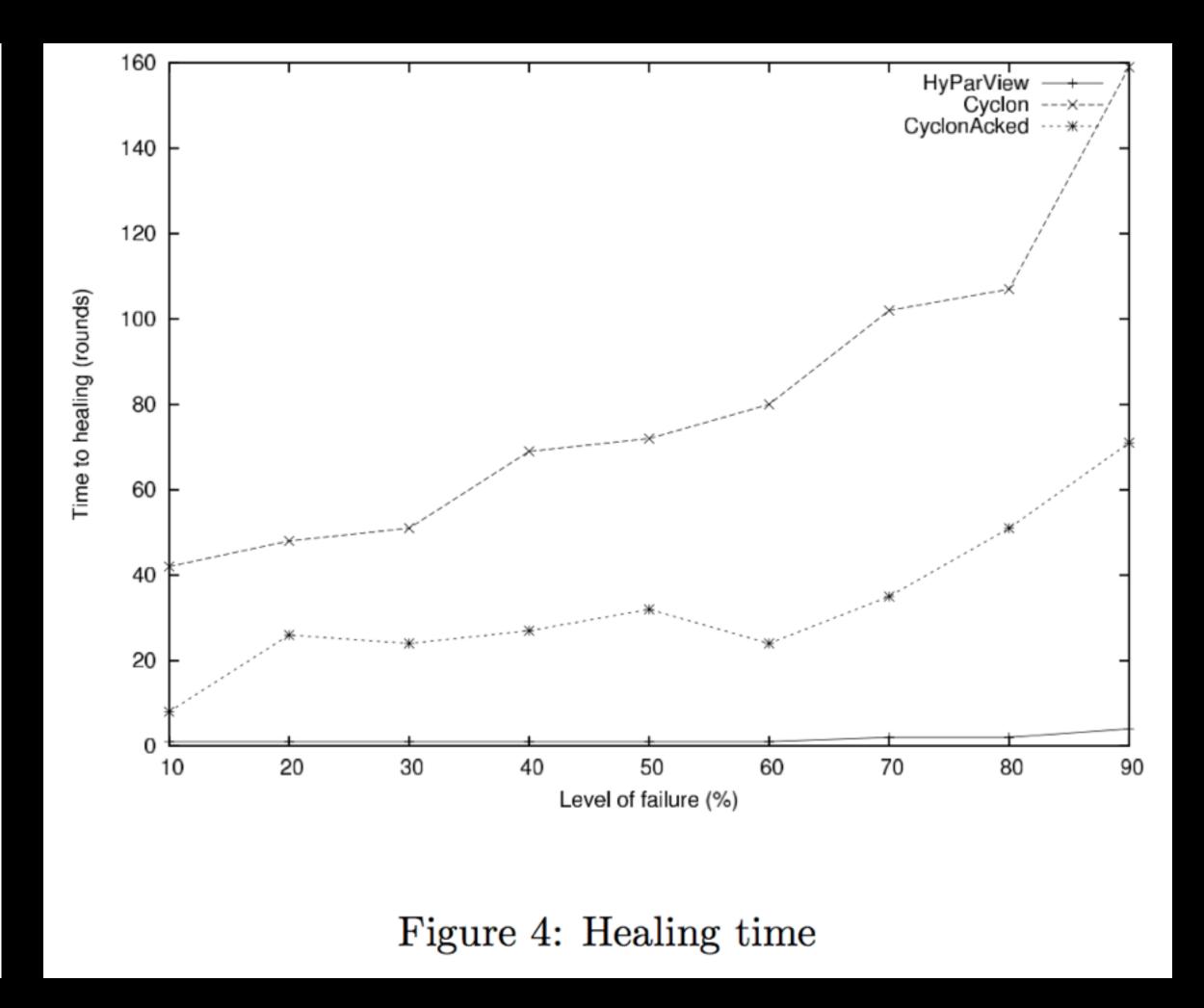


Figure 2: Reliability for 1000 messages

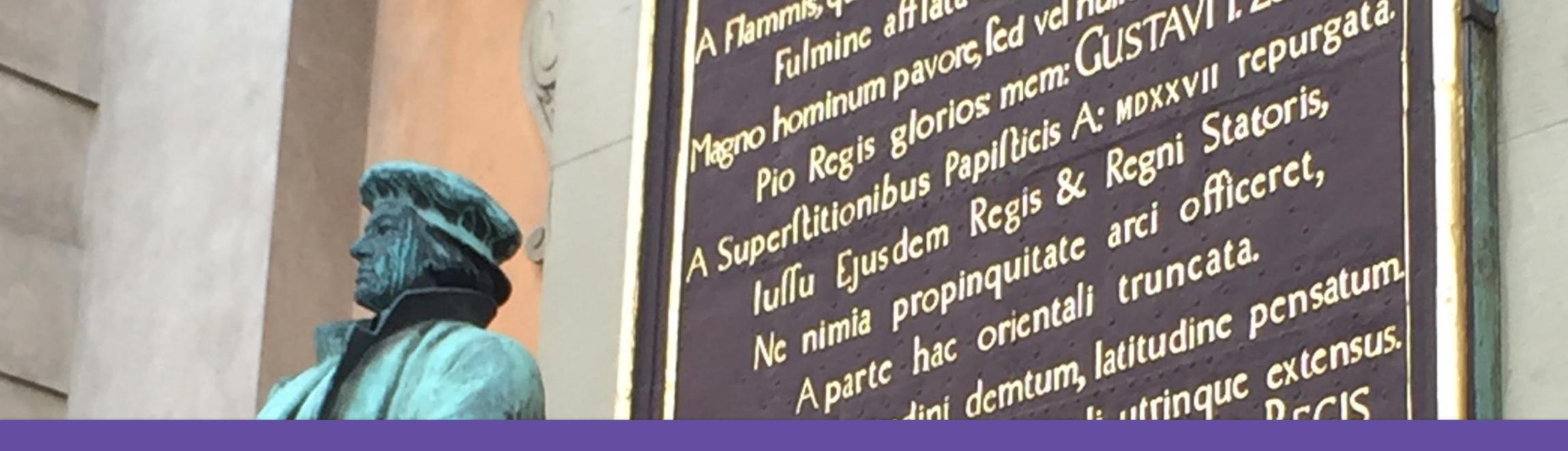




WE CHOSE HYPARVIEW

- Only active view maintenance
- Passive view maintains full membership (unbounded)
- Later: switch to complete passive maintenance





DISSEMINATION PROTOCOLS





Reliability

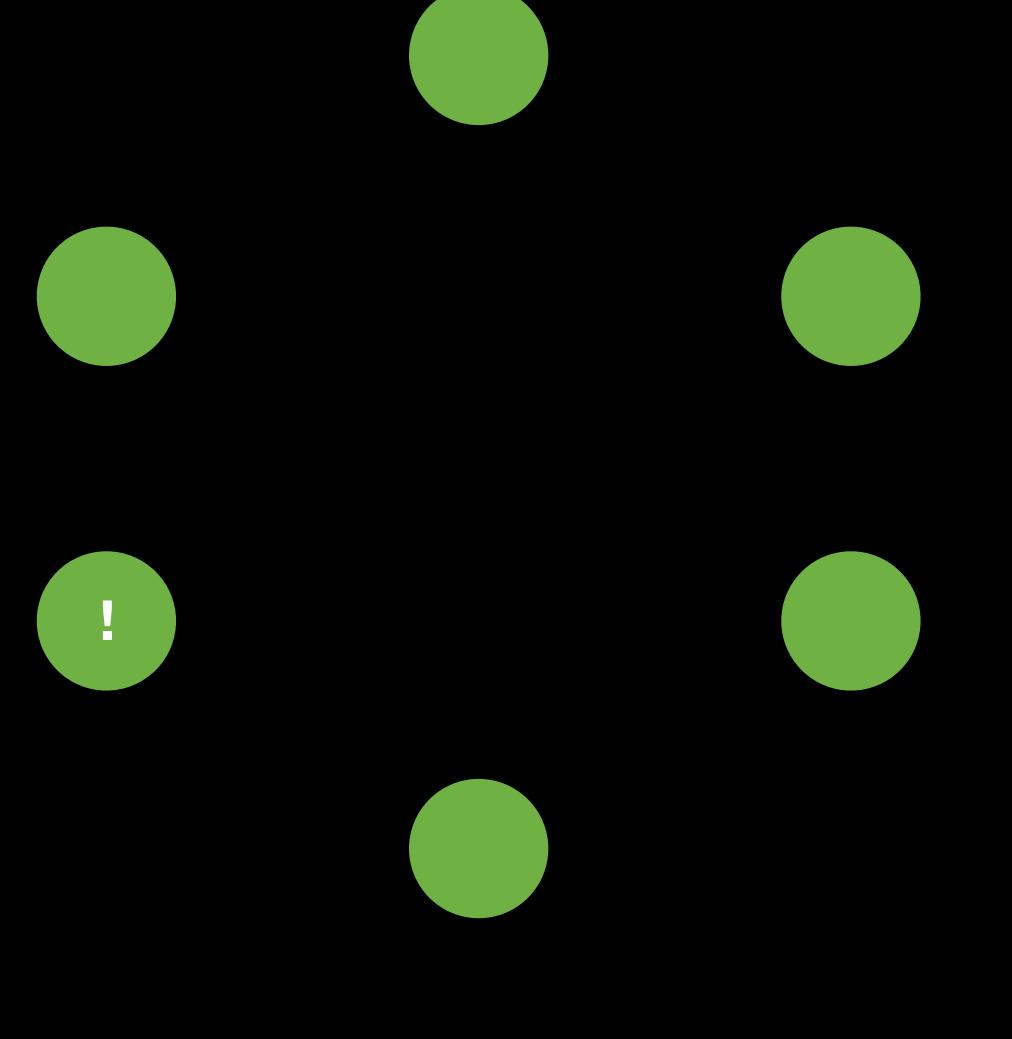


- Reliability
- Scalability



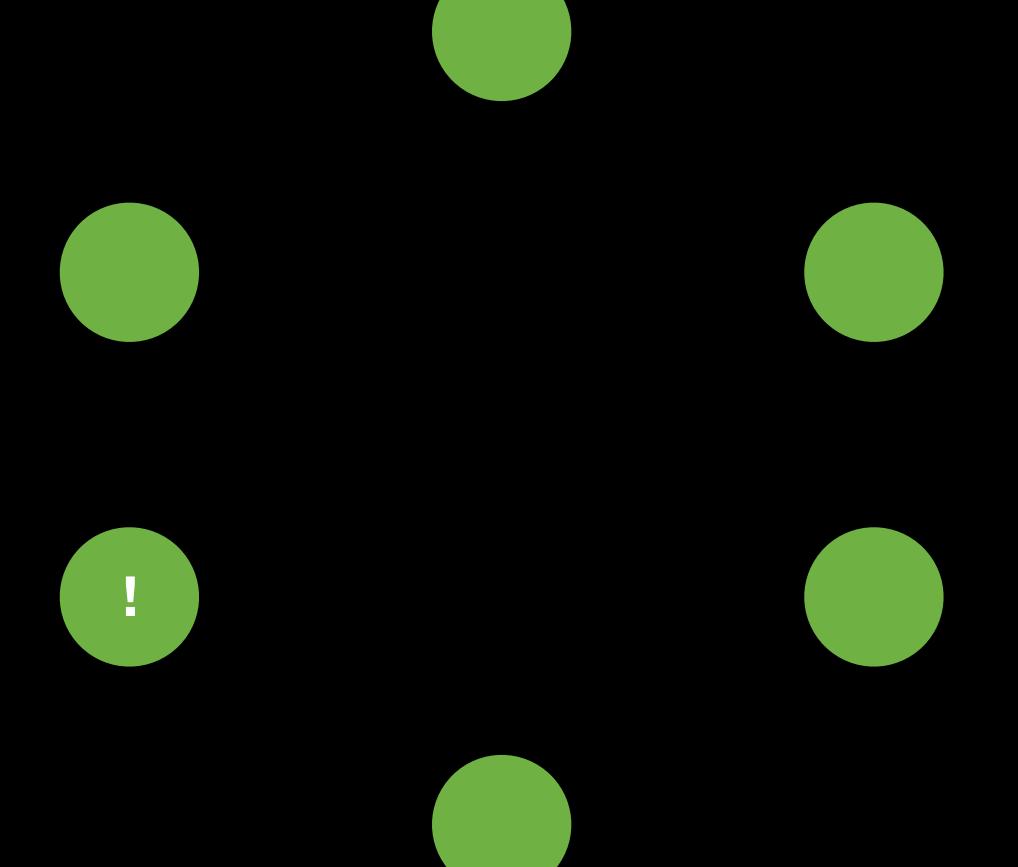
- Reliability
- Scalability
- Efficiency





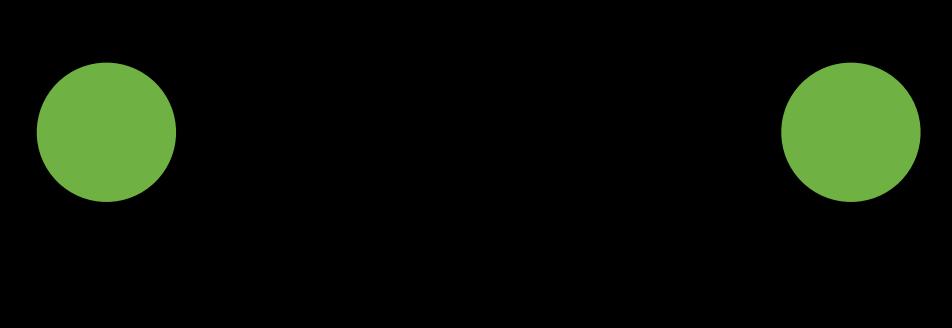


Send to random peers



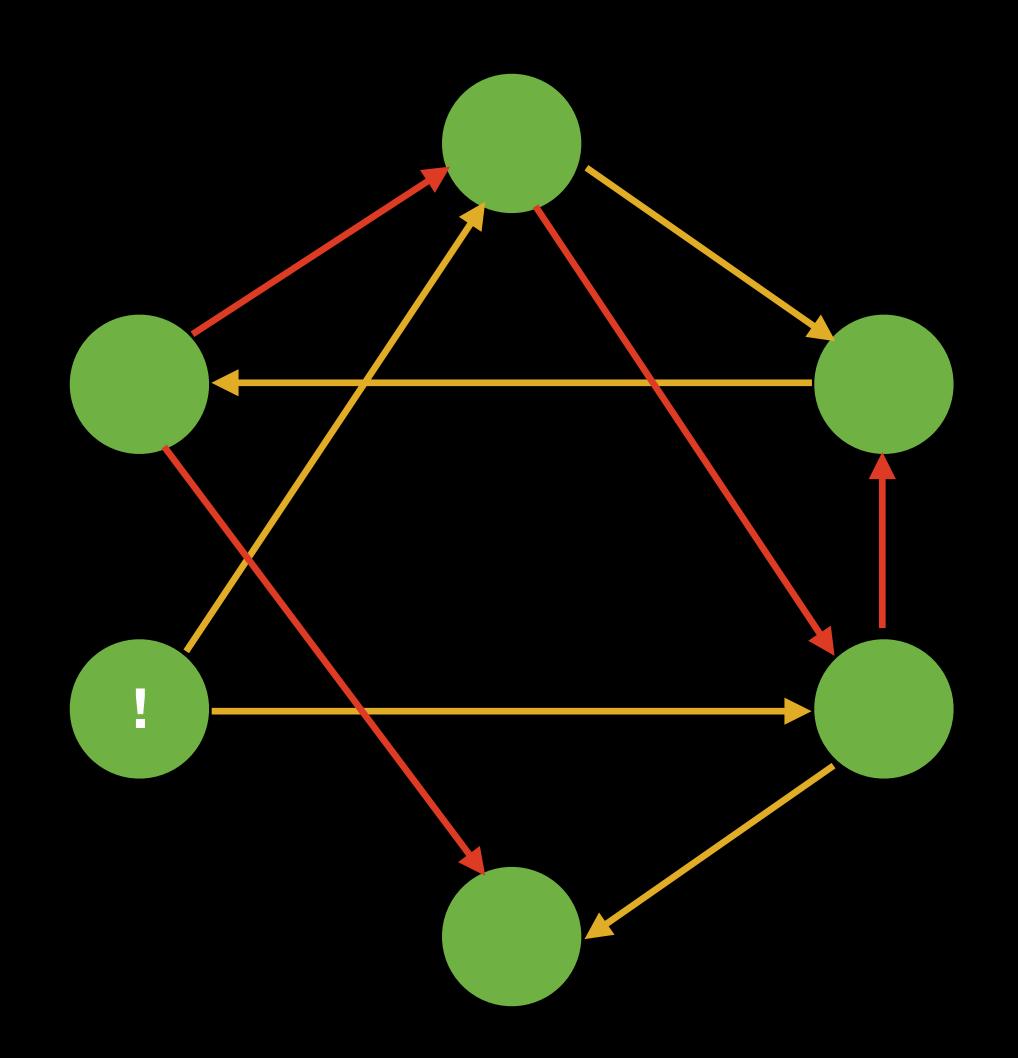


- Send to random peers
- Messages rebroadcast by recipients



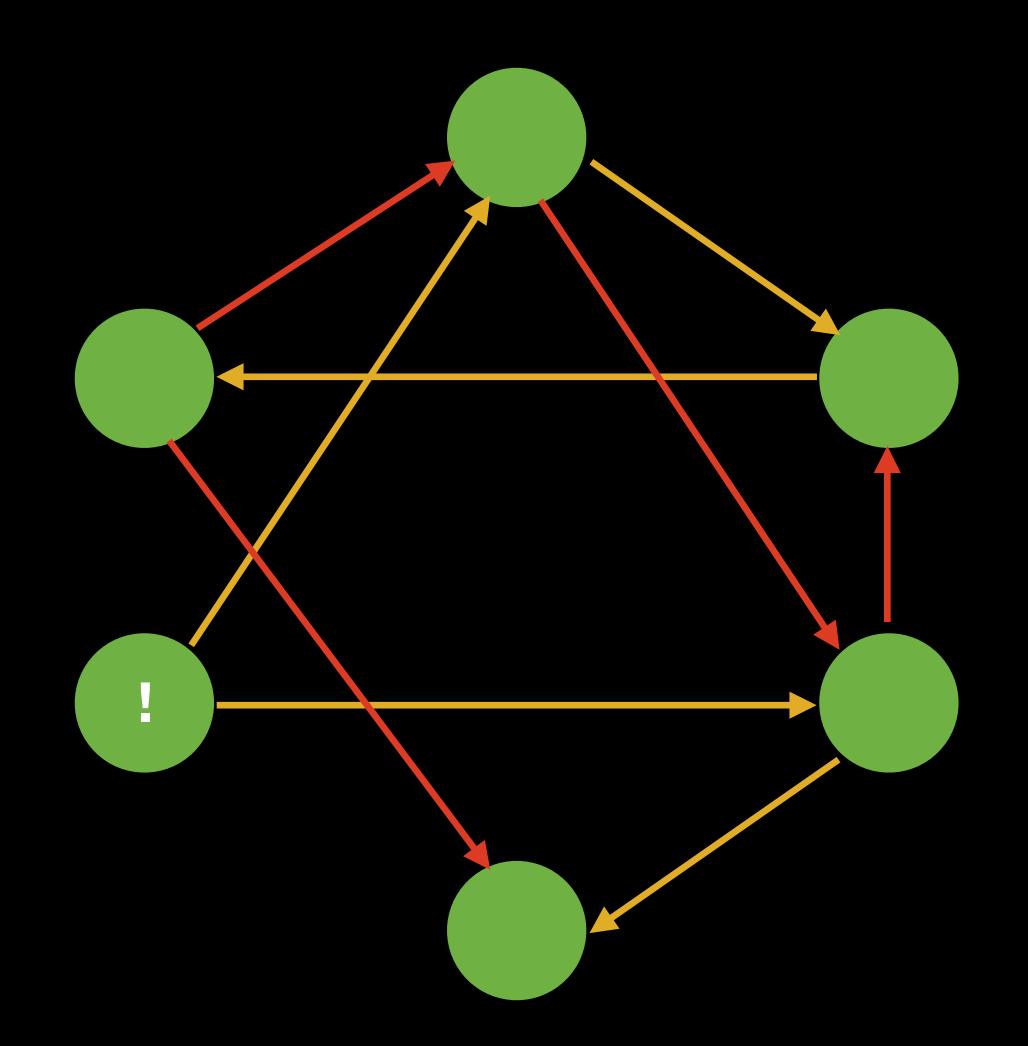


- Send to random peers
- Messages rebroadcast by recipients



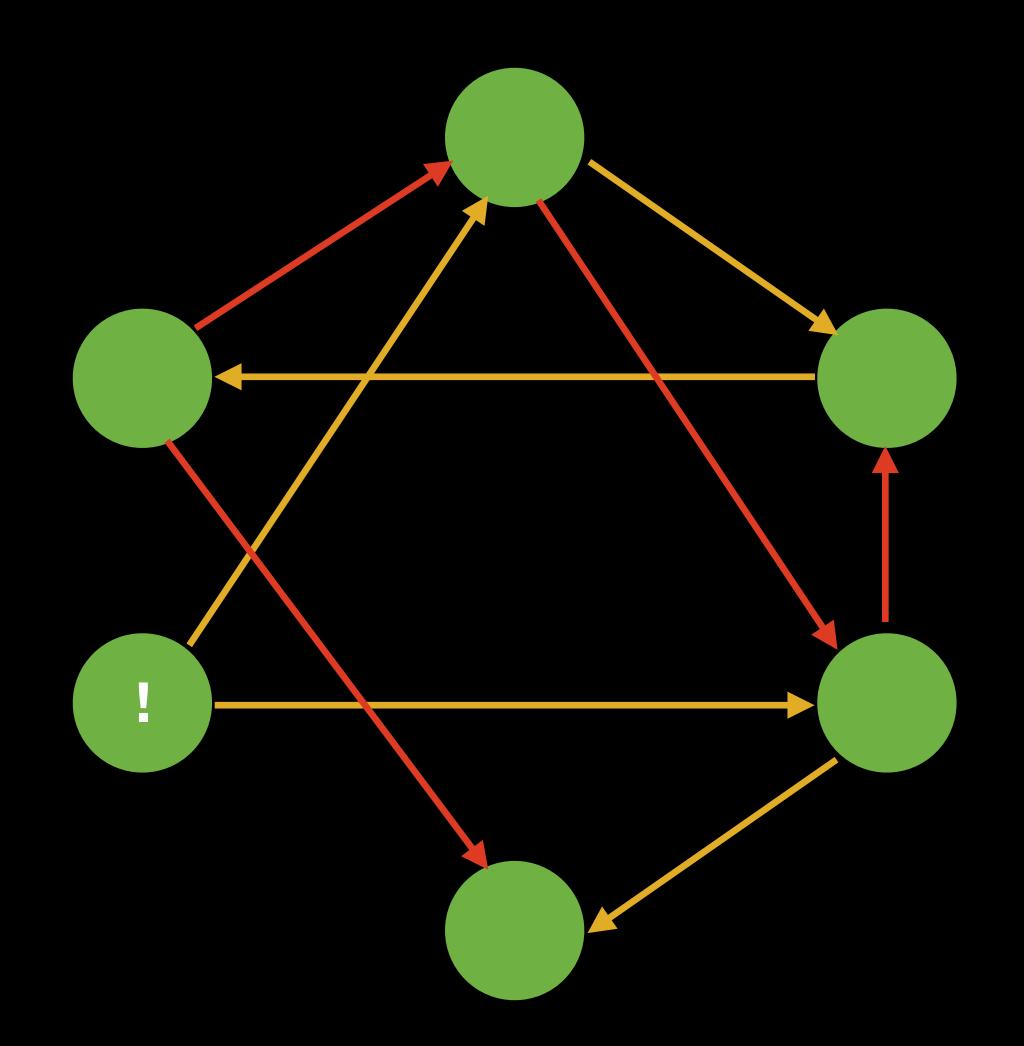


- Send to random peers
- Messages rebroadcast by recipients
- High redundancy





- Send to random peers
- Messages rebroadcast by recipients
- High redundancy
- Low scalability





INCREASED EFFICIENCY



Epidemic Broadcast Trees

João Leitão

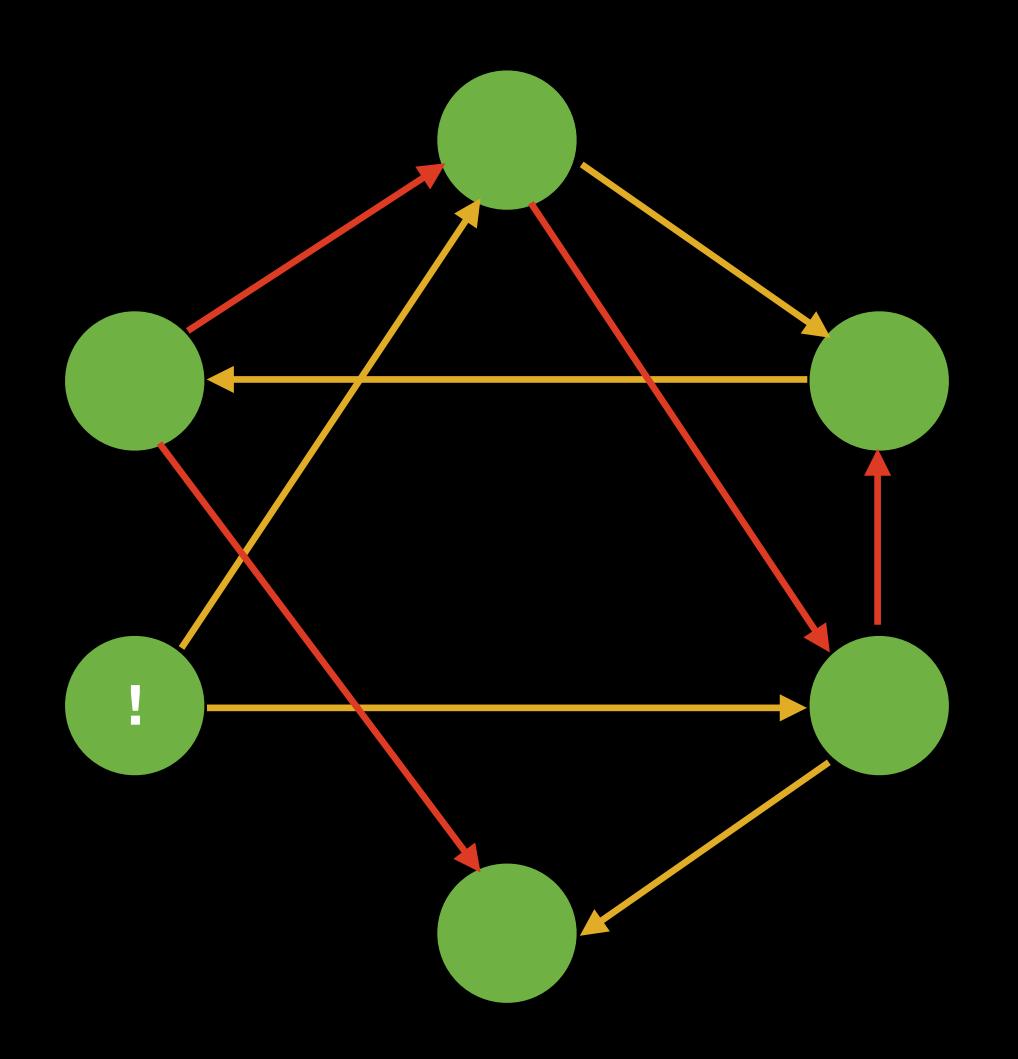
José Pereira

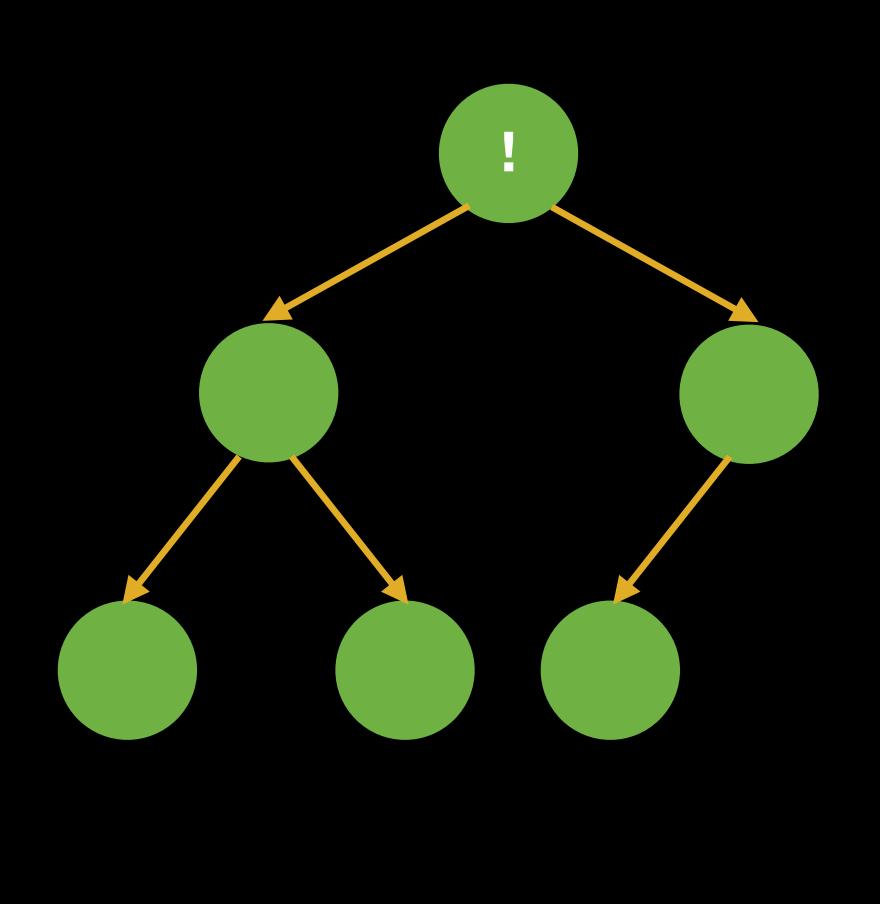
Luís Rodrigues

DI-FCUL TR-07-14

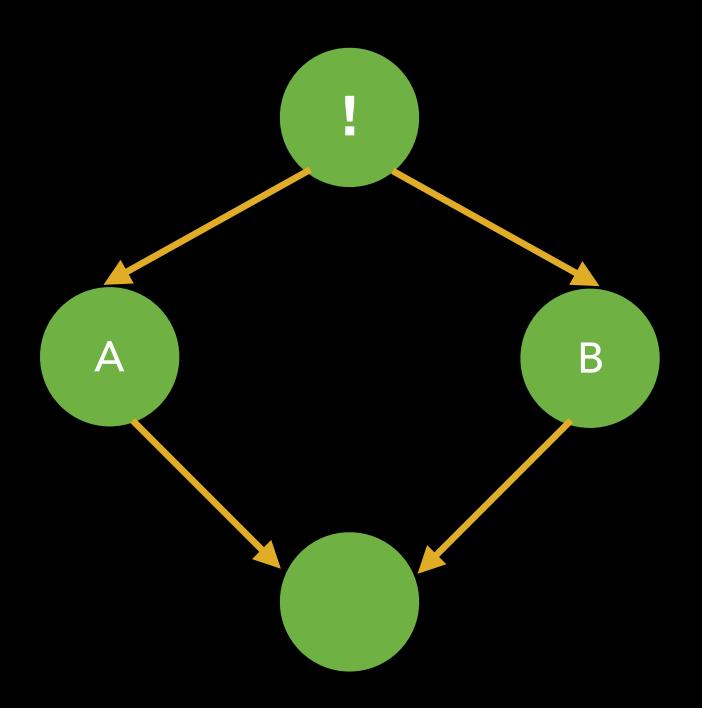


PLUMTREE - 2009



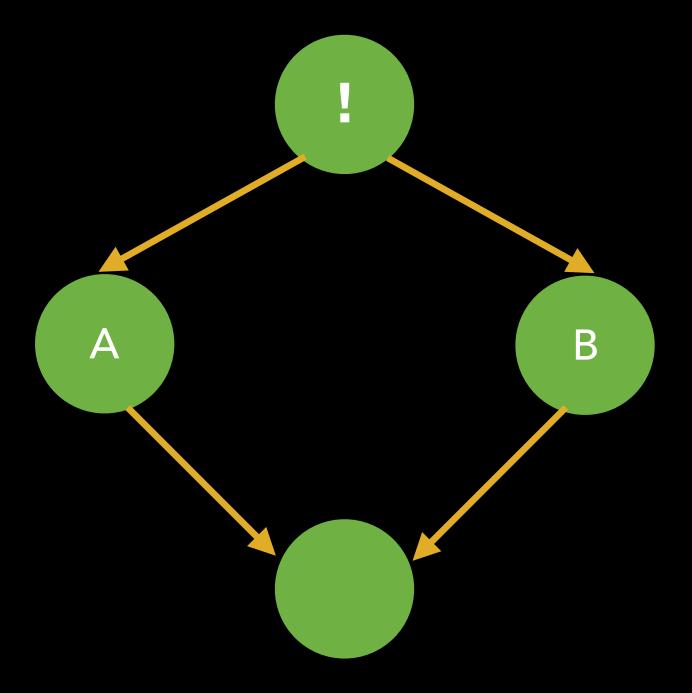






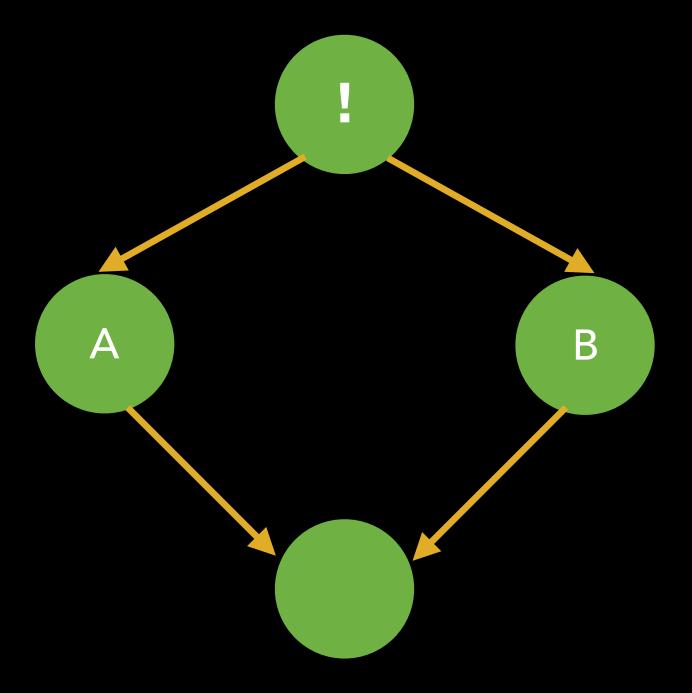


All nodes start with full "eager" set



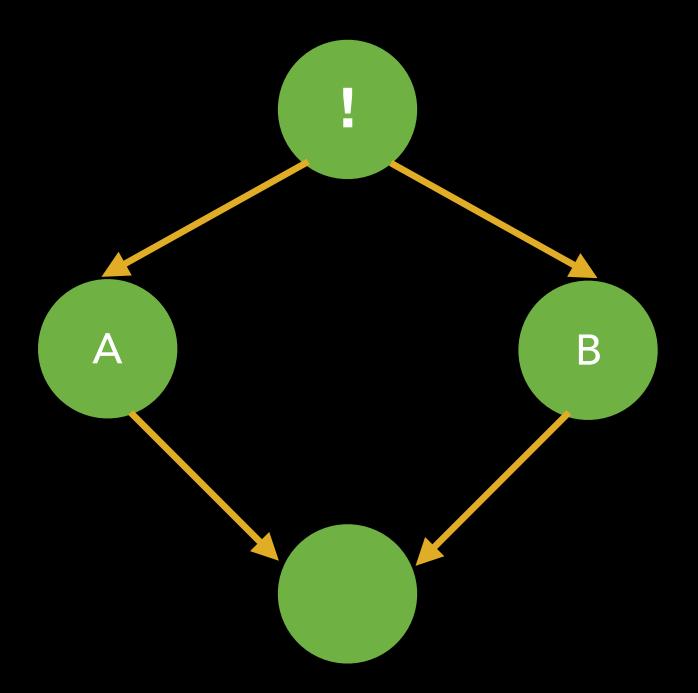


- All nodes start with full "eager" set
- Broadcast triggers eager-push



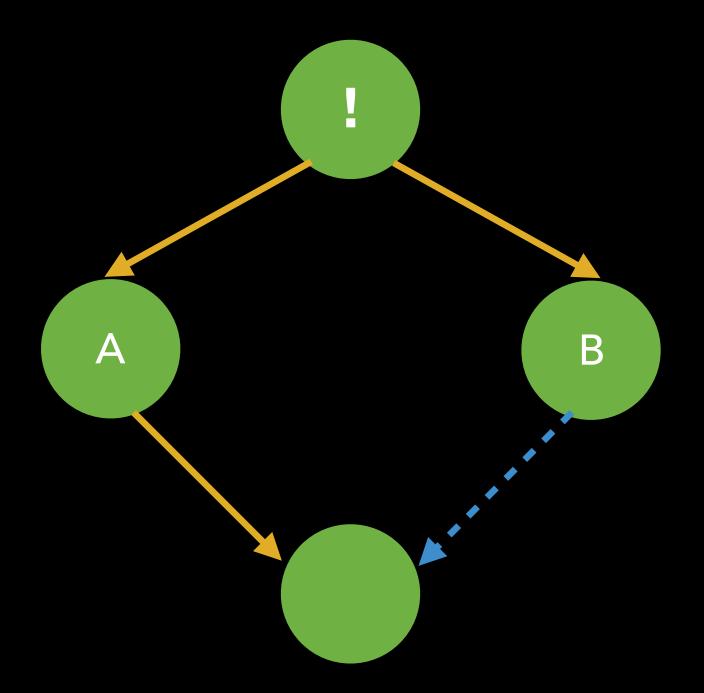


- All nodes start with full "eager" set
- Broadcast triggers eager-push
- Duplicate messages cause "pruning" (move to "lazy")



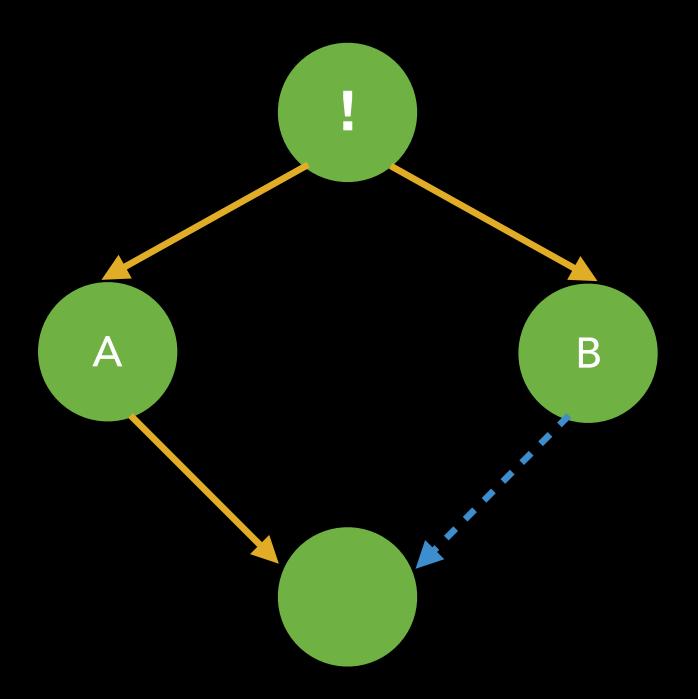


- All nodes start with full "eager" set
- Broadcast triggers eager-push
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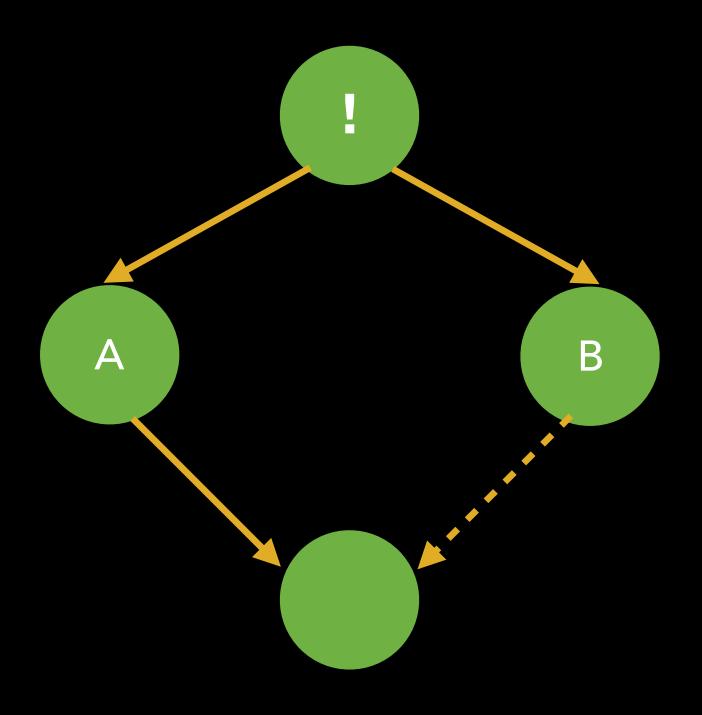




- All nodes start with full "eager" set
- Broadcast triggers eager-push
- Duplicate messages cause "pruning" (move to "lazy")
- Regular broadcasts proceed with new "eager" sets

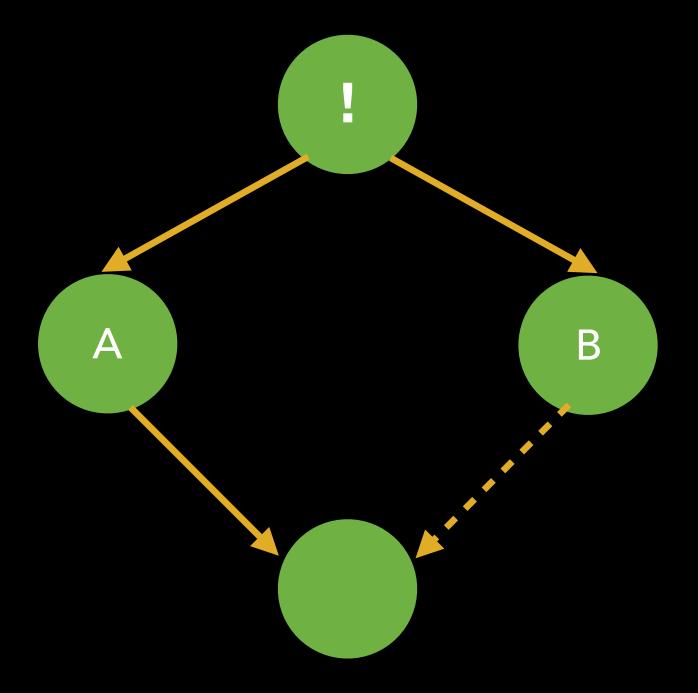






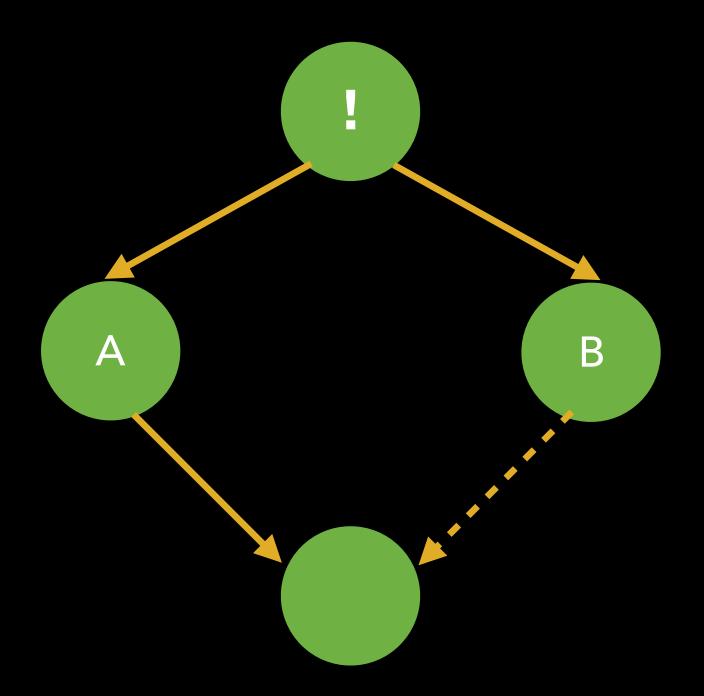


Lazy-push sends "I Have" messages



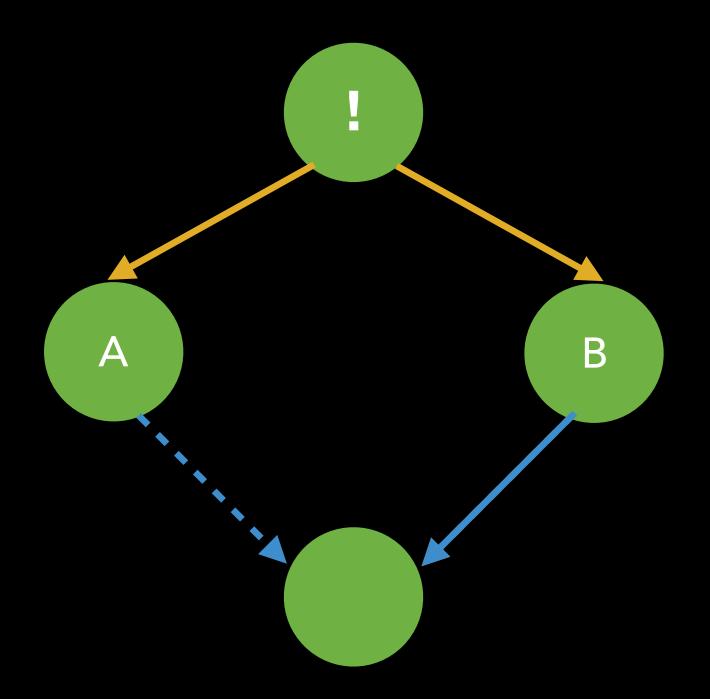


- Lazy-push sends "I Have" messages
- Timeout triggers "grafting" (move to "eager")



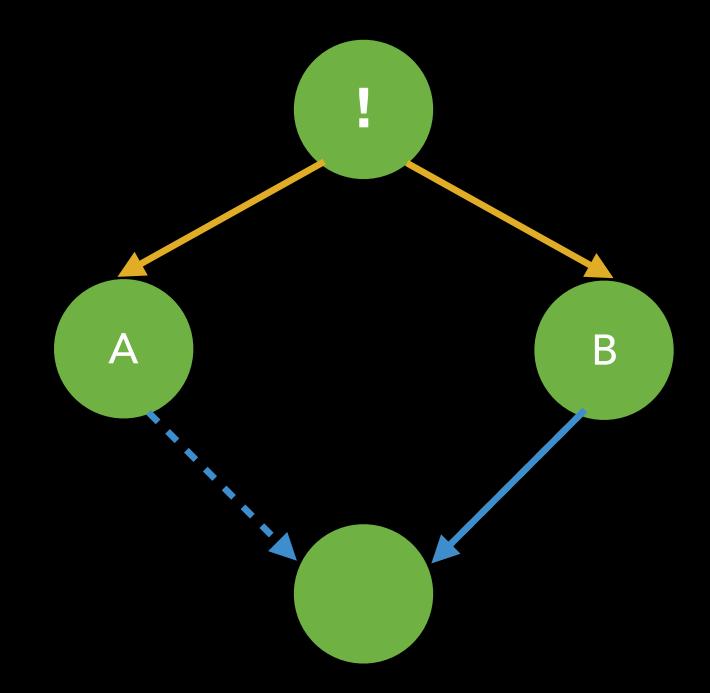


- Lazy-push sends "I Have" messages
- Timeout triggers "grafting" (move to "eager")





- Lazy-push sends "I Have" messages
- Timeout triggers "grafting" (move to "eager")
- Lazy-push batched to reduce overhead







Good tradeoff between reliability and redundancy



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- Optimizes for lowest-latency paths



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- Good tradeoff between reliability and redundancy
- Optimizes for lowest-latency paths
- Existing open-source implementations
- Excellent fit with HyParView





POPULATION PROTOCOLS



RANDOMIZED INTERACTIONS



Logical Physical Clocks and Consistent Snapshots in Globally Distributed Databases

Sandeep Kulkarni*, Murat Demirbas**, Deepak Madeppa**, Bharadwaj Avva**, and Marcelo Leone*

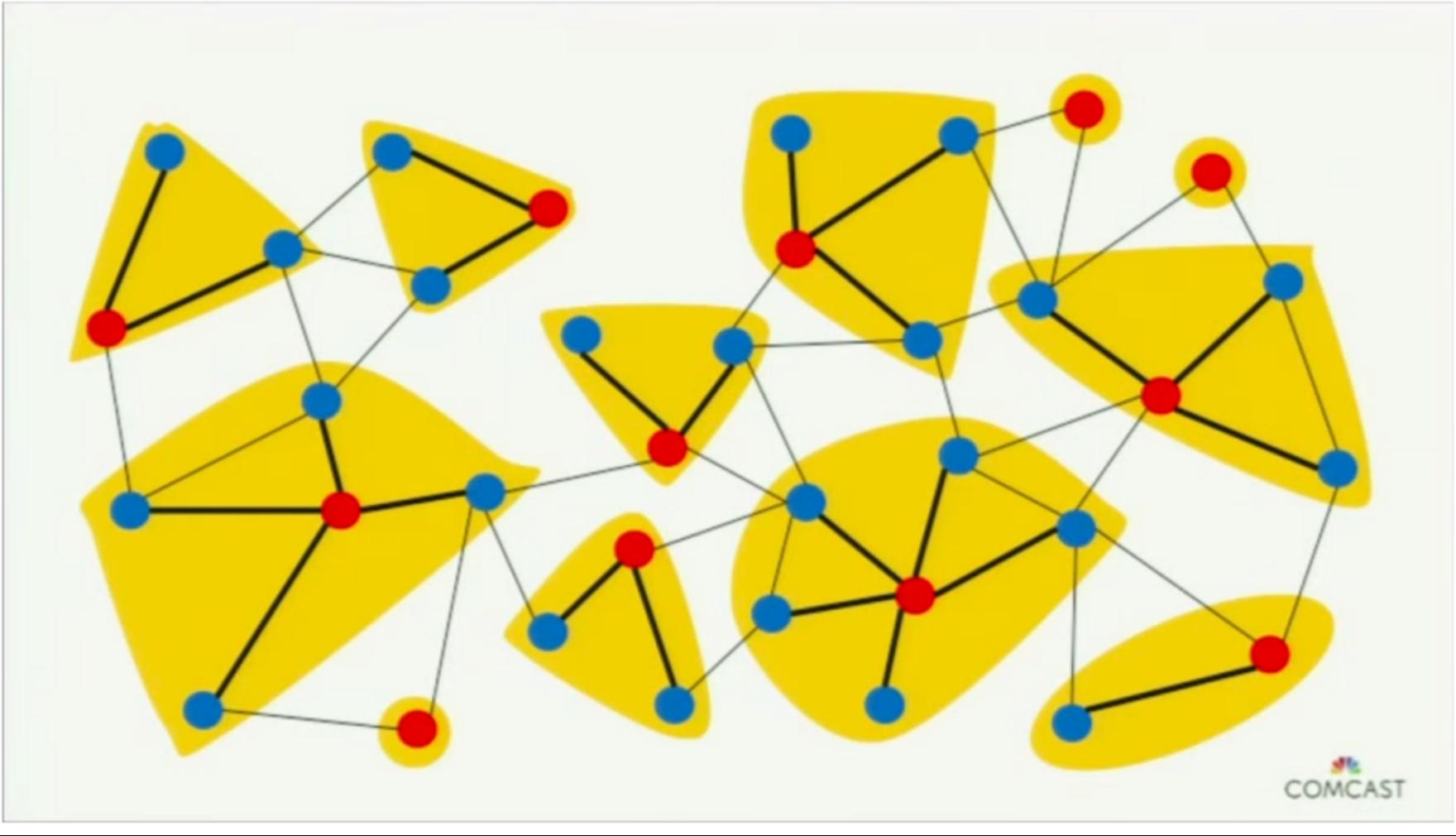
**Michigan State University
**University at Buffalo, SUNY



DISTRIBUTED MONOTONIC CLOCKS







Vidcap from StrangeLoop 2015: https://youtu.be/YqNGbvFHoKM

DMC PROBLEMS

- "Wacky clock mode"
- Hierarchy imbalances load
- Long-lived partitions
- No convergence proof



Bridging the Gap between Population and Gossip-based Protocols

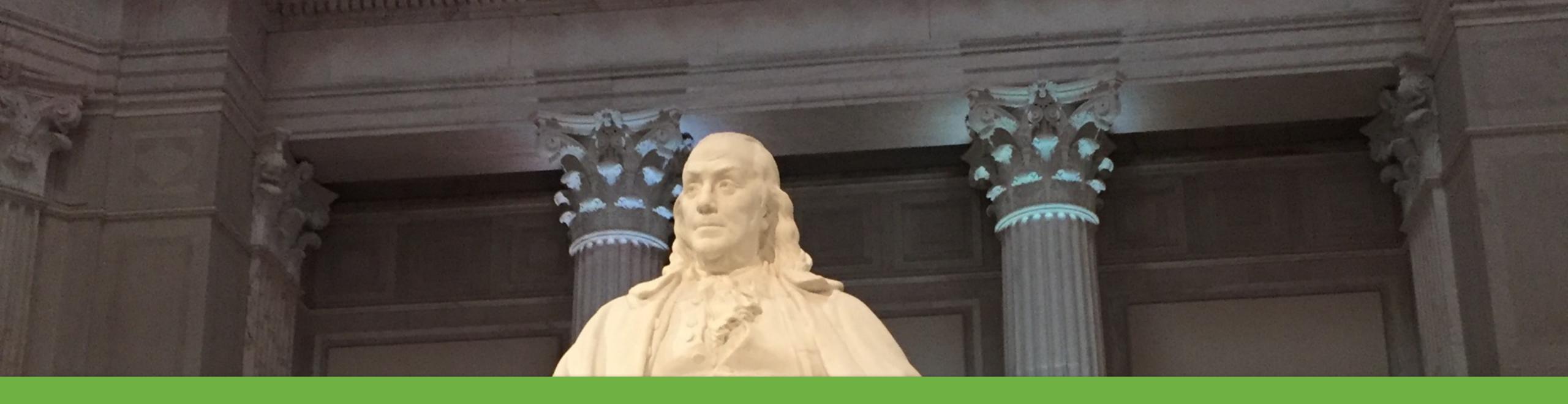
Yann Busnel, Marin Bertier, Anne-Marie Kermarrec



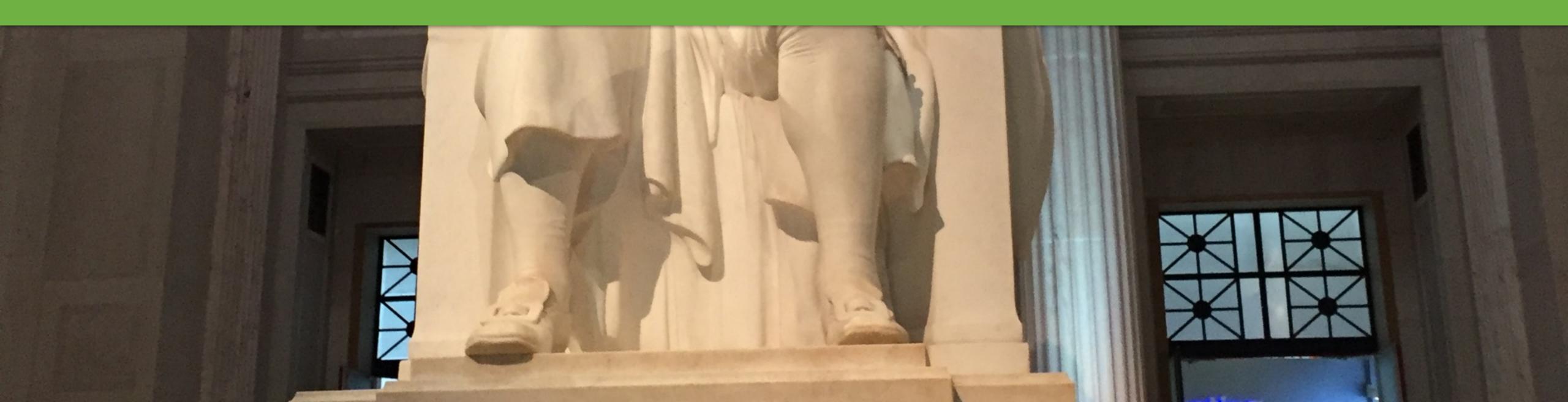
APPLYING DMC

- Use existing dissemination with DMC
- Transmit clocks along with other messages
- Use monotonic clocks as a drift-detection mechanism





LESSONS LEARNED





THANK YOU!

@ S E A N C R I B B S

